

Surface Code Quantum Error Correction

Rodney Van Meter, Keio University

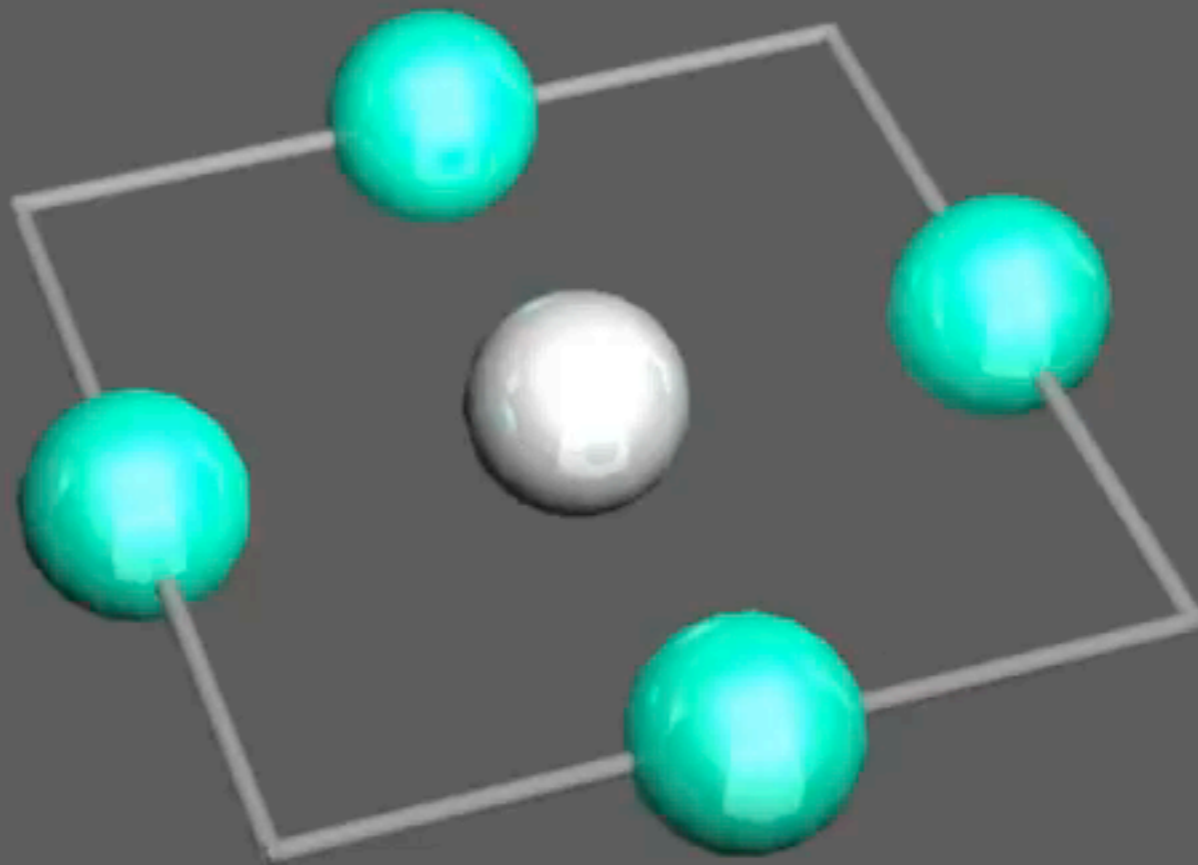
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FIRST Project Summer School @ Kyoto U.

2011 Aug 16

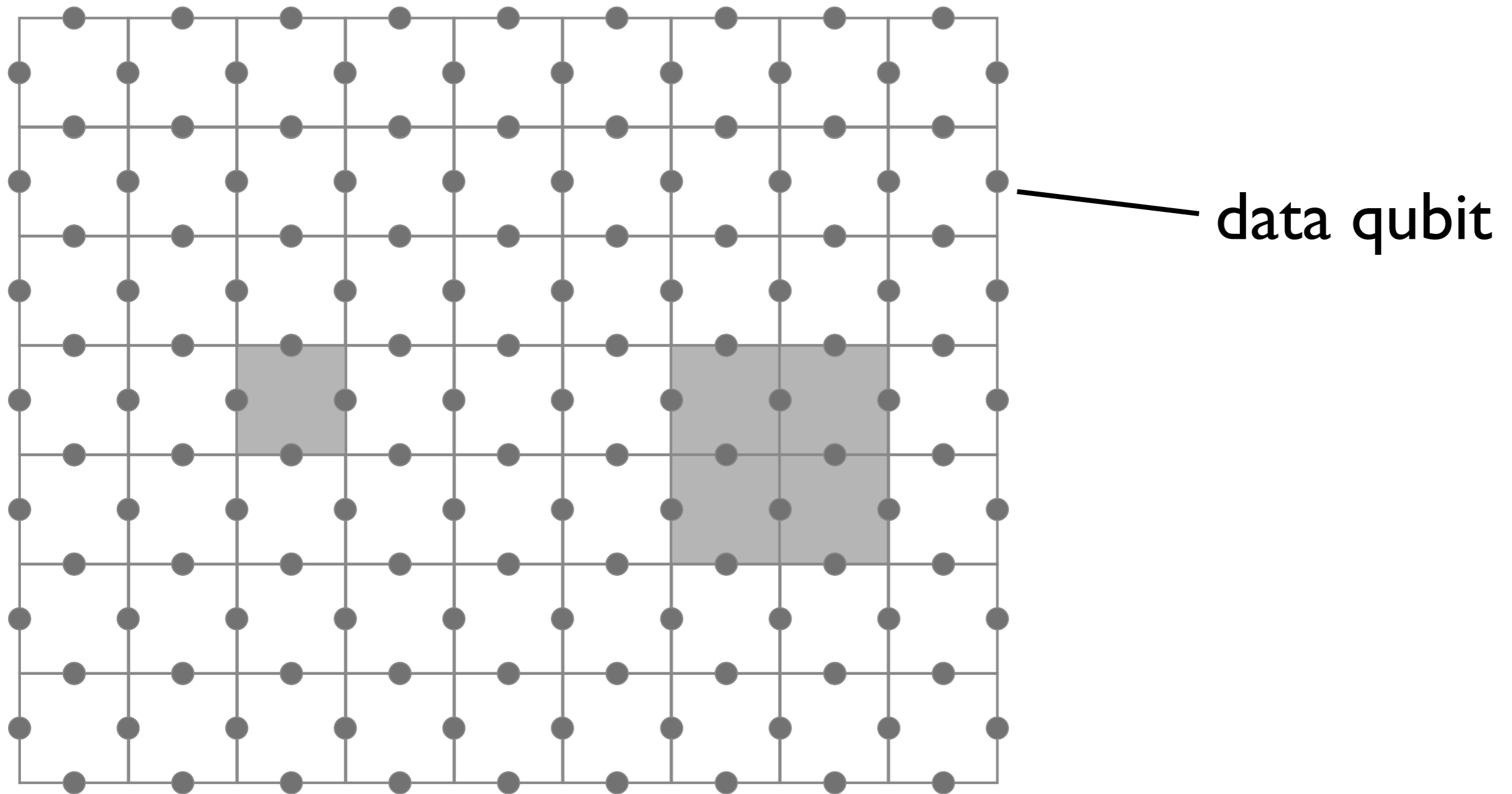
(all the good slides and animations are by
Shota Nagayama, Austin Fowler and Clare Horsman)





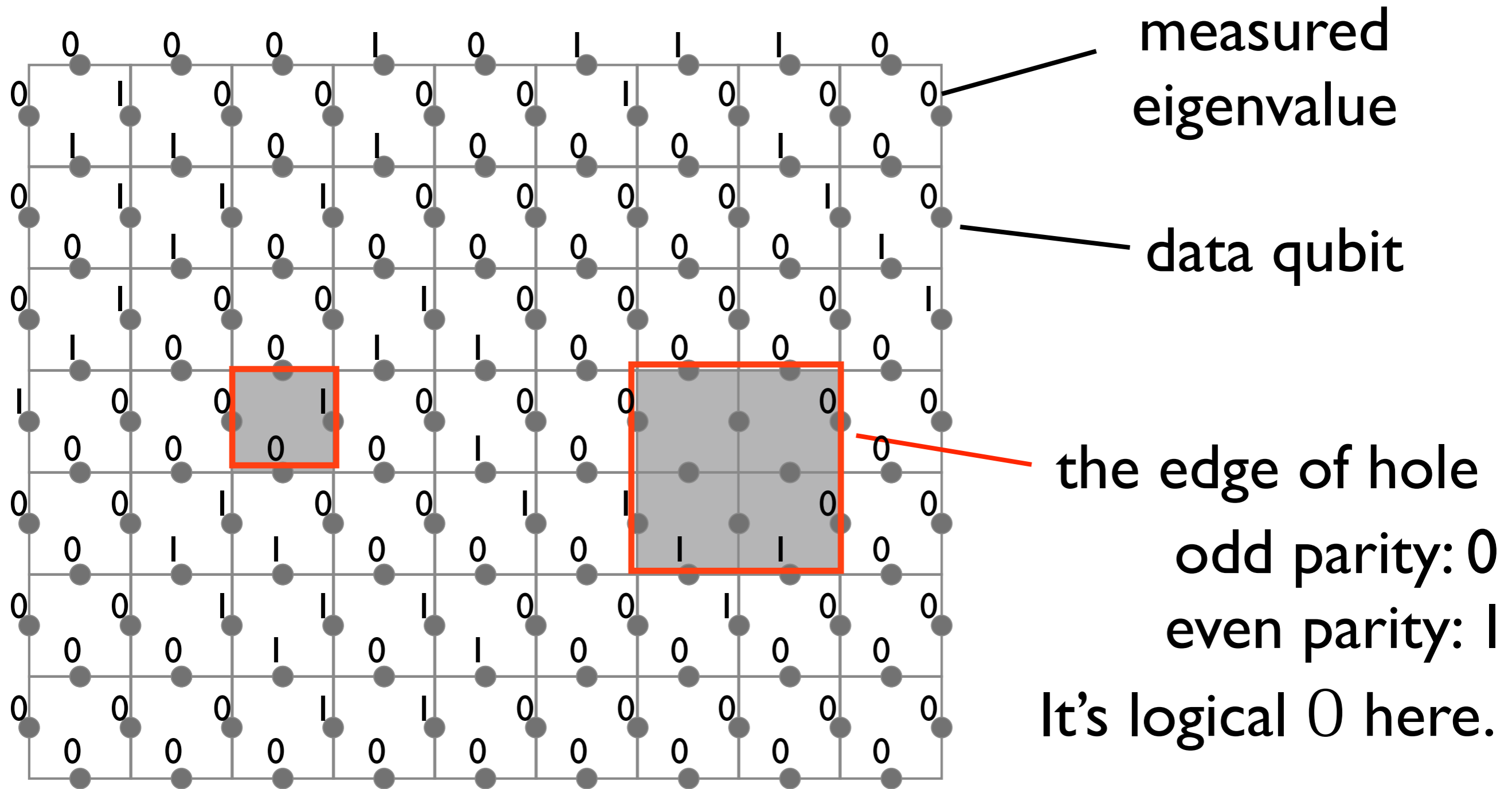
Simplest Example

~Encoding logical qubit~



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Surface Code Strengths



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- Simple, 2-D or 3-D nearest-neighbor-only operation (physical feasibility high!)





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 - strength of EC grows incrementally (compare to concatenated CSS codes)
 - software-assigned resources
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- High threshold: 1.4% for gate, memory, measurement errors
- Flexible:
 - strength of EC grows incrementally (compare to concatenated CSS codes)
 - software-assigned resources
 - easy movement of logical qubits
- Supporting classical processing achievable





Scalability: fault-tolerance

- Trade-off between resources and threshold
- Thresholds
 - unlimited range, unlimited qubits: $\sim 10^{-2}$
Knill, quant-ph/0410199
 - unlimited range, many qubits: $\sim 10^{-3}$ – 10^{-4}
Steane, Phys. Rev. A 68, 042322 (2003)
 - 2D lattice, nearest neighbor: $\sim 10^{-5}$
Svore, QIC 7, 297 (2007)
 - bilinear nearest neighbor: $\sim 10^{-6}$
Stephens, QIC 8, 330 (2008)
 - linear nearest neighbor: $\sim 10^{-8}$
Stephens, in preparation





Surface Code Drawbacks





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- Non-Clifford group operations difficult





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Surface Code Drawbacks

- Non-Clifford group operations difficult
- Direct calculation of residual error rates difficult due to many error chains; determined via simulation
- Almost uniform set of operations across whole device, but not quite!
- *Extremely* difficult to explain to classical computer engineers!





Outline

- Stabilizers
- Surface Code Operation
- Theory of the Surface Code
- Advanced topics
- (system architecture reserved for next lecture)





Stabilizers

$$|0\rangle \equiv \begin{pmatrix} 1 \\ 0 \end{pmatrix} \equiv +1 \text{ eigenstate of } Z = \begin{pmatrix} 1 & 0 \\ 0 & -1 \end{pmatrix}$$

$$|+\rangle \equiv \frac{1}{\sqrt{2}} \begin{pmatrix} 1 \\ 1 \end{pmatrix} \equiv +1 \text{ eigenstate of } X = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}$$

$$|0\rangle \quad \text{or} \quad Z$$

$$|+\rangle \quad \text{or} \quad X$$

$$\frac{1}{\sqrt{2}}(|00\rangle + |11\rangle) \quad \text{or} \quad Z_1 Z_2, X_1 X_2$$

- n qubits, n independent commuting stabilizers \Rightarrow unique state





Stabilizers

- $n-k$ stabilizers on set of n qubits leaves k degrees of freedom, can encode k logical qubits
- Measure set of stabilizers to get error syndrome
- See Gottesman PhD. thesis, and the set of notes by Clare Horsman in download materials

Introduction to stabilizer theory

June 16, 2011

1 Stabilizer definition

In quantum mechanics, a **state** is given by a **vector**, and an **operator** is given by a **matrix**. The state of N qubits is a 1×2^N vector, and an operation on the state is a $2^N \times 2^N$ matrix.

For example, a state of 1 qubit could be

$$|+\rangle = \frac{1}{\sqrt{2}} \begin{pmatrix} 1 \\ 1 \end{pmatrix}$$

and an operation on it could be

$$X = \sigma_x = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}$$

In linear algebra, if for a matrix M there is a **vector** V and a **scalar** v such that

$$M.V = v.V$$

then V is an **eigenvector** of M and v is the associated **eigenvalue**.

The vector V can be the eigenvector of more than one matrix. The vector V can be **fully defined** by the set $\{M, v\}$ of matrices that V is an eigenvector of.

The **stabilizers** of a state $|\psi\rangle$ are the operators $\frac{1}{v}.M$ where $|\psi\rangle$ is an eigenvector of M with eigenvalue v .

We see that

$$X|+\rangle = \frac{1}{\sqrt{2}} \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} \begin{pmatrix} 1 \\ 1 \end{pmatrix} = \frac{1}{\sqrt{2}} \begin{pmatrix} 1 \\ 1 \end{pmatrix} = |+\rangle$$

so X is a stabilizer of $|+\rangle$.

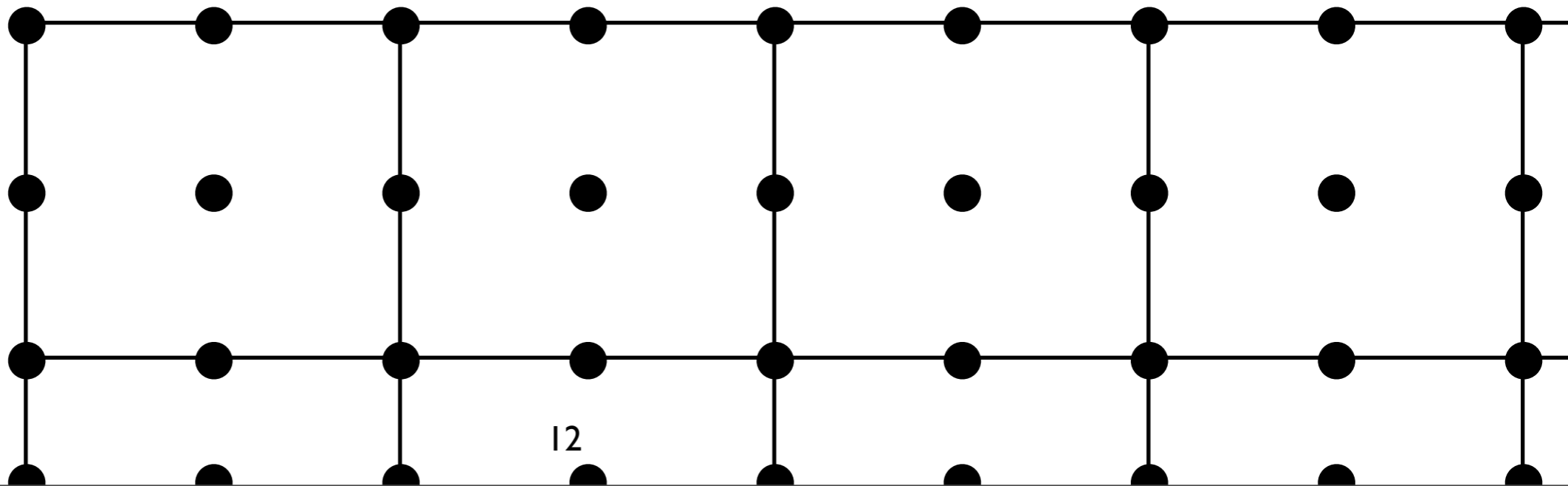


Operations

- Lattice cycle
- Detecting & correcting errors
- Holes as logical qubits
- Single-qubit X and Z gates
- Moving holes
- Braiding for CNOT (Primal & dual lattice)
- Non-Clifford gates using singular qubits
- “Singular factories”
- Measuring a logical qubit



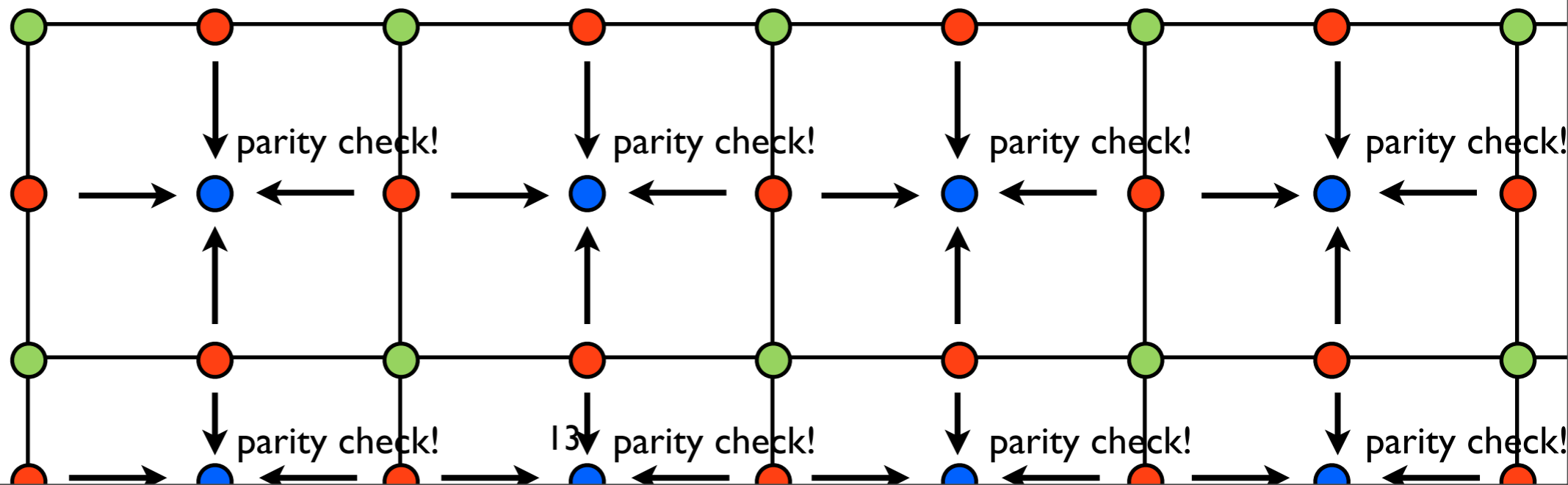
Surface Code



Surface Code

edges of two holes = logical qubit

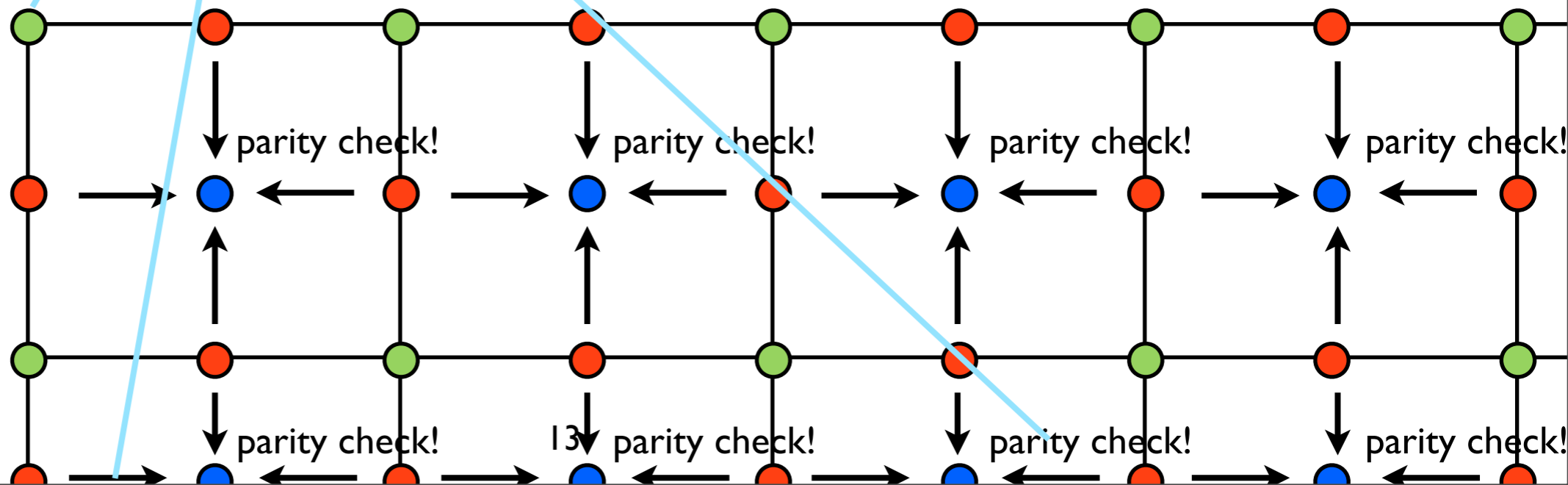
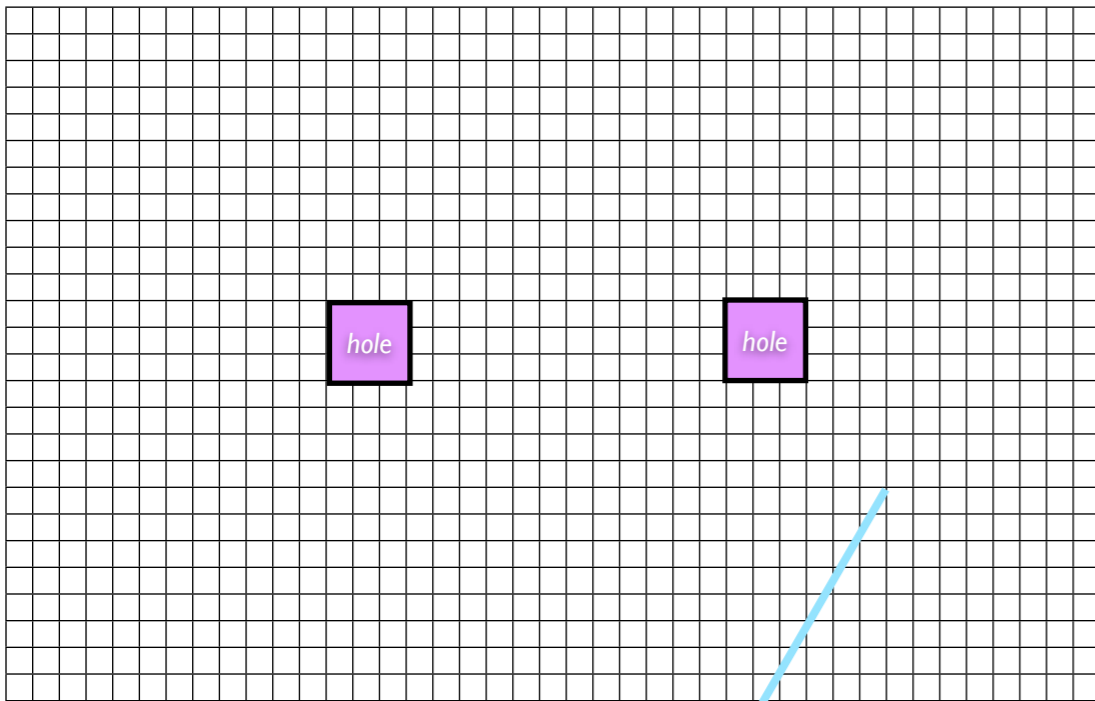
- data qubit
- phase-flip checker
- bit-flip checker



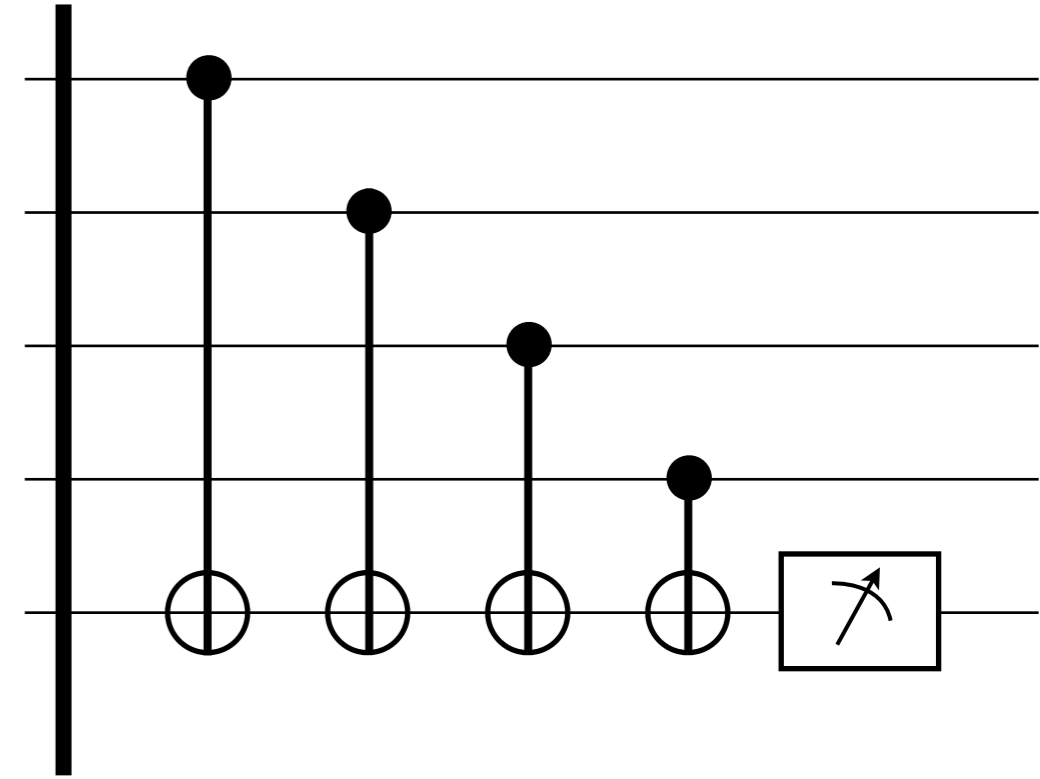
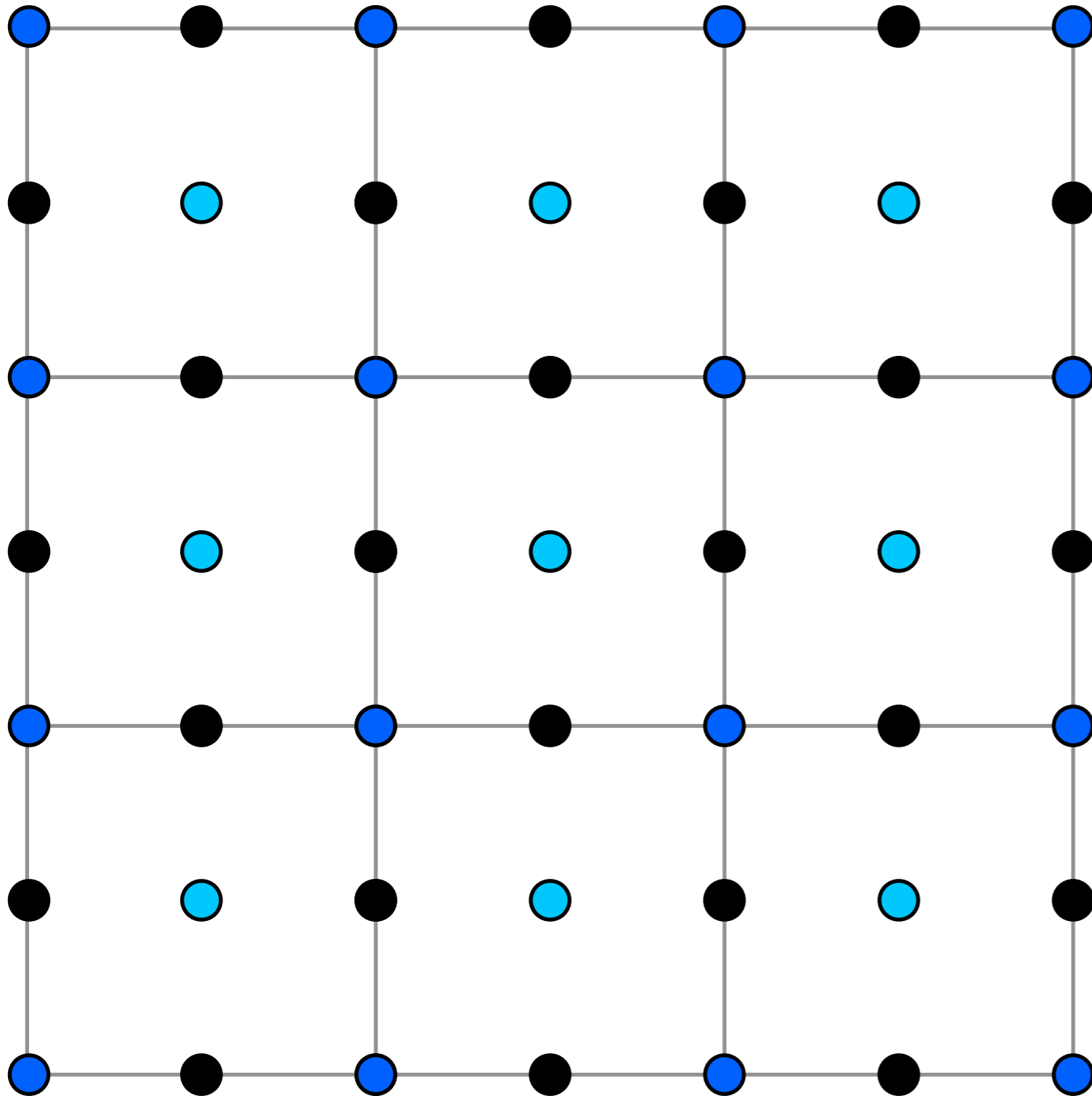
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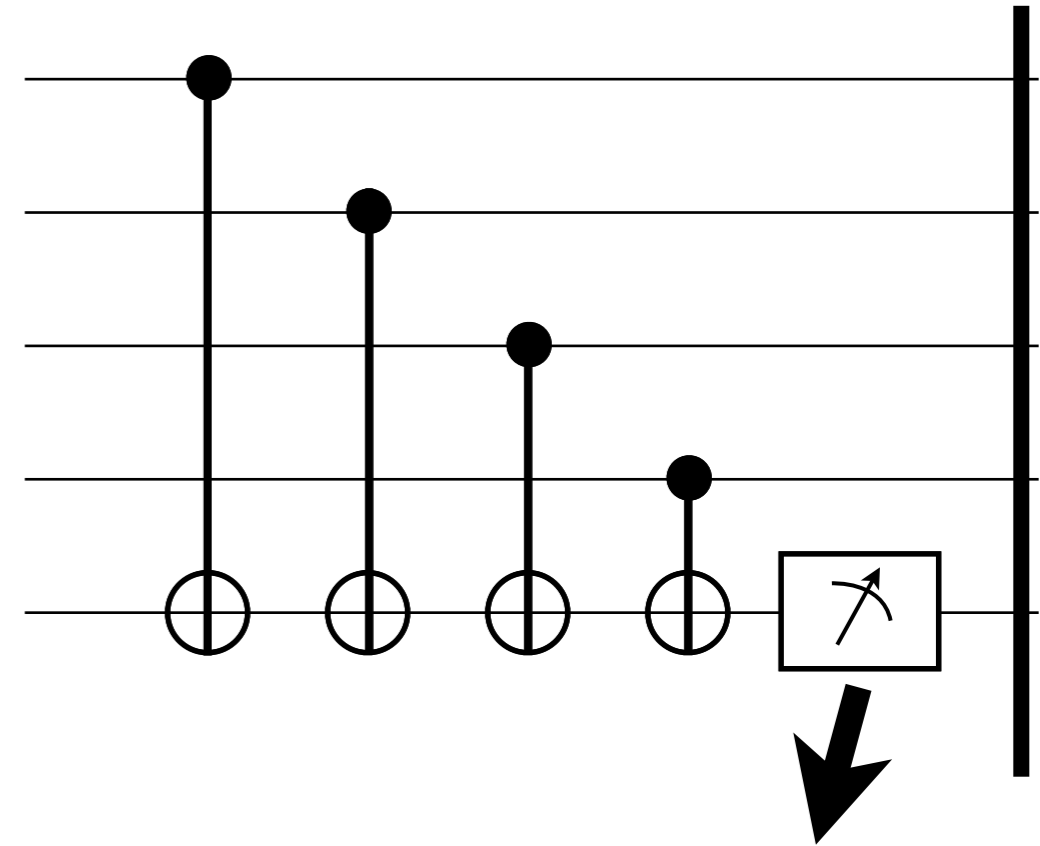
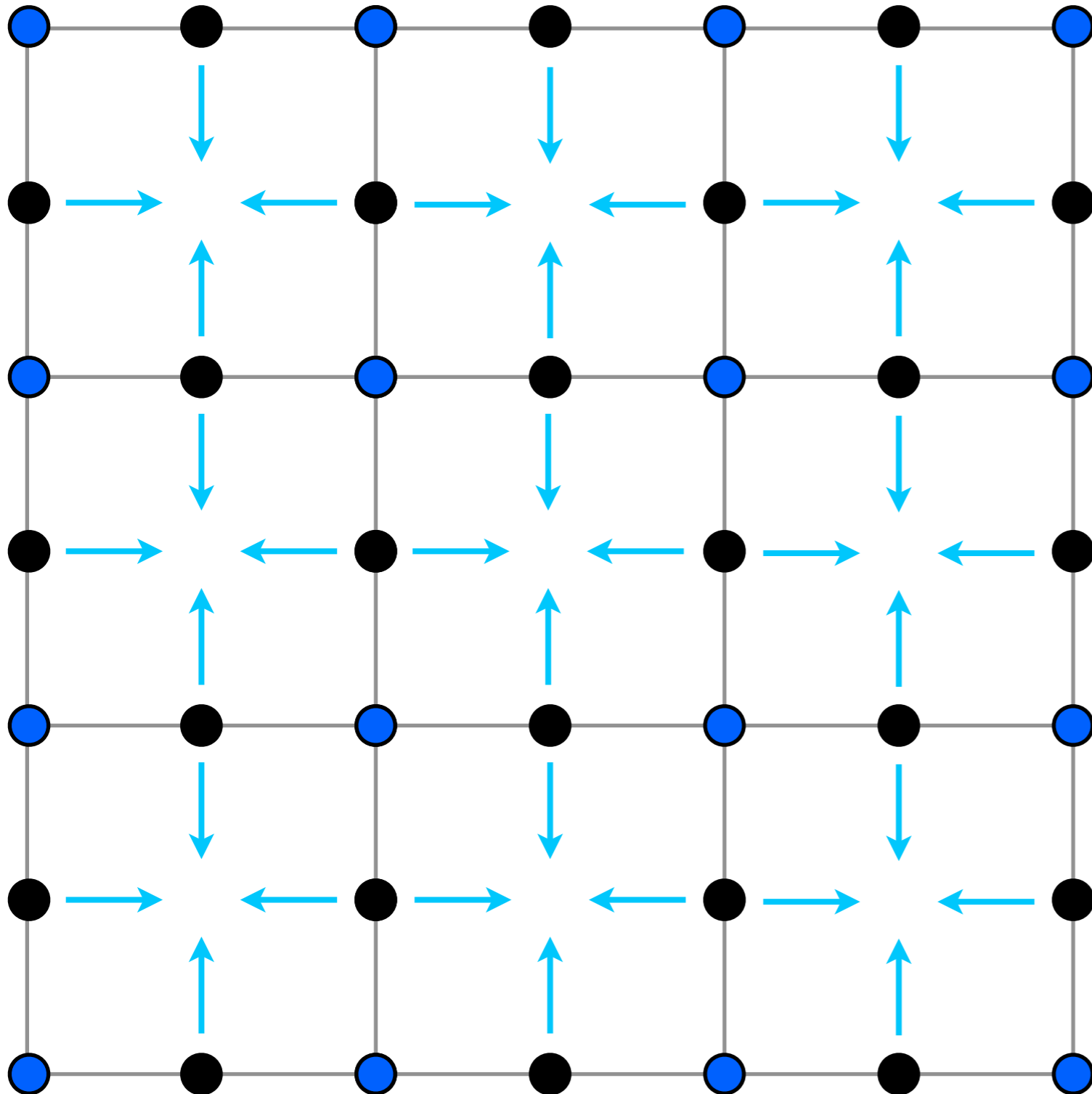
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Entangling the Lattice

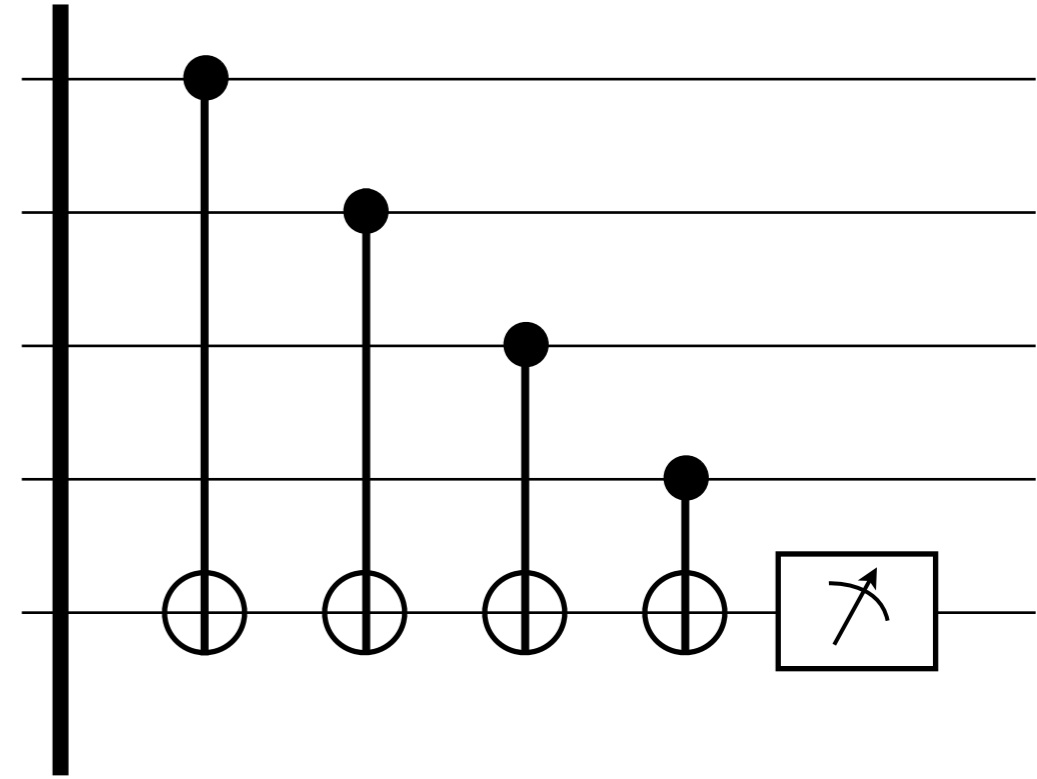
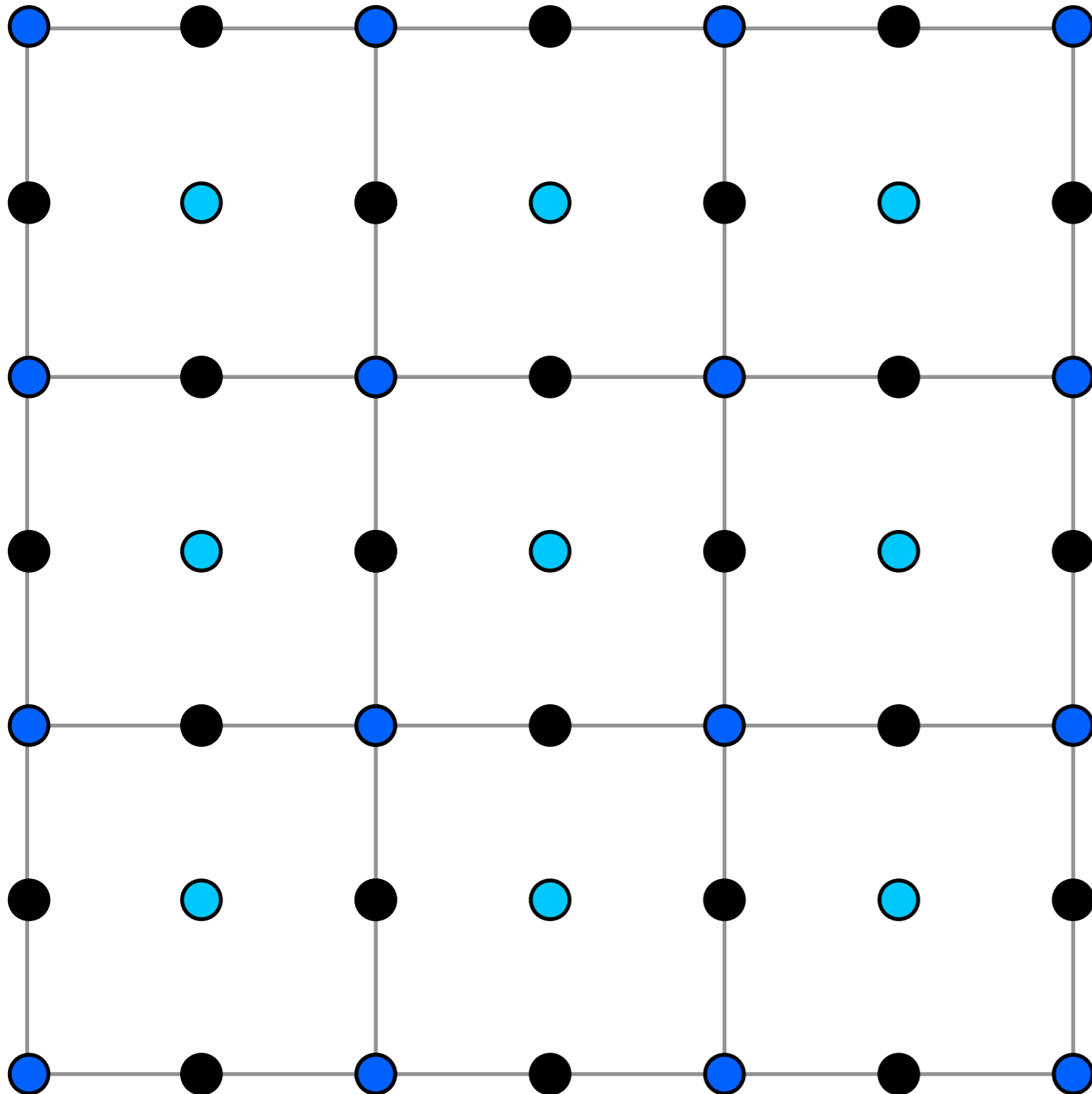


Entangling the Lattice

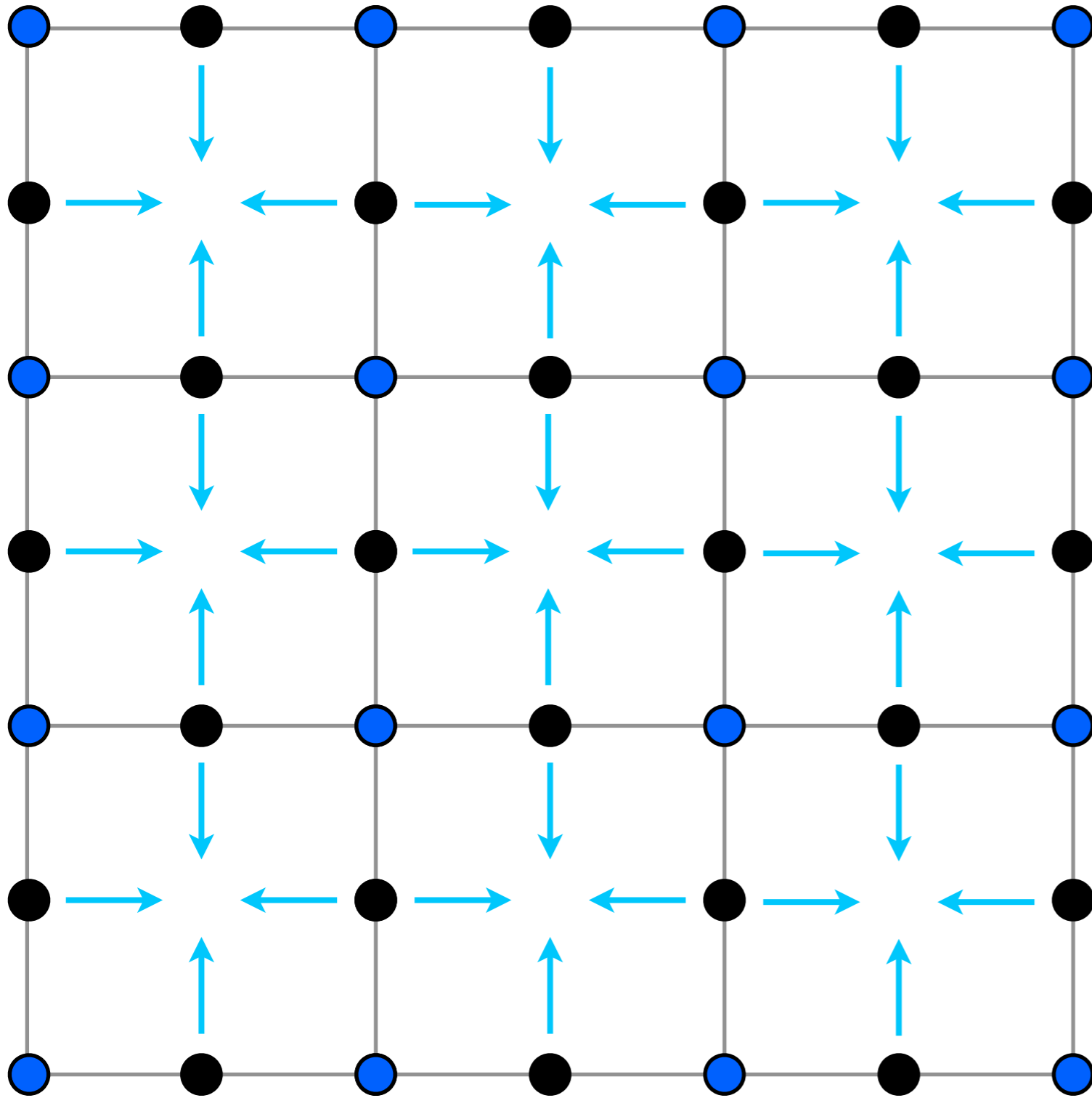


0	0	0
0	0	0
0	0	0

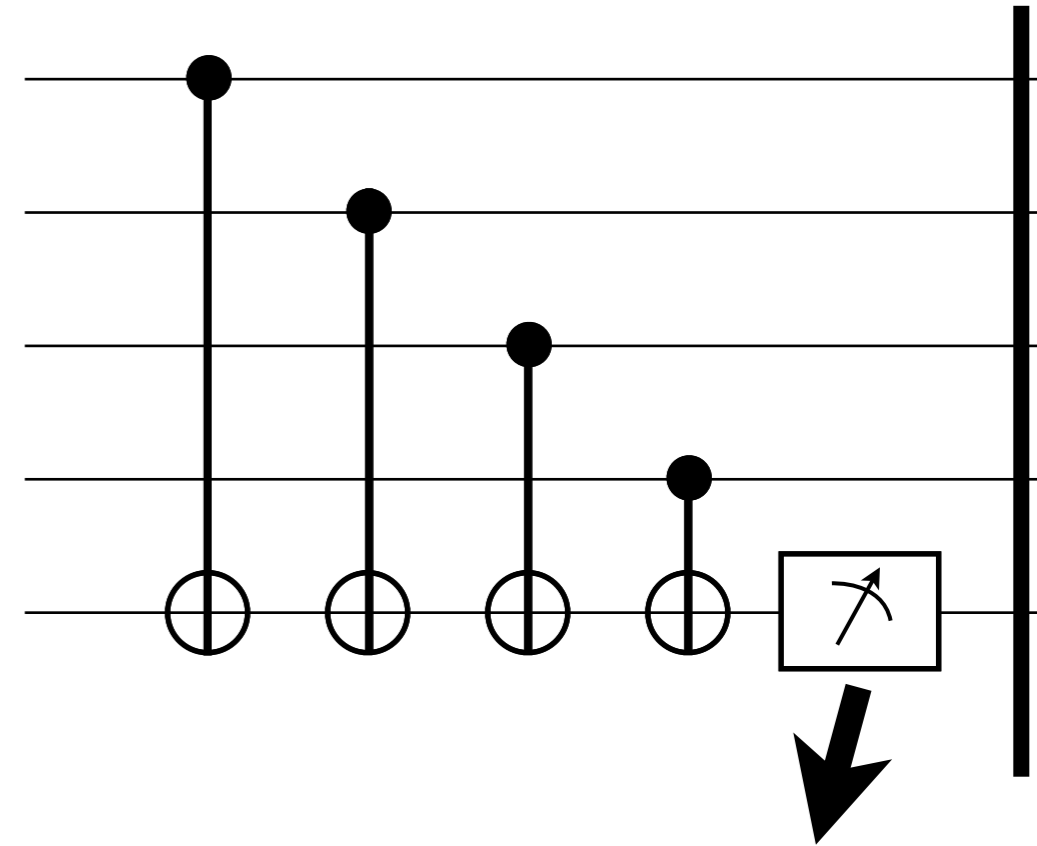
Measuring Z stabilizers



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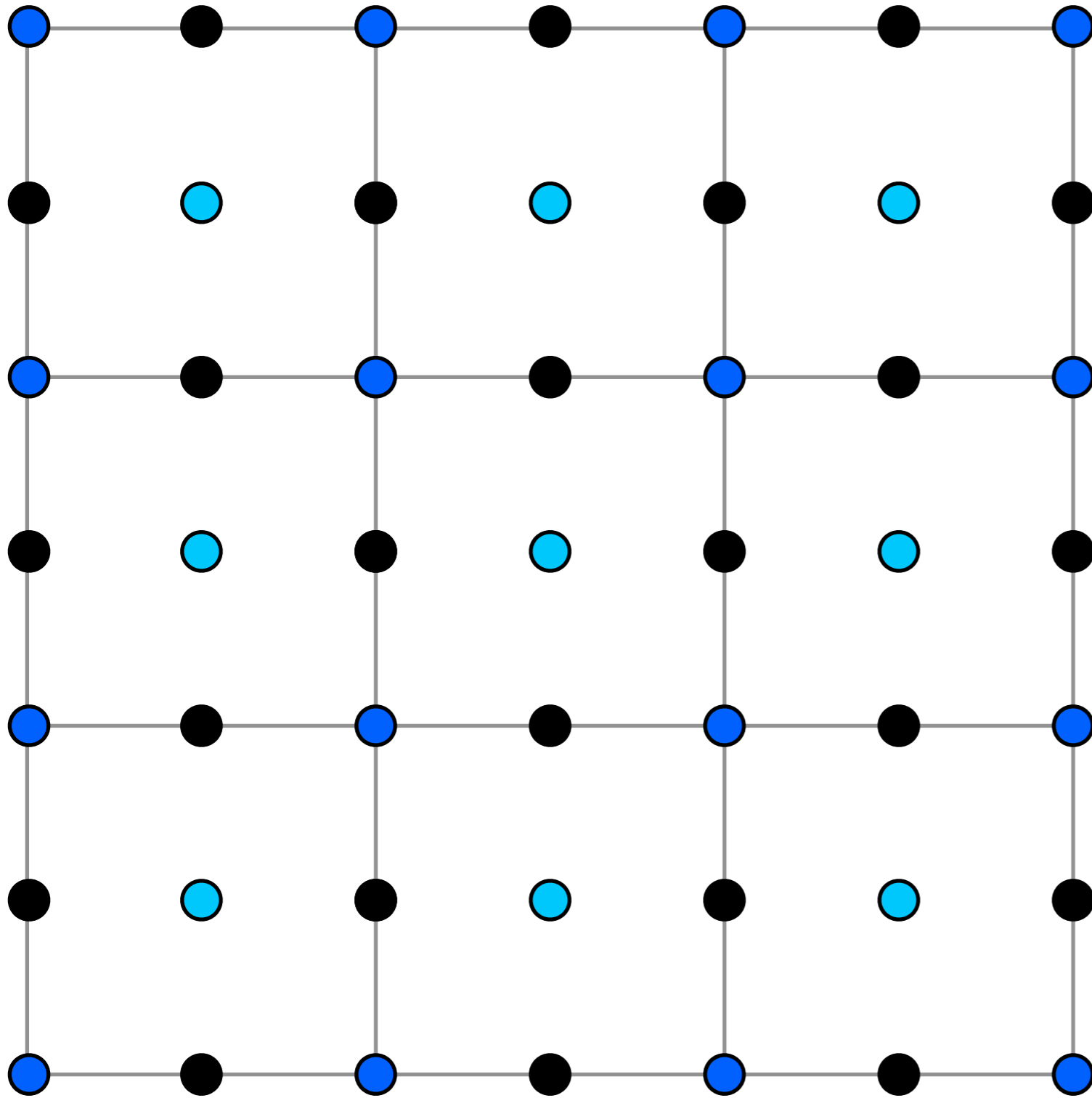


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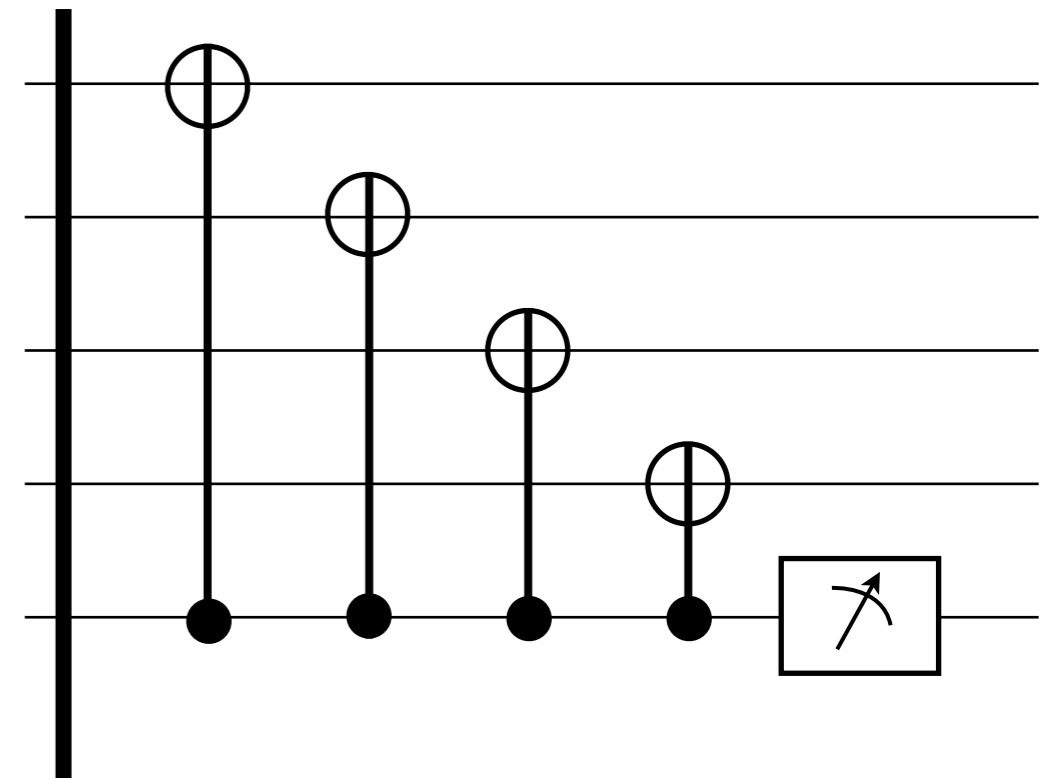


0	0	0
0	0	0
0	0	0

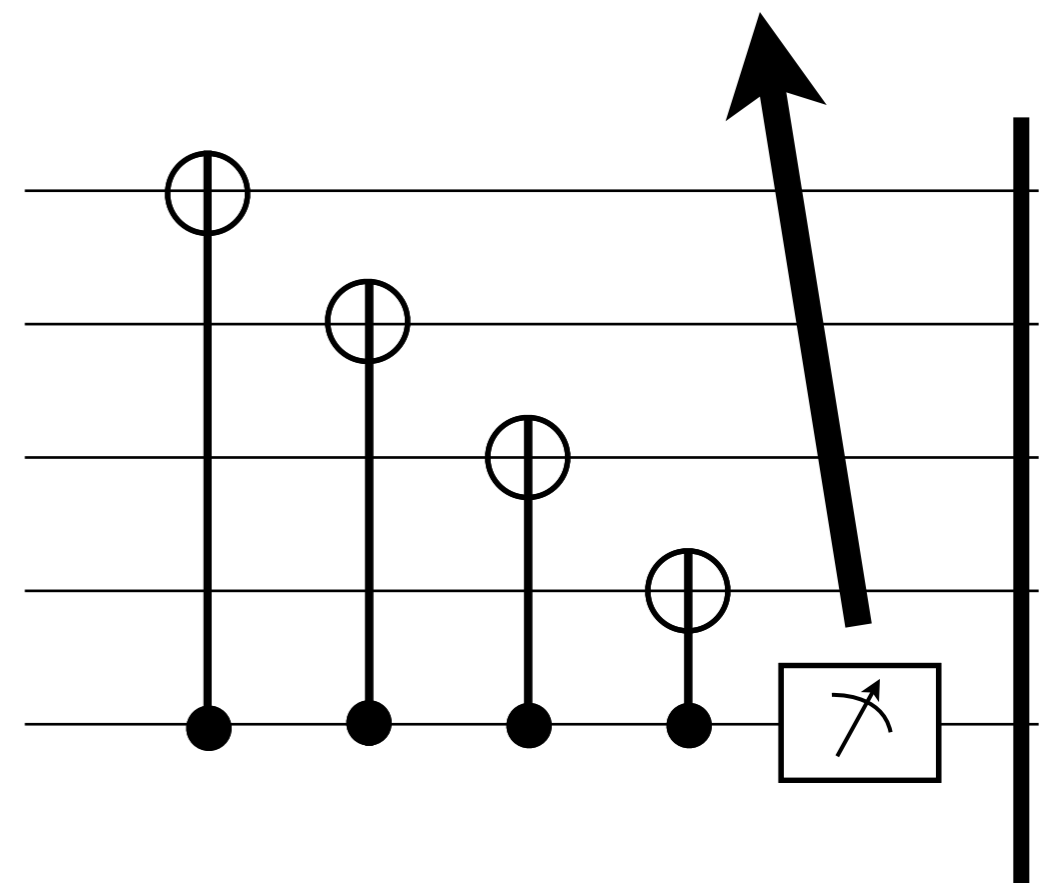
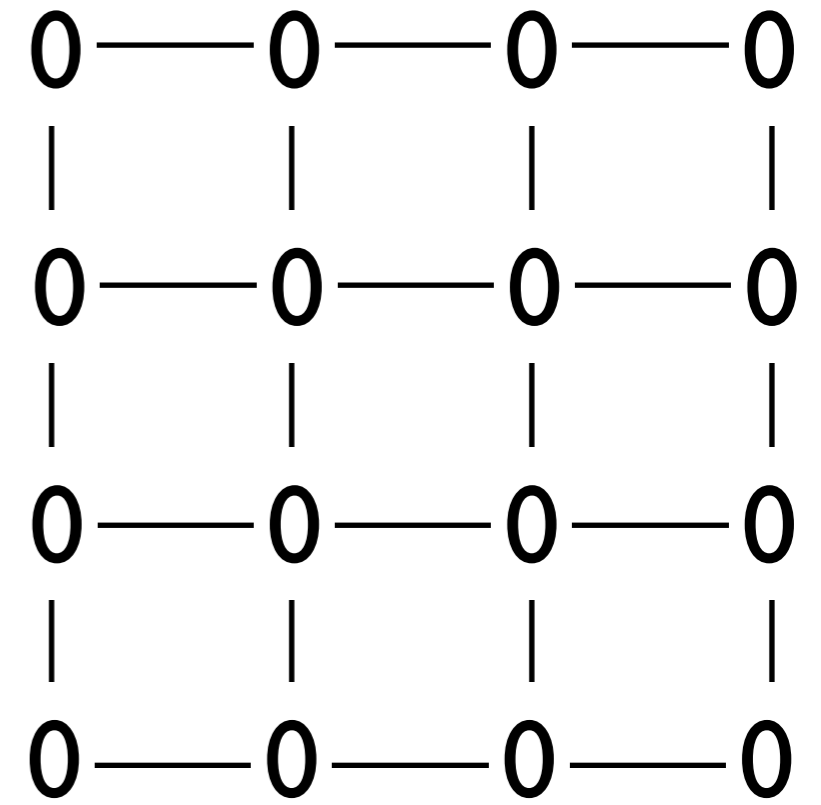
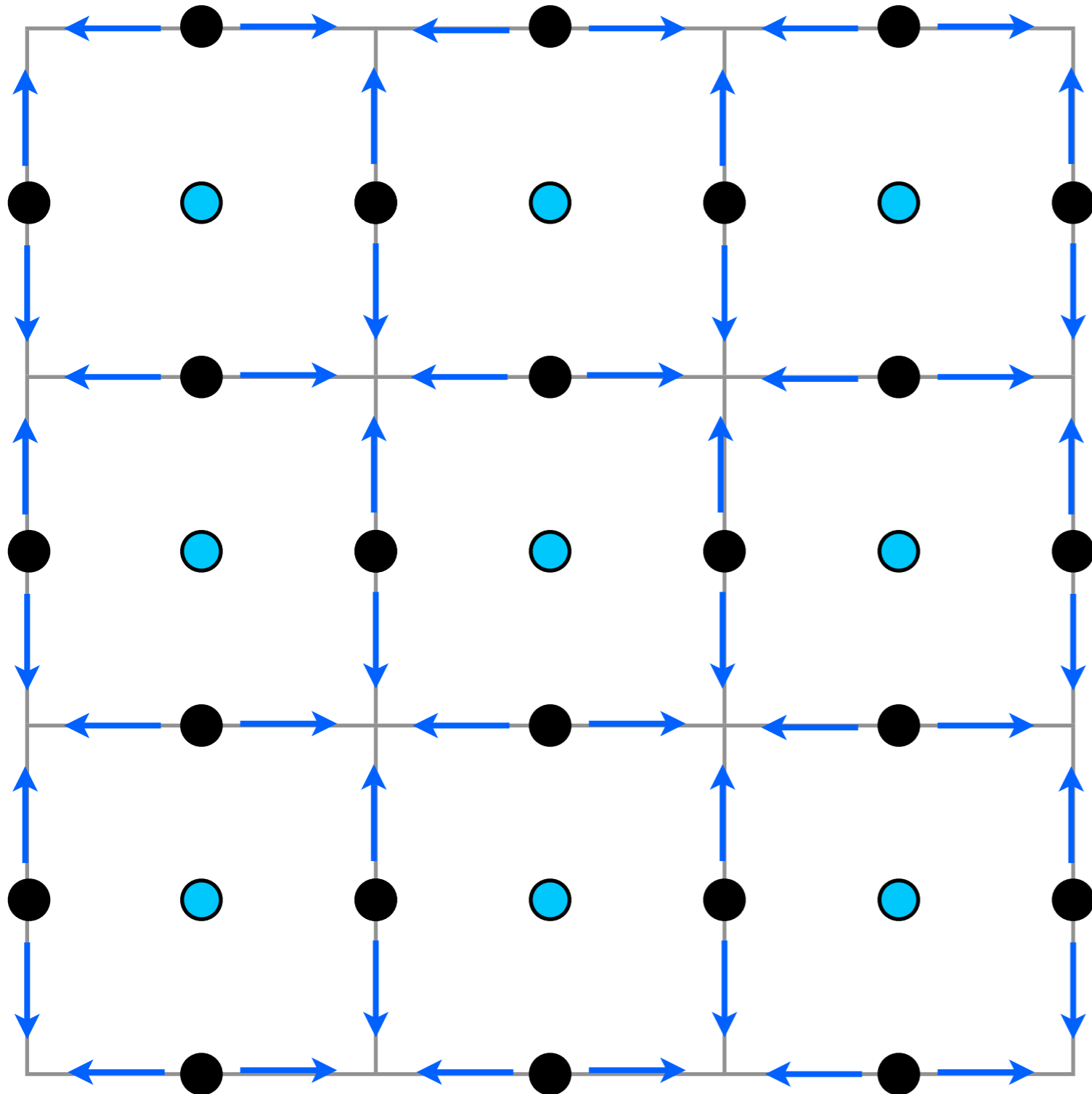
Measuring X stabilizers



16

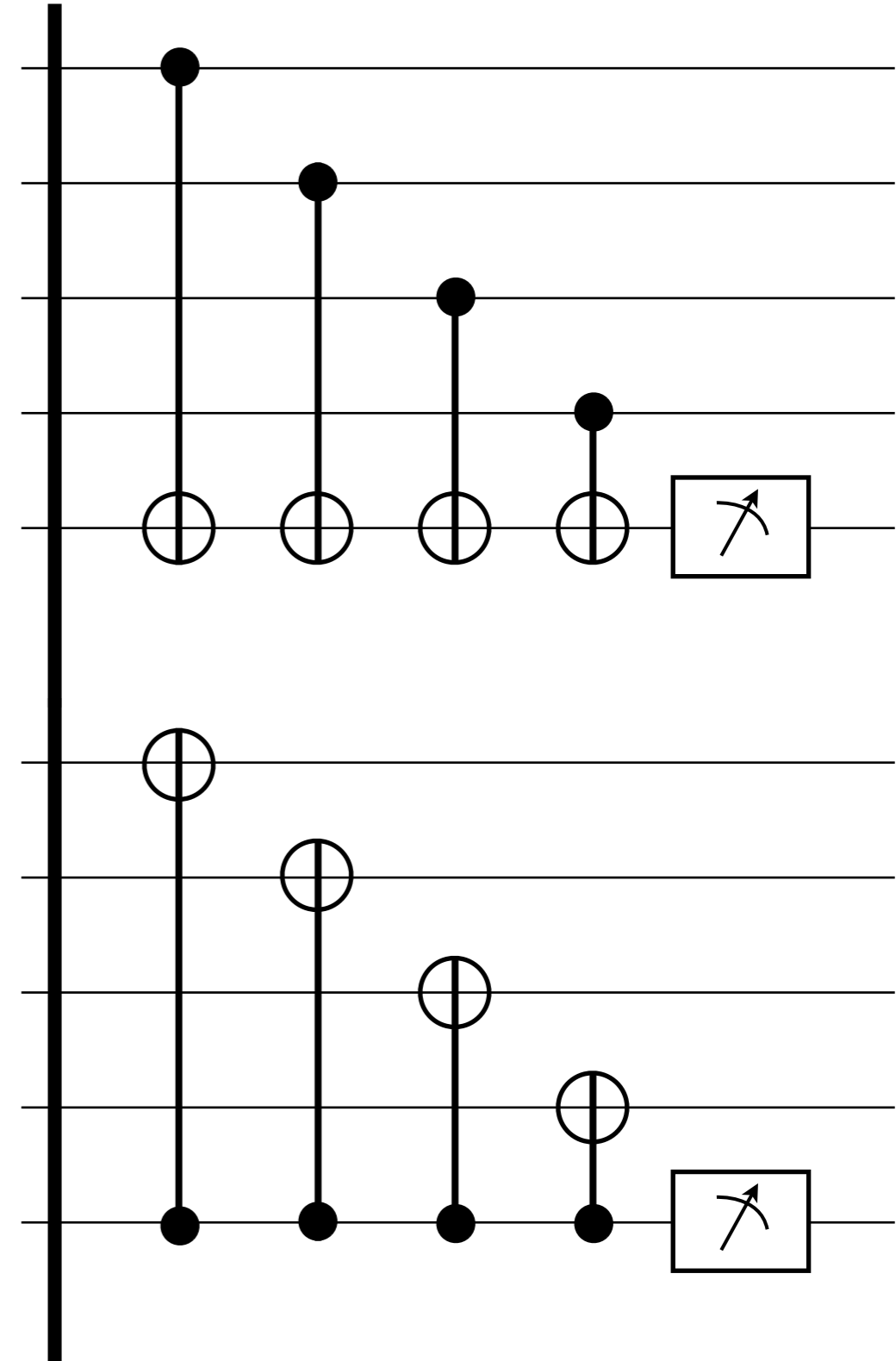
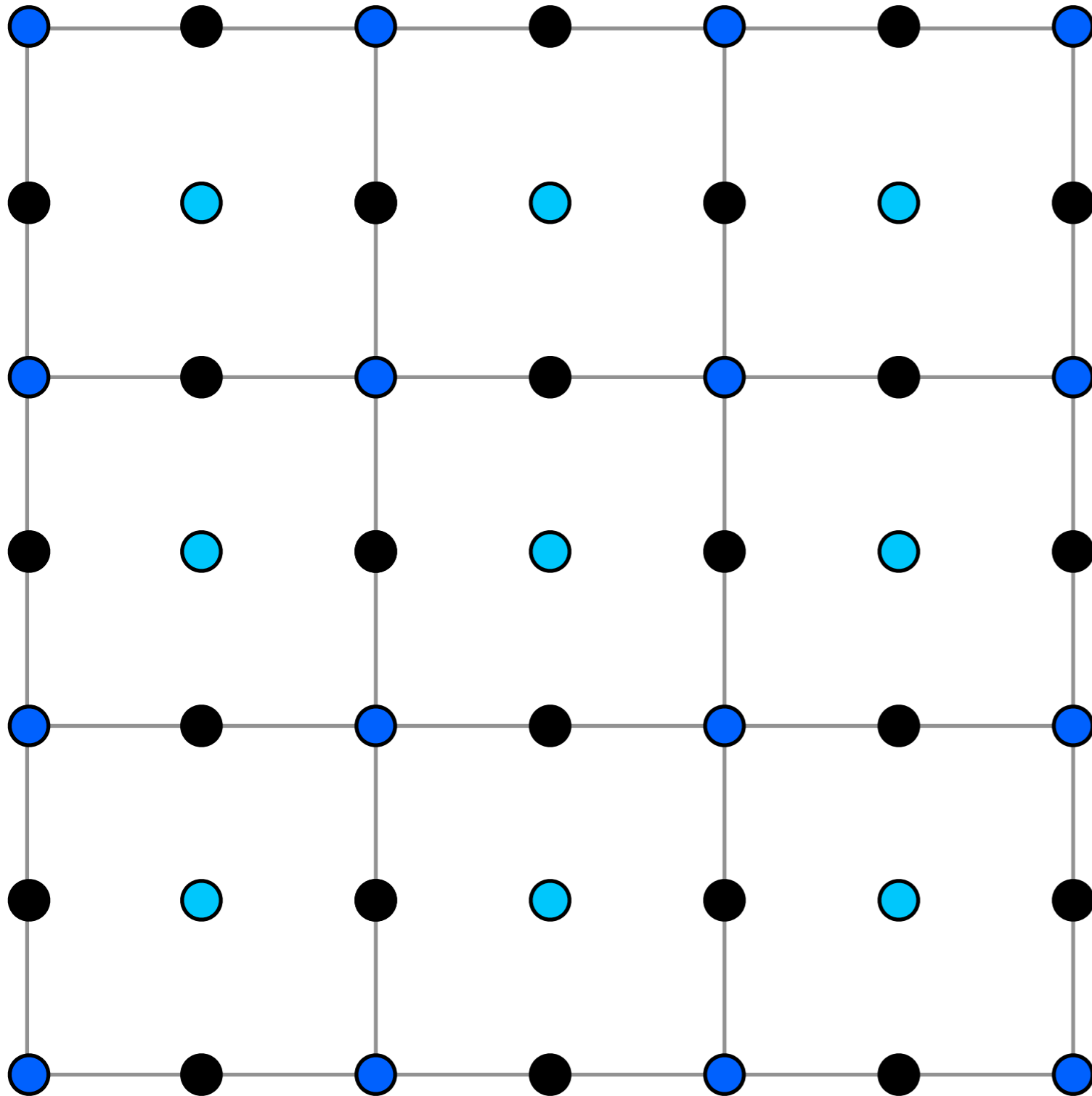


Measuring X stabilizers



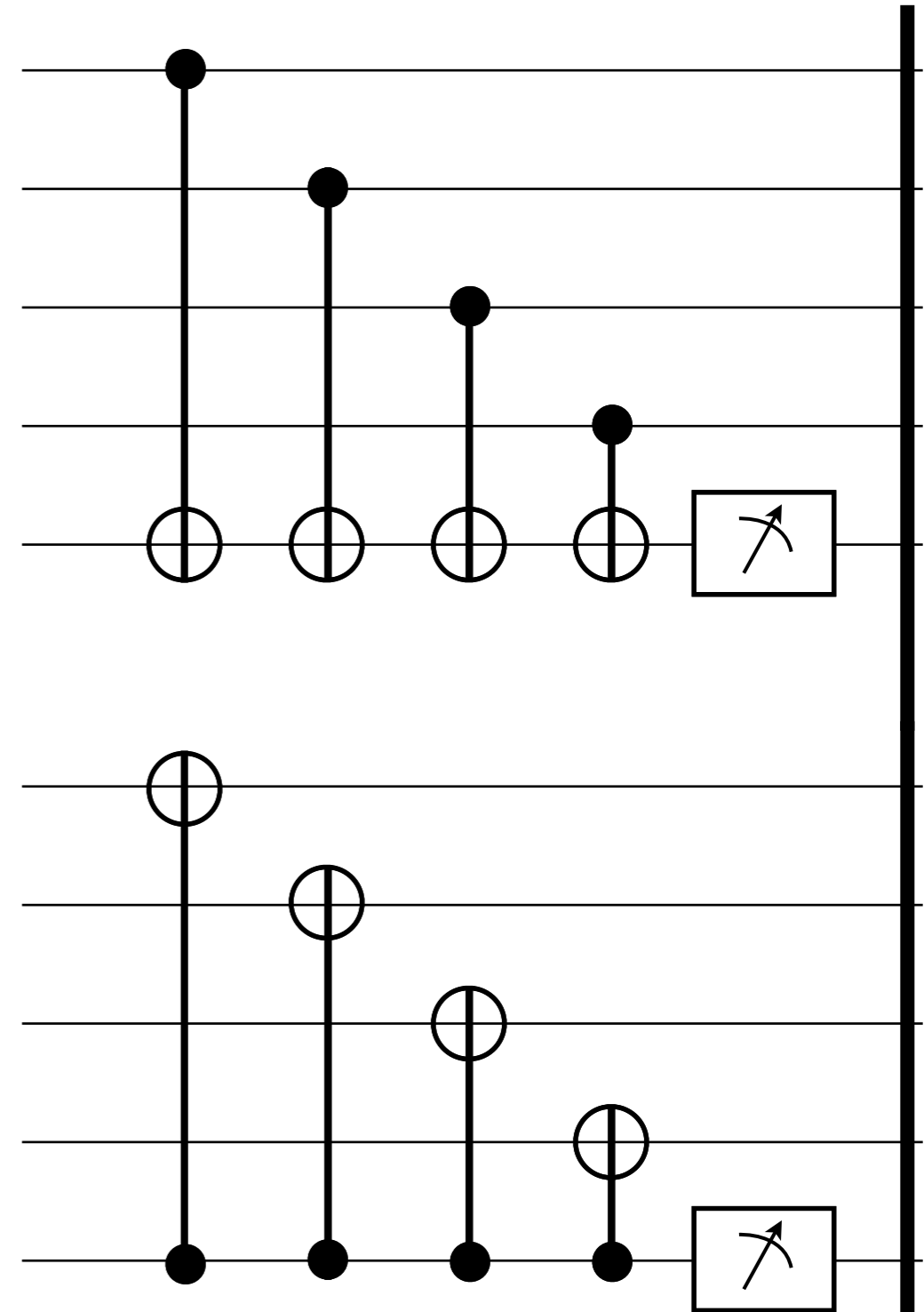
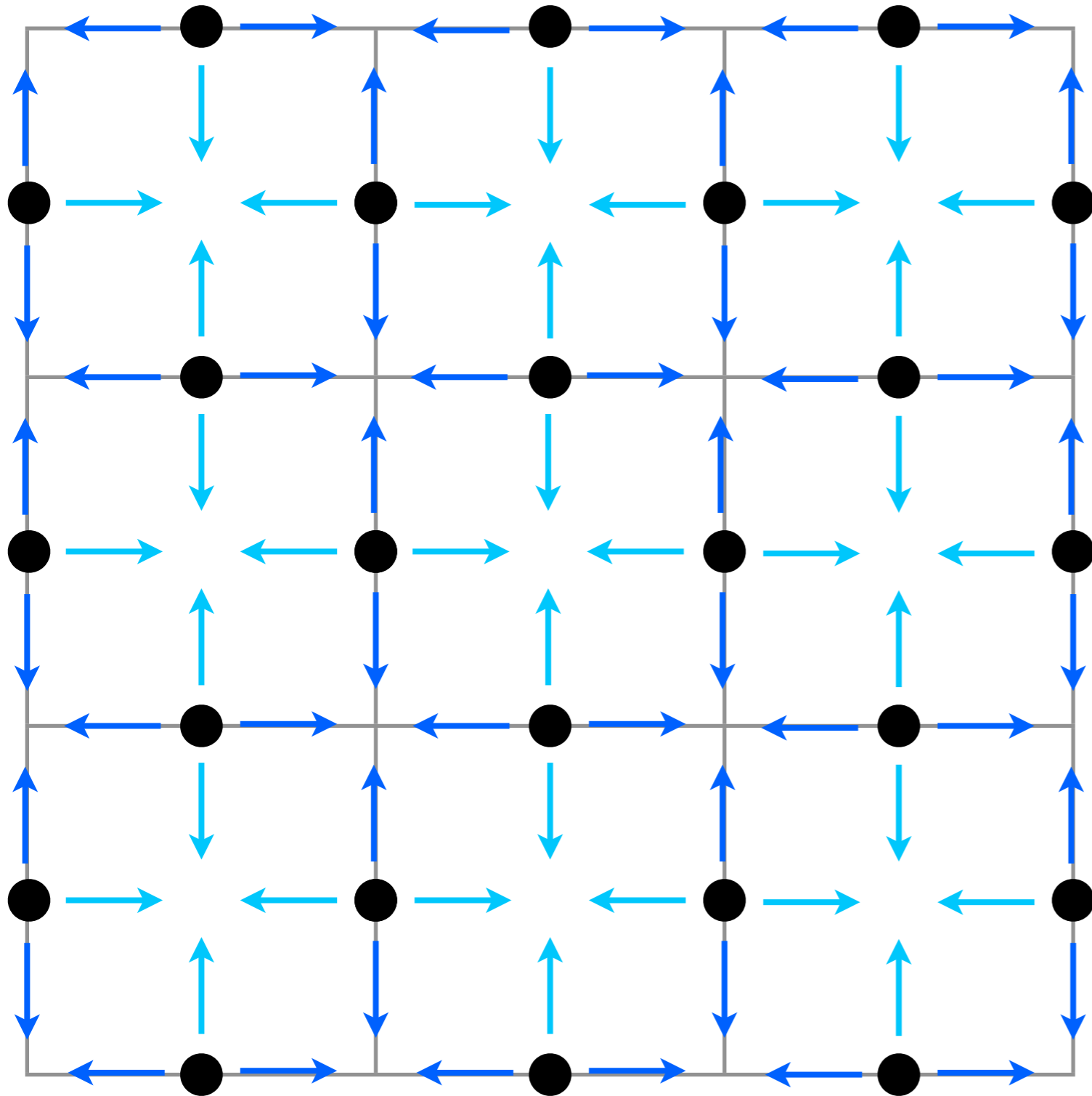
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Full lattice cycle

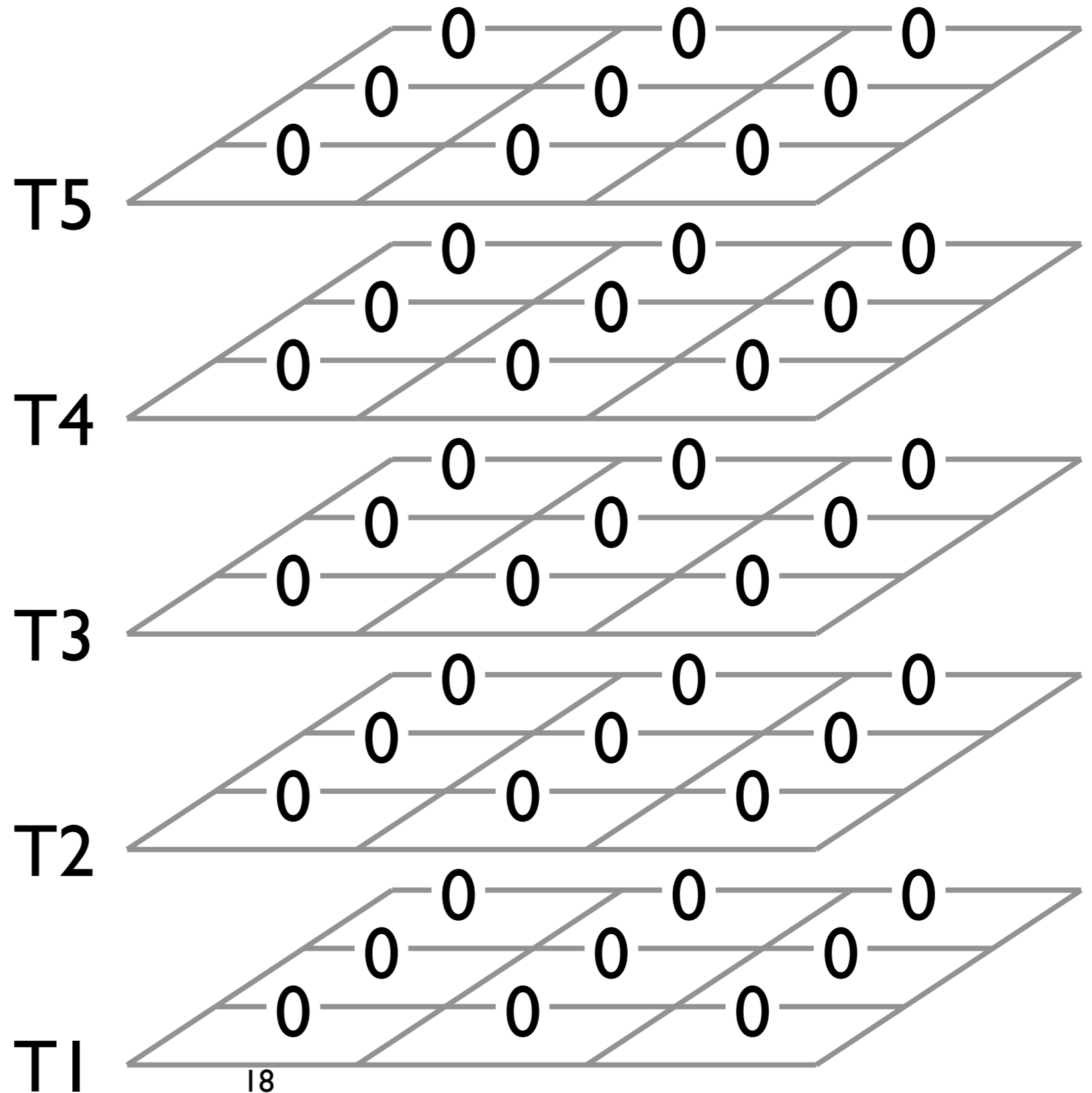


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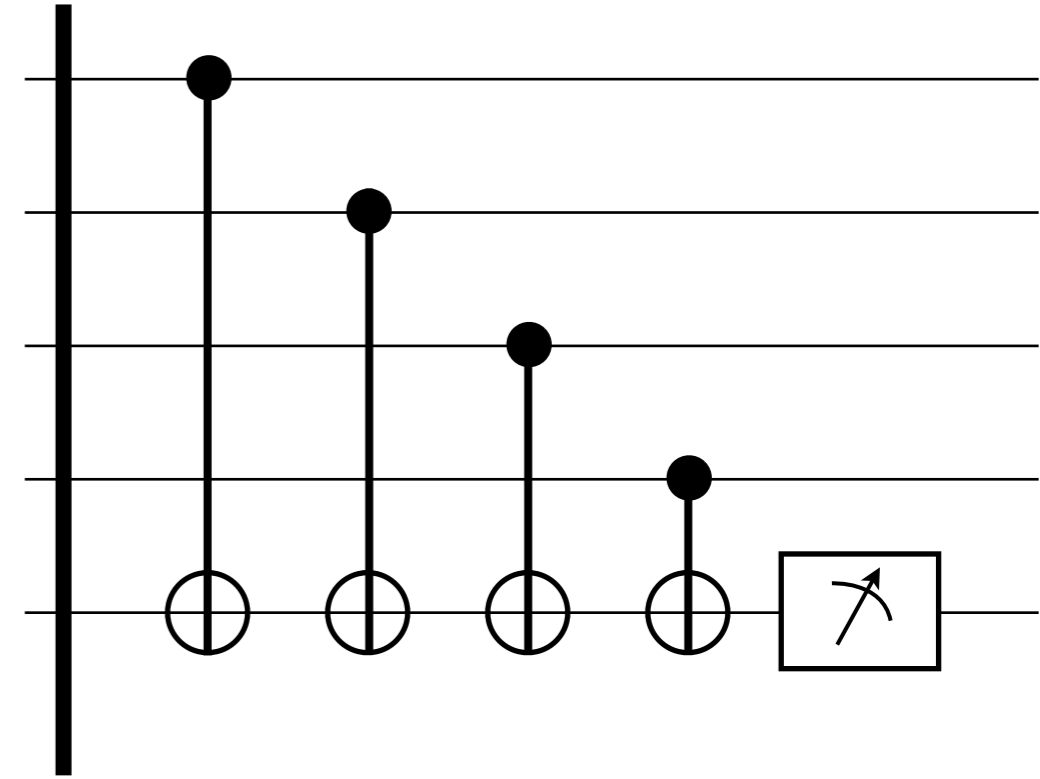
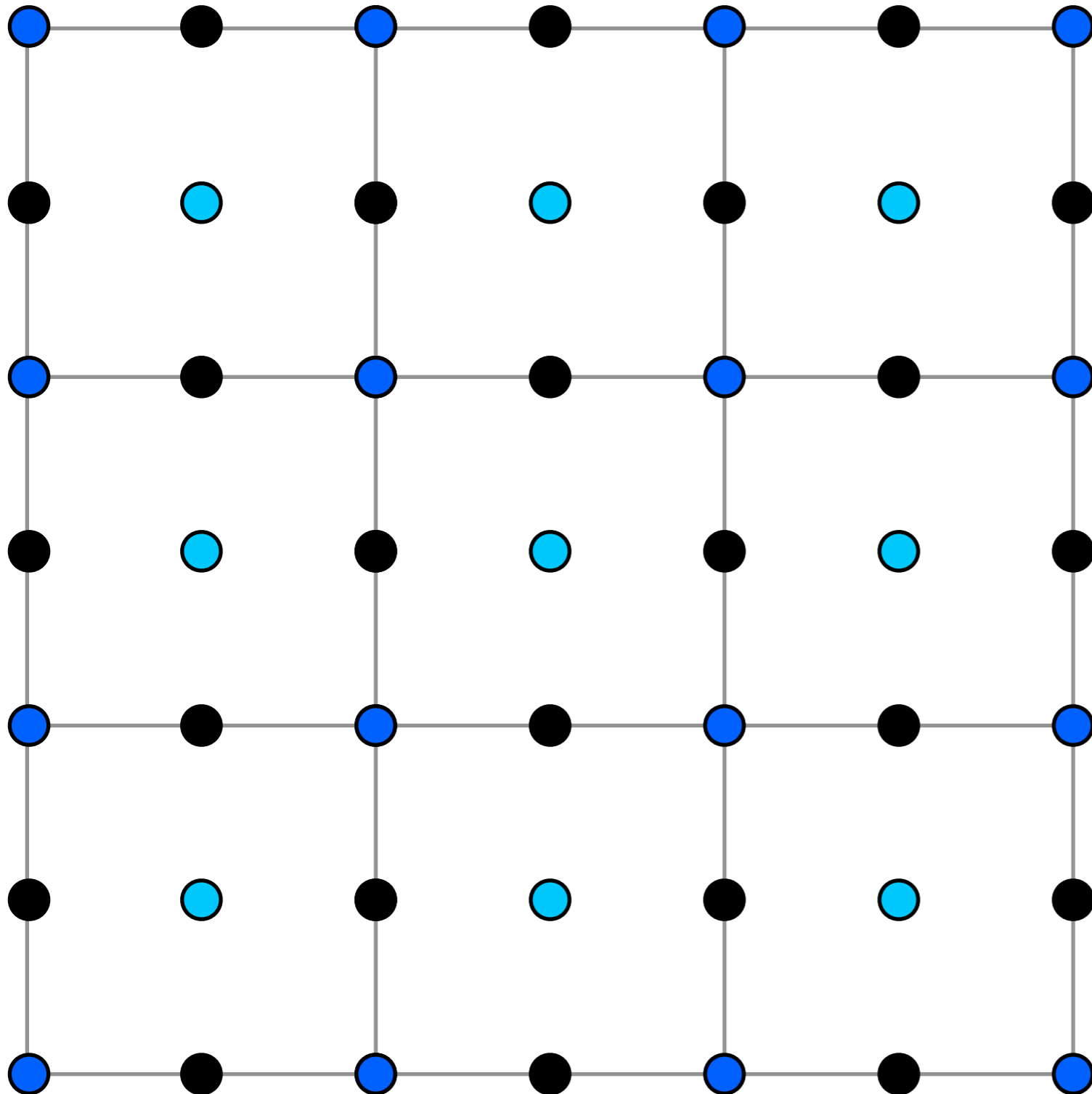
Full lattice cycle



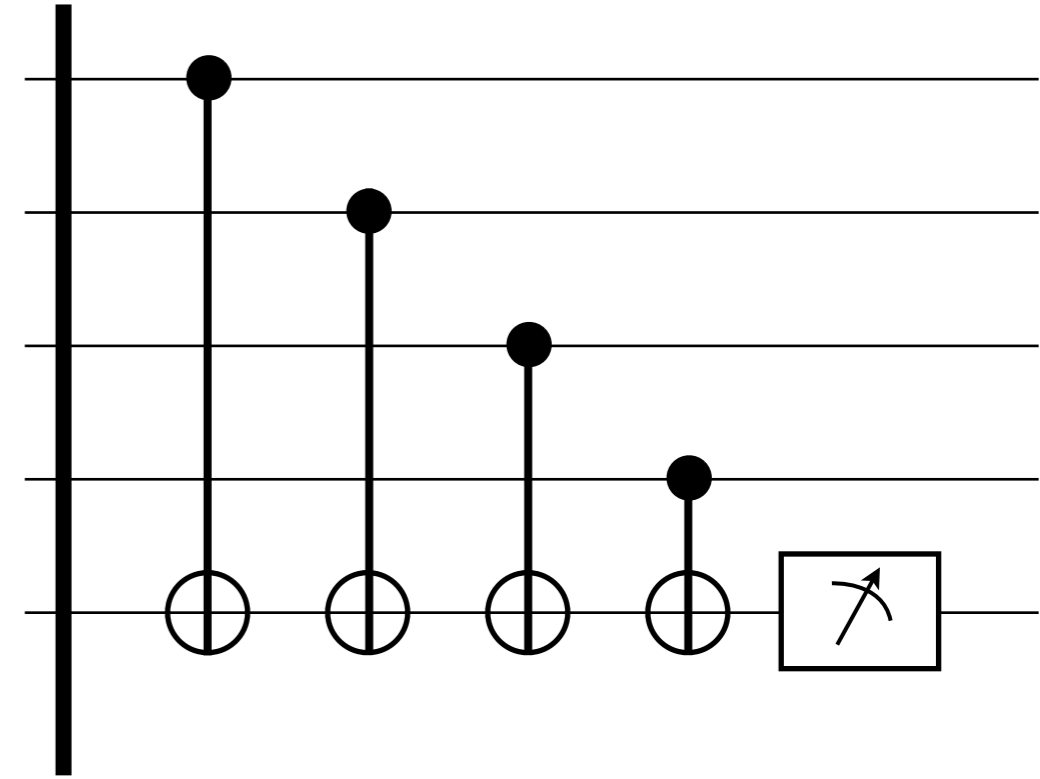
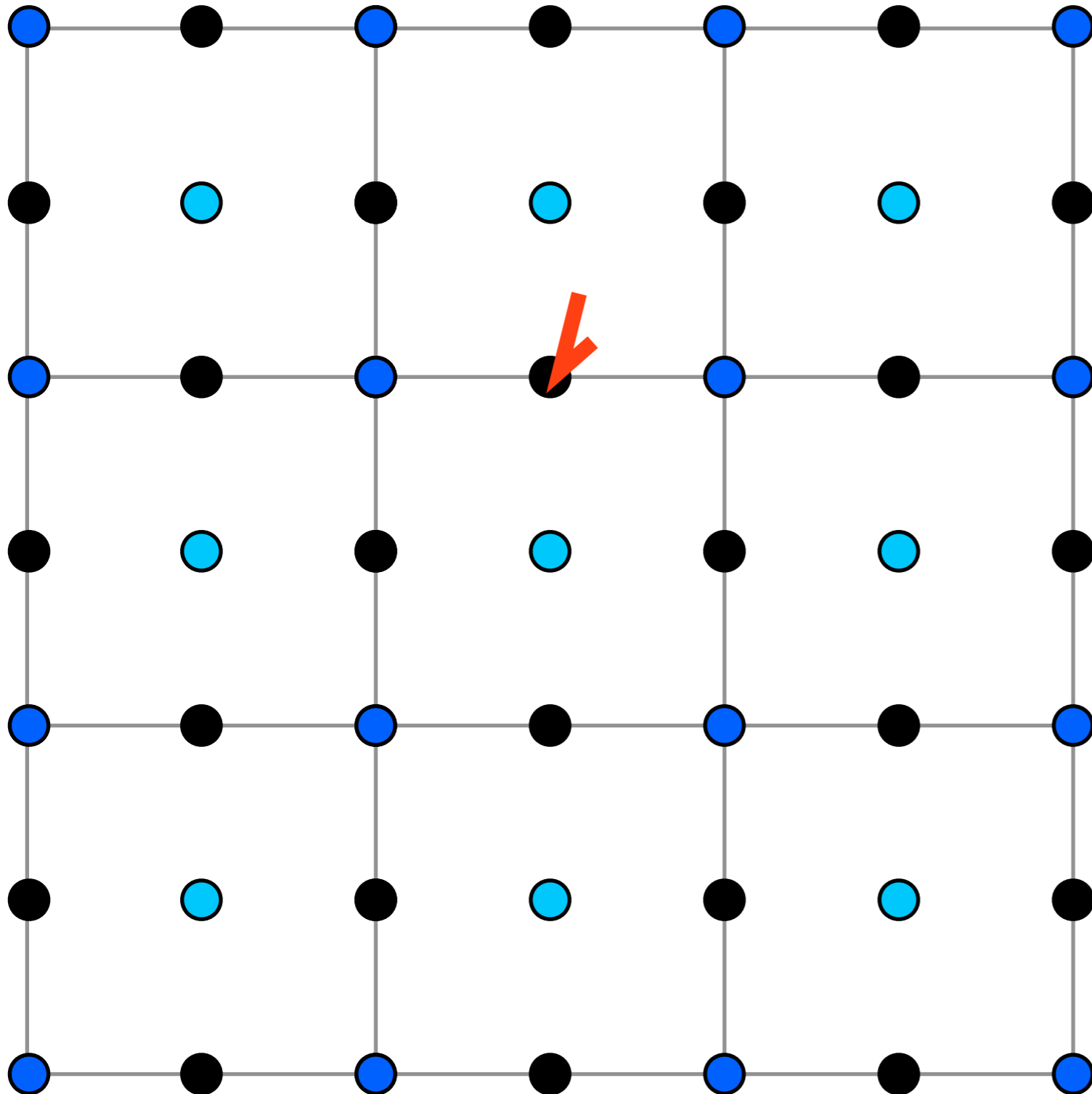
Error syndromes



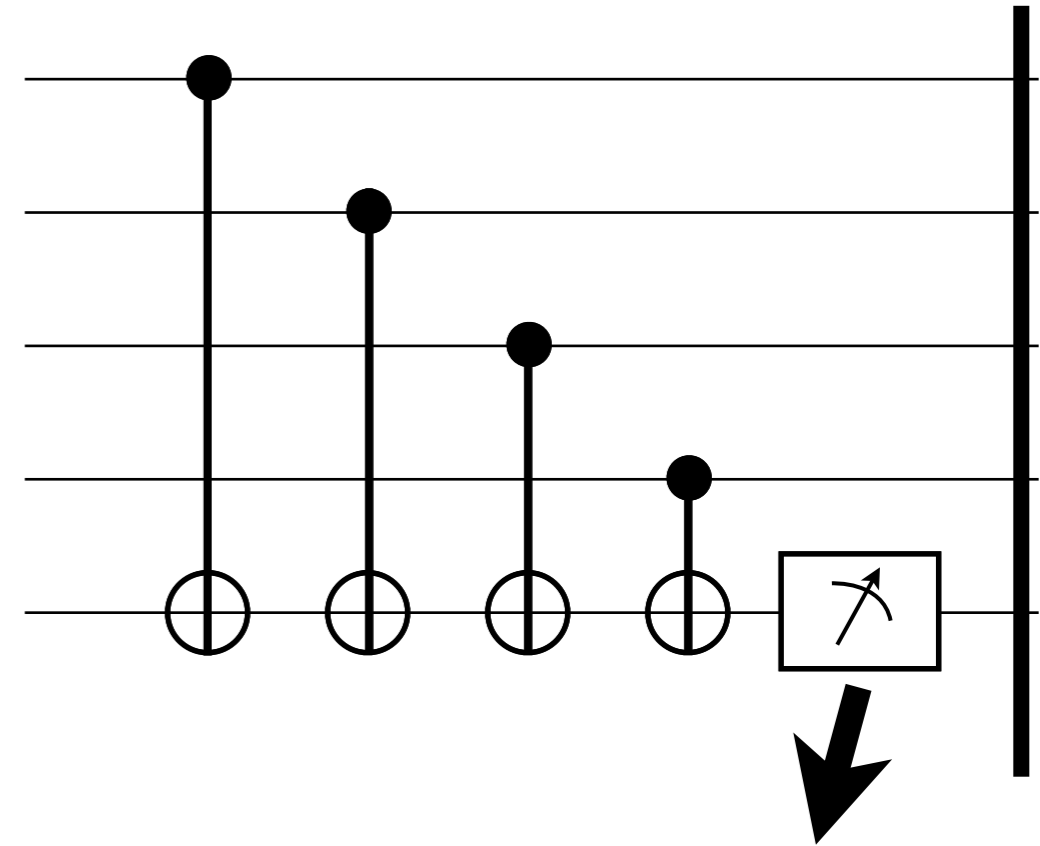
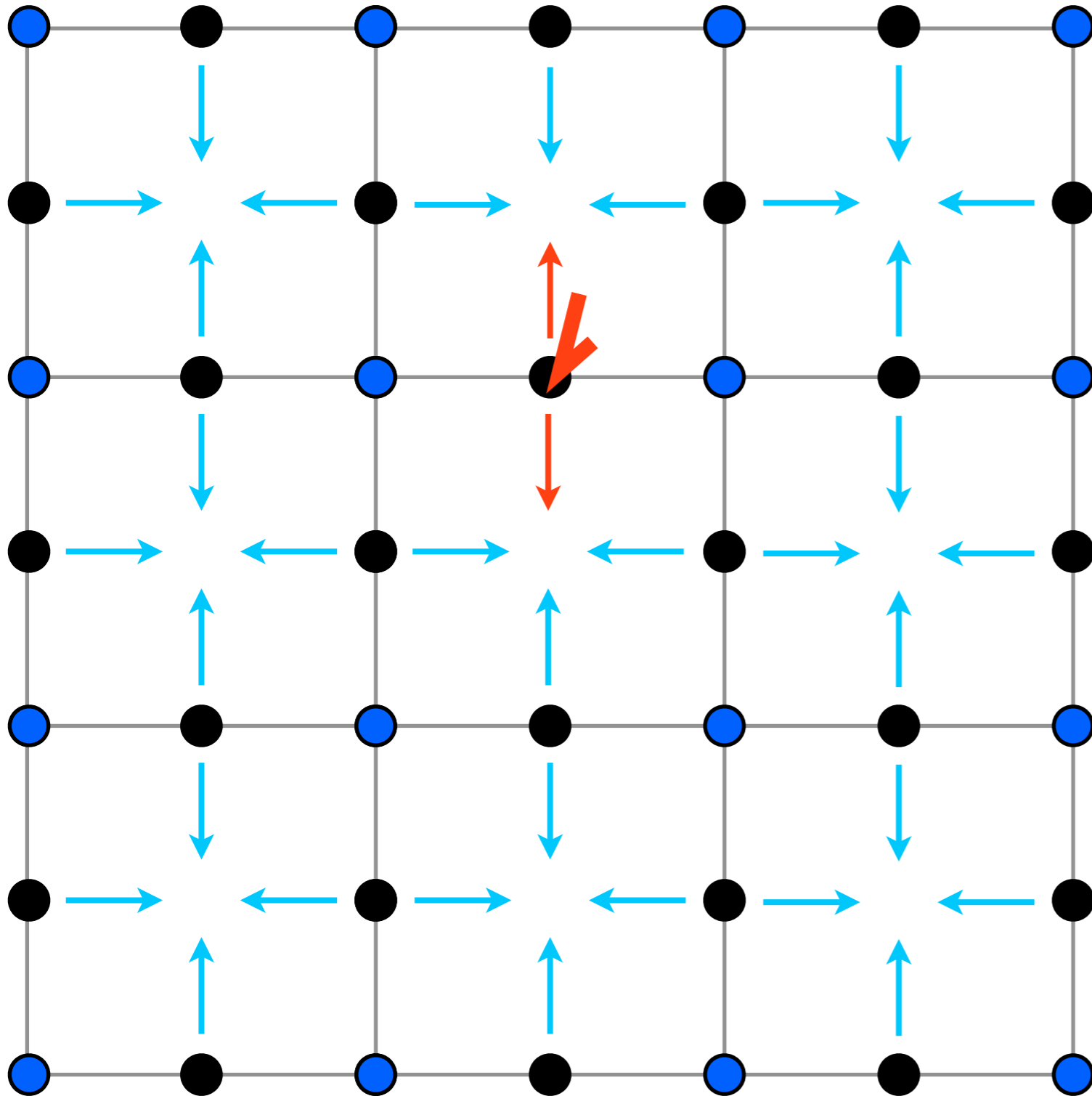
Memory Error I



Memory Error I

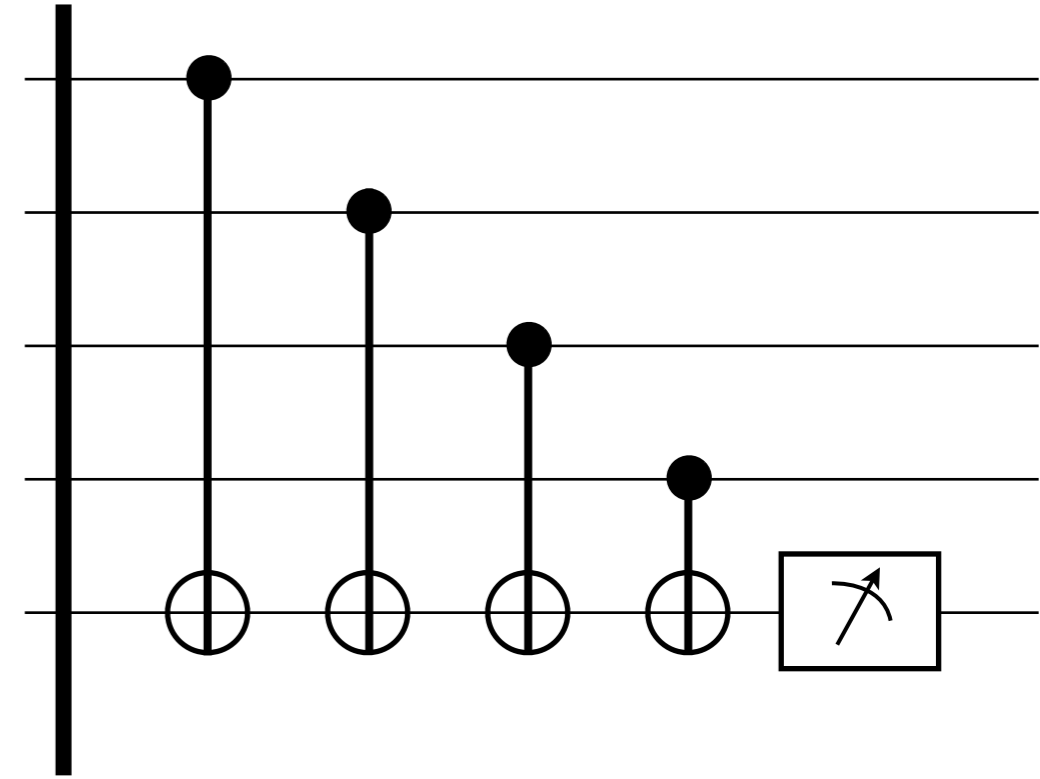
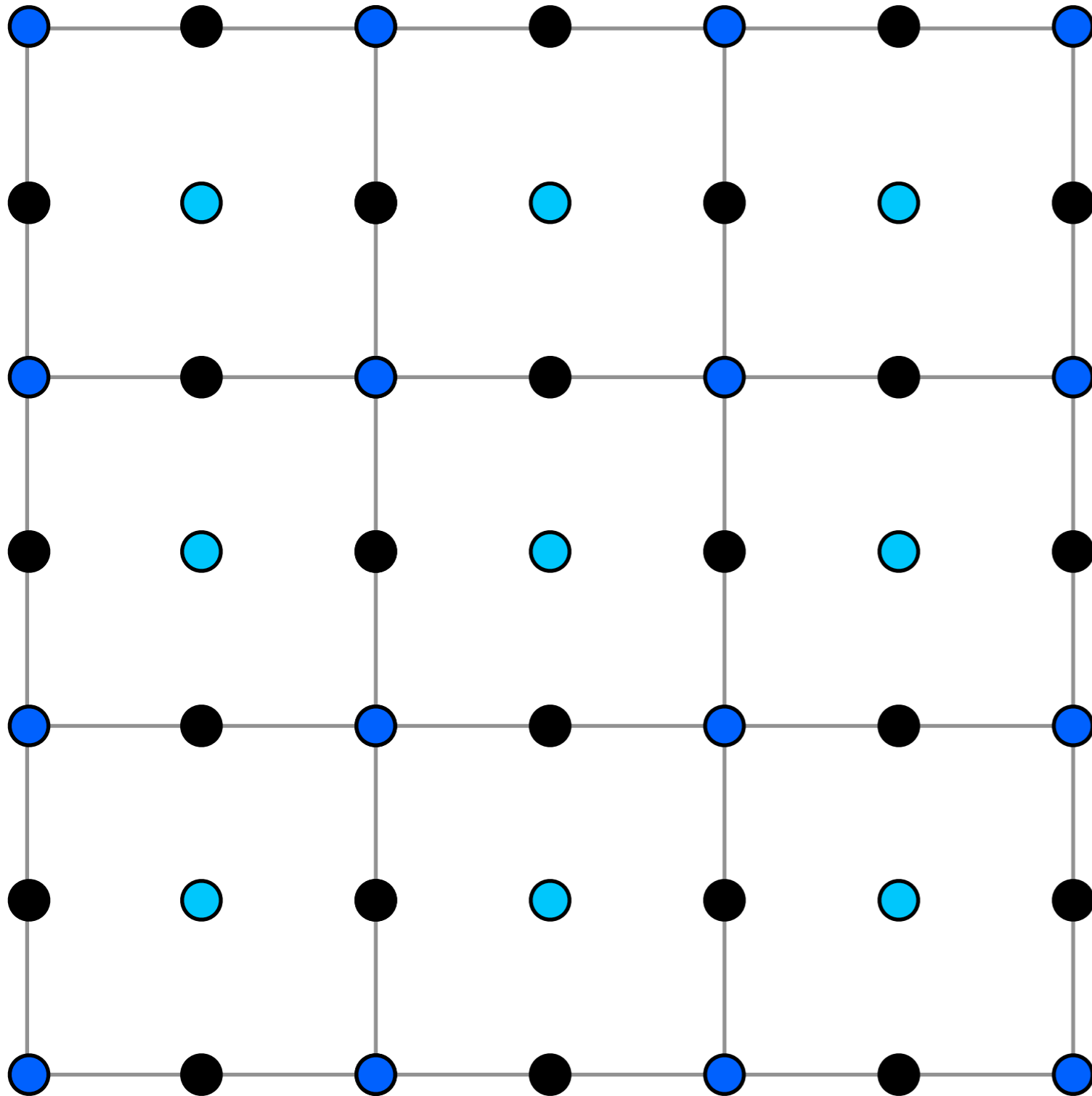


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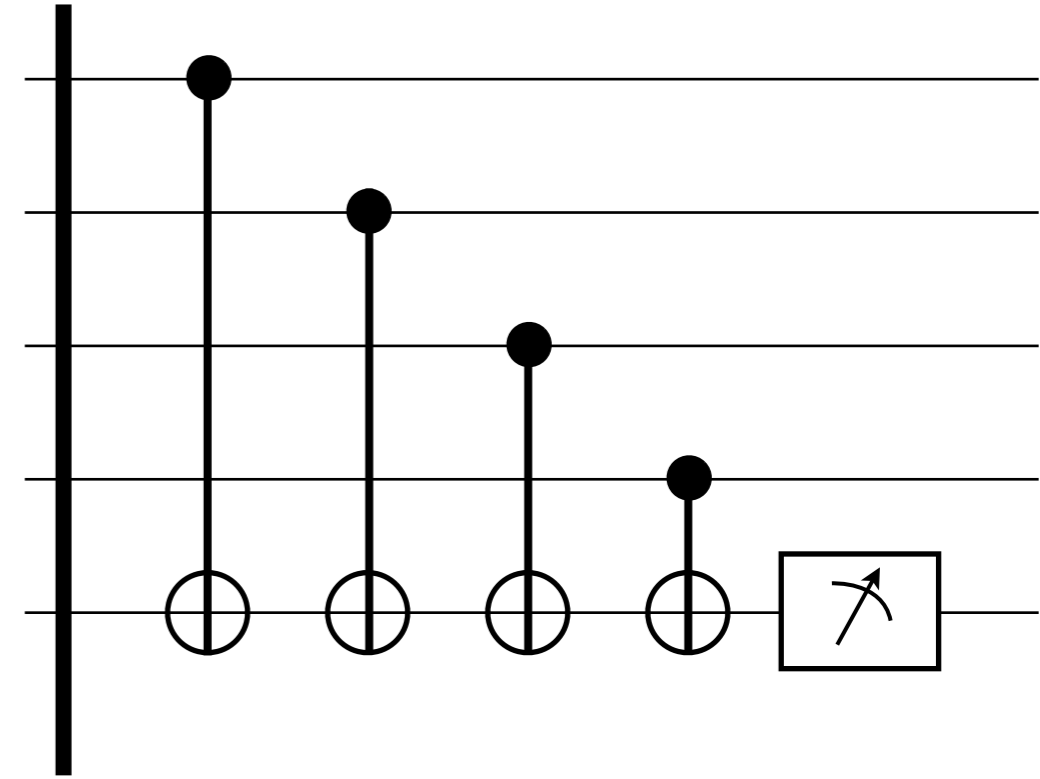
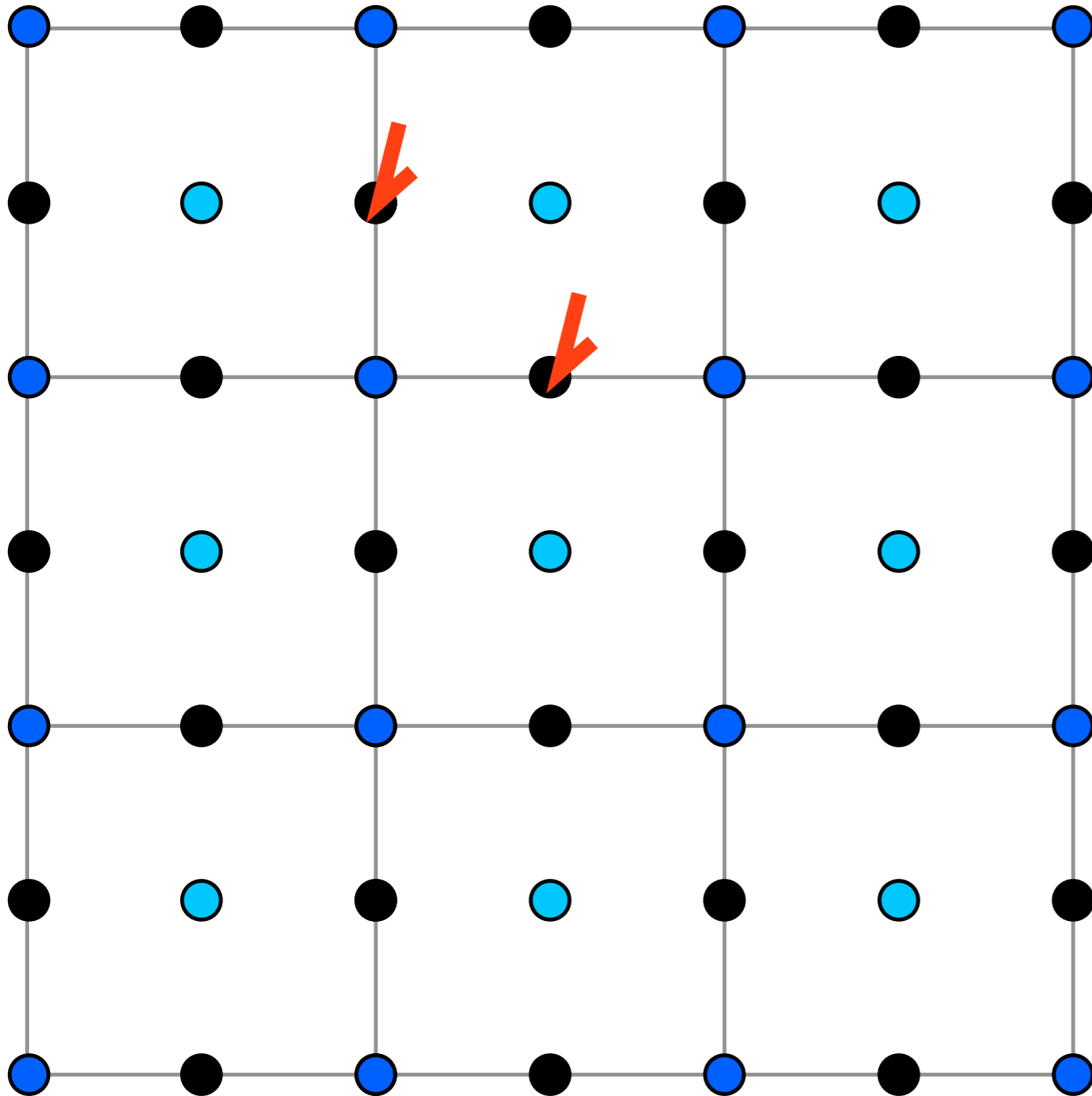


0		0
0		0
0	0	0

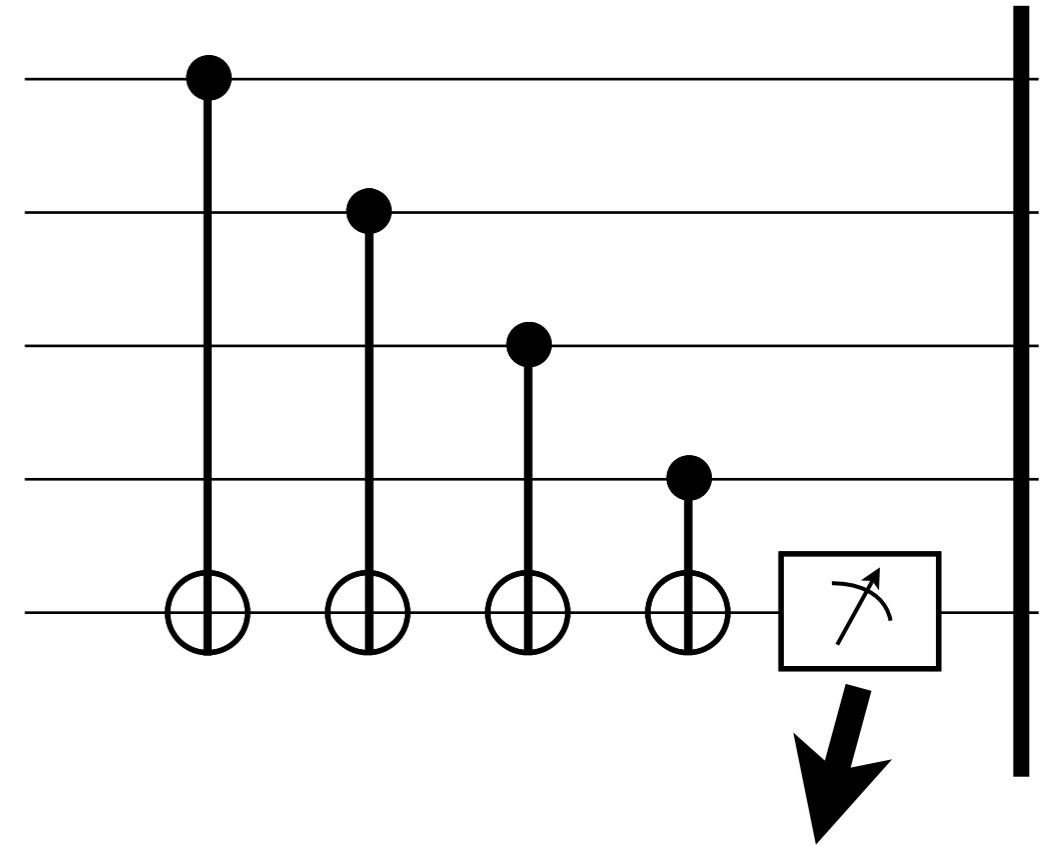
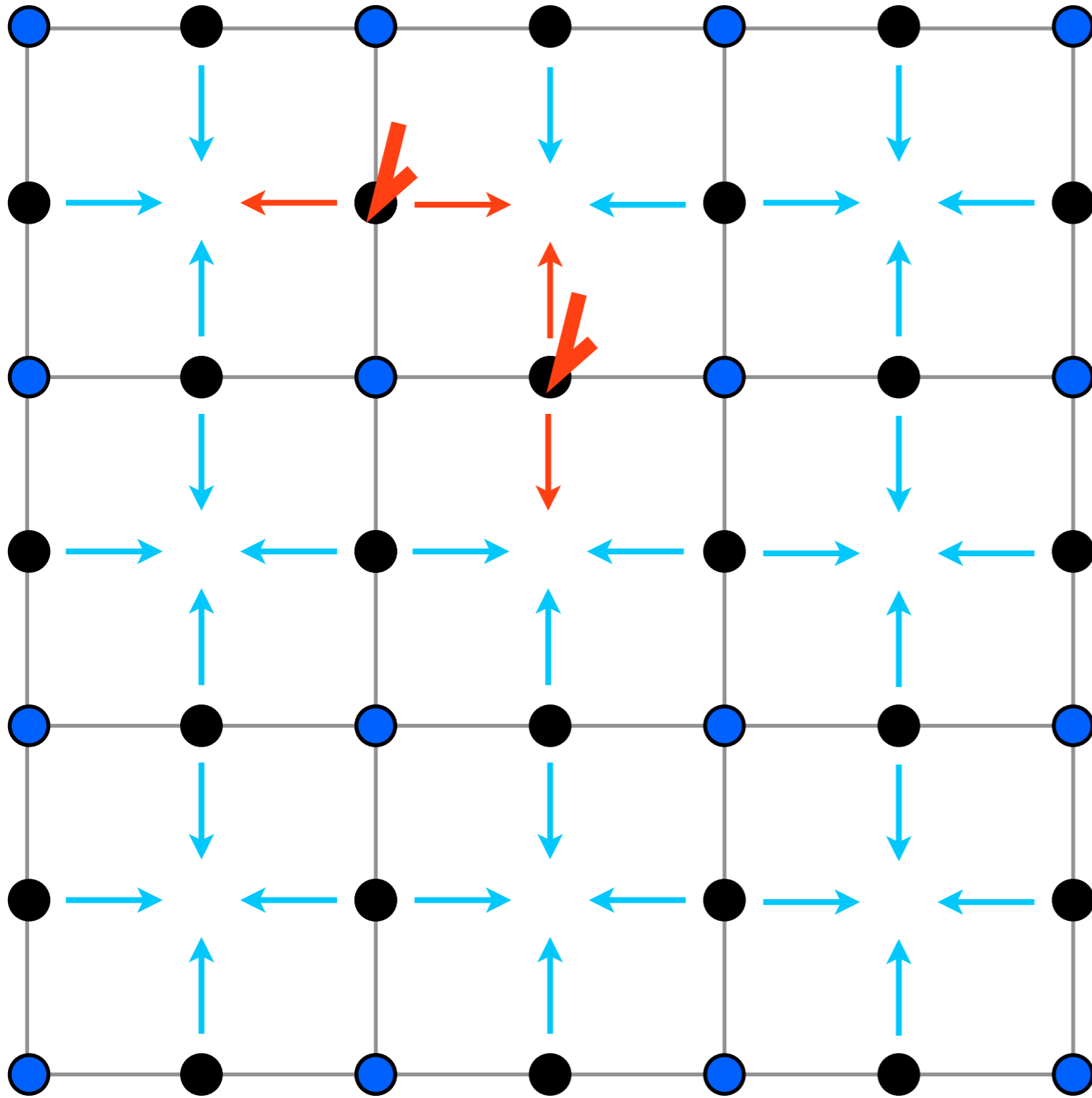
Memory Error 2



Memory Error 2

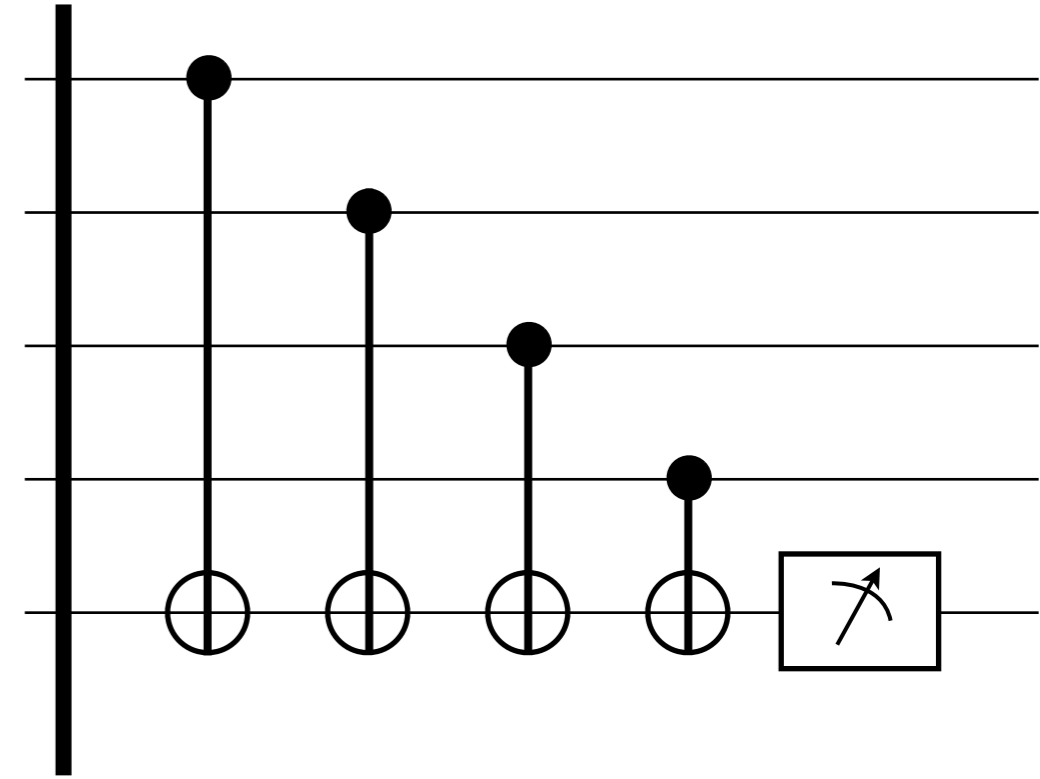
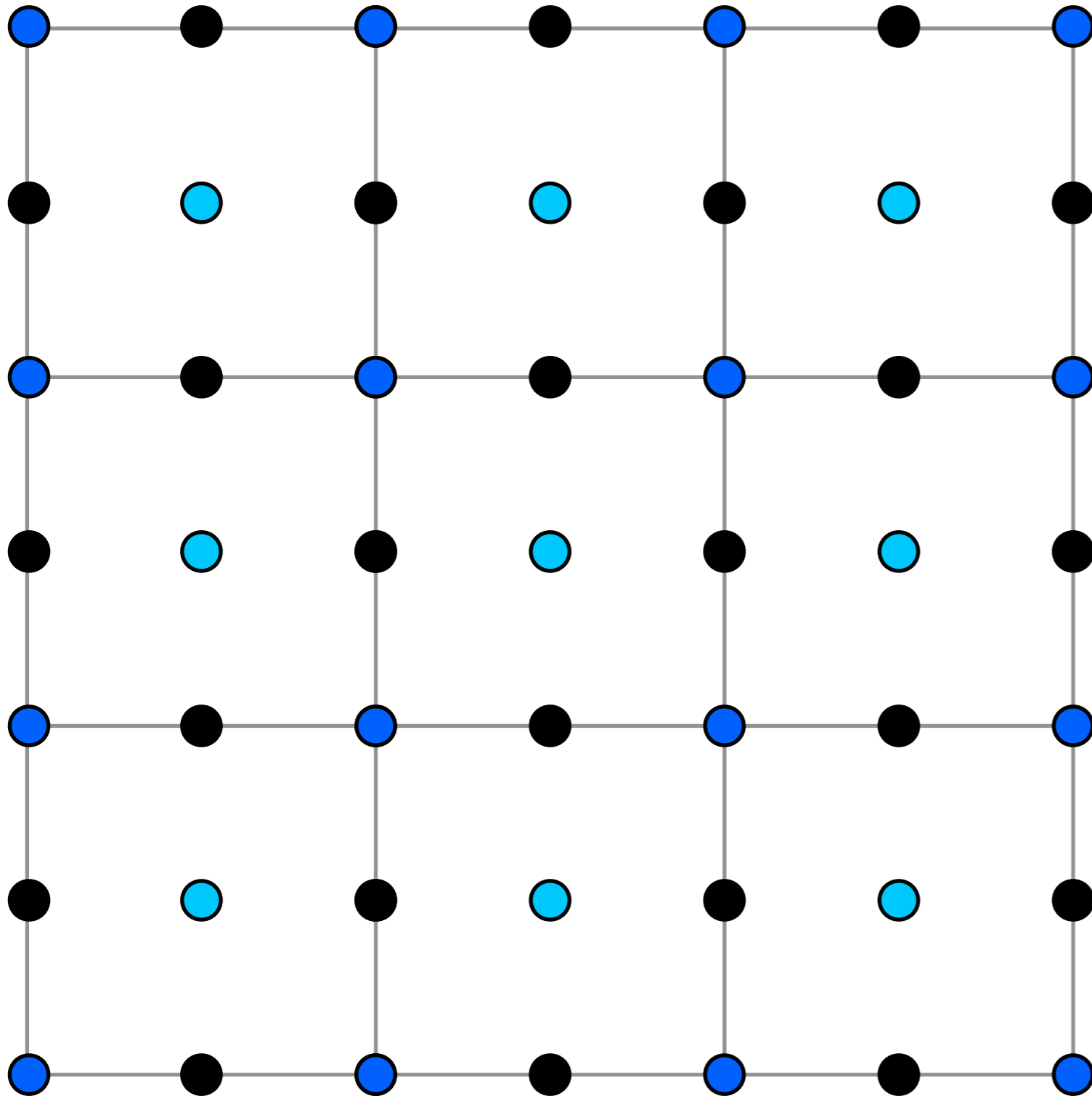


Memory Error 2

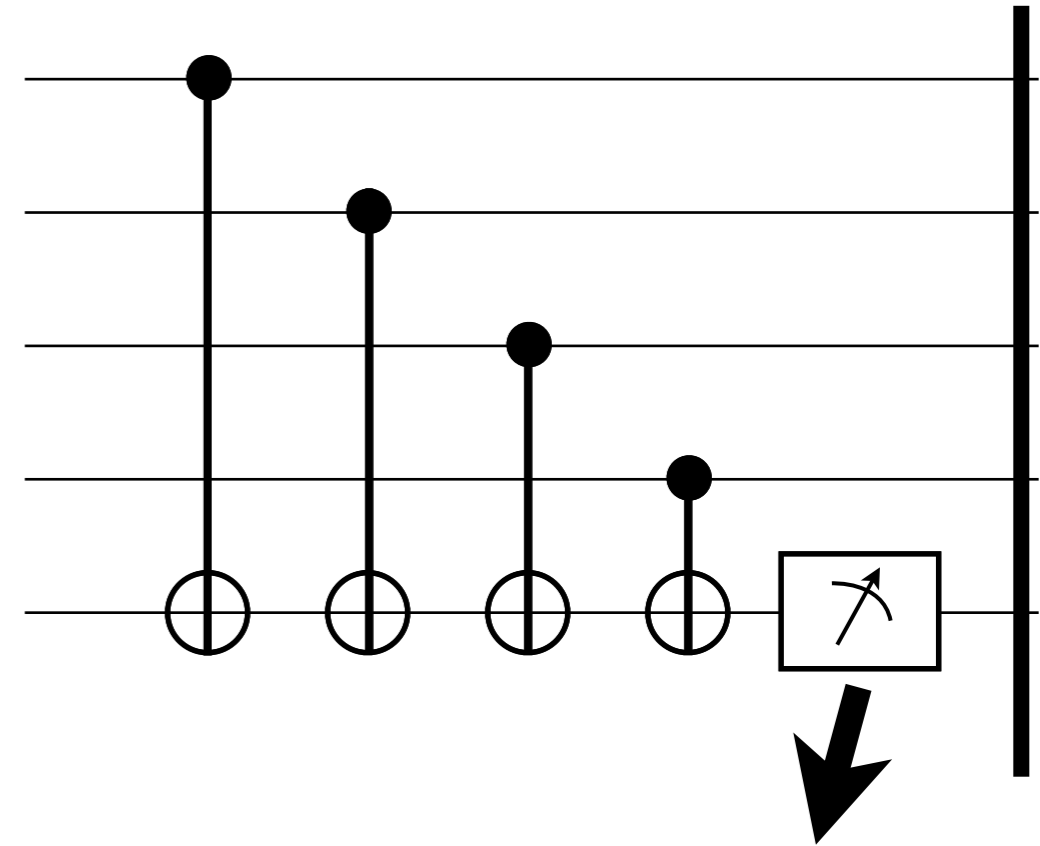
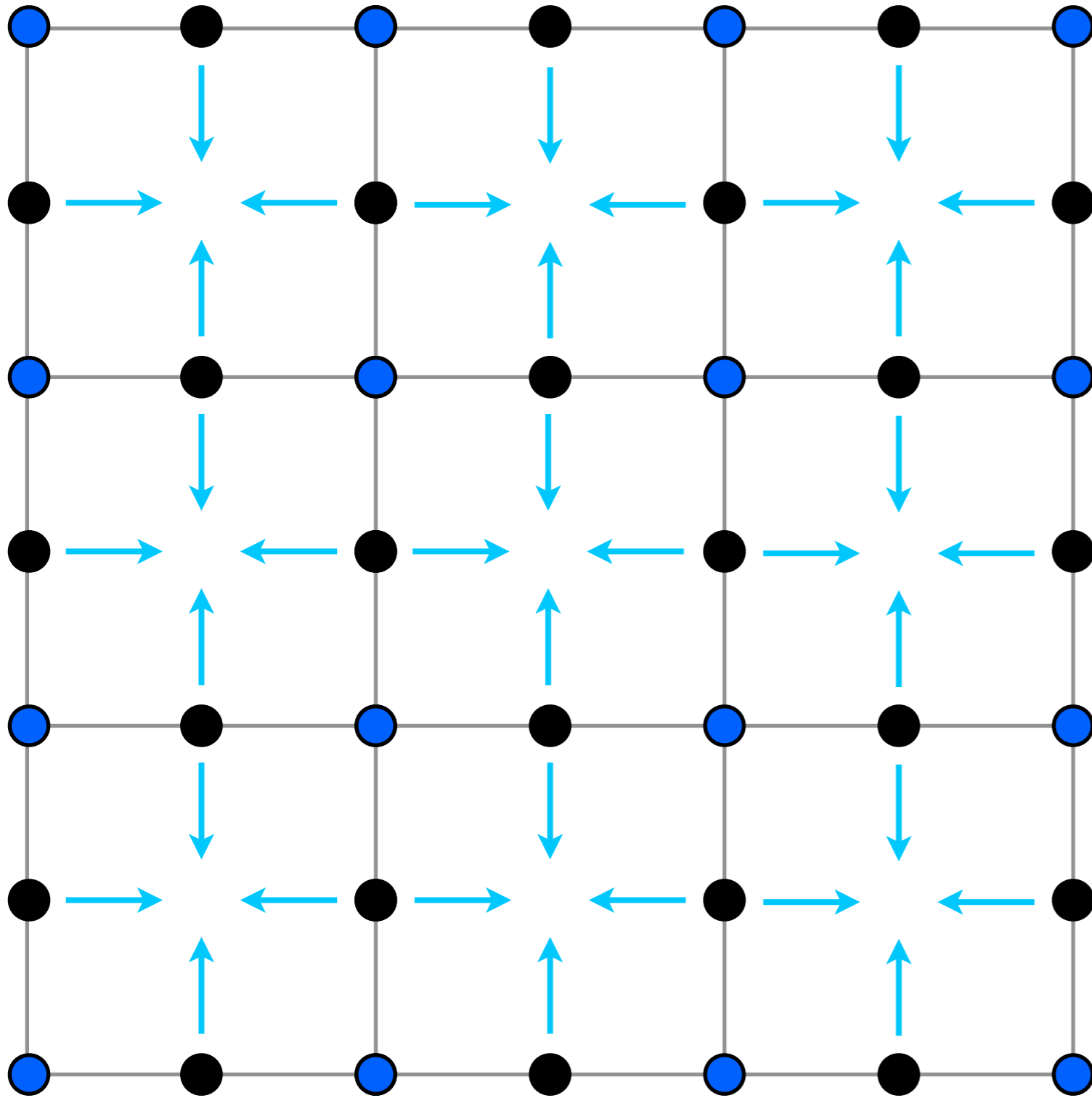


	0	0
0		0
0	0	0

Measurement Error I



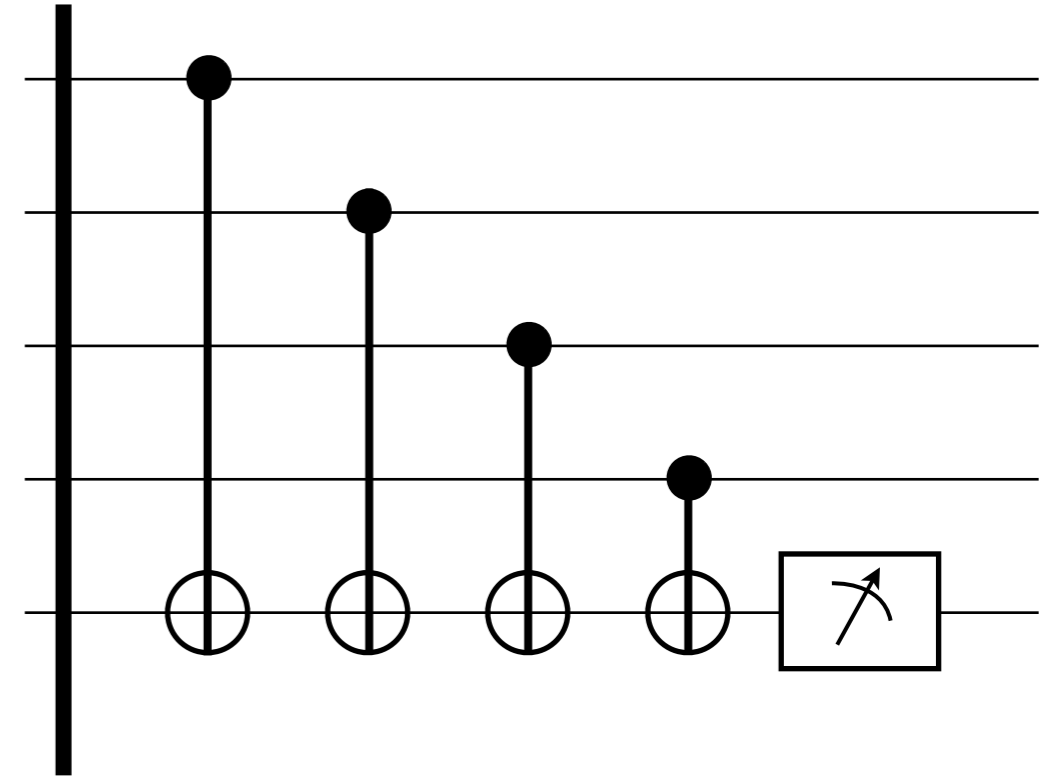
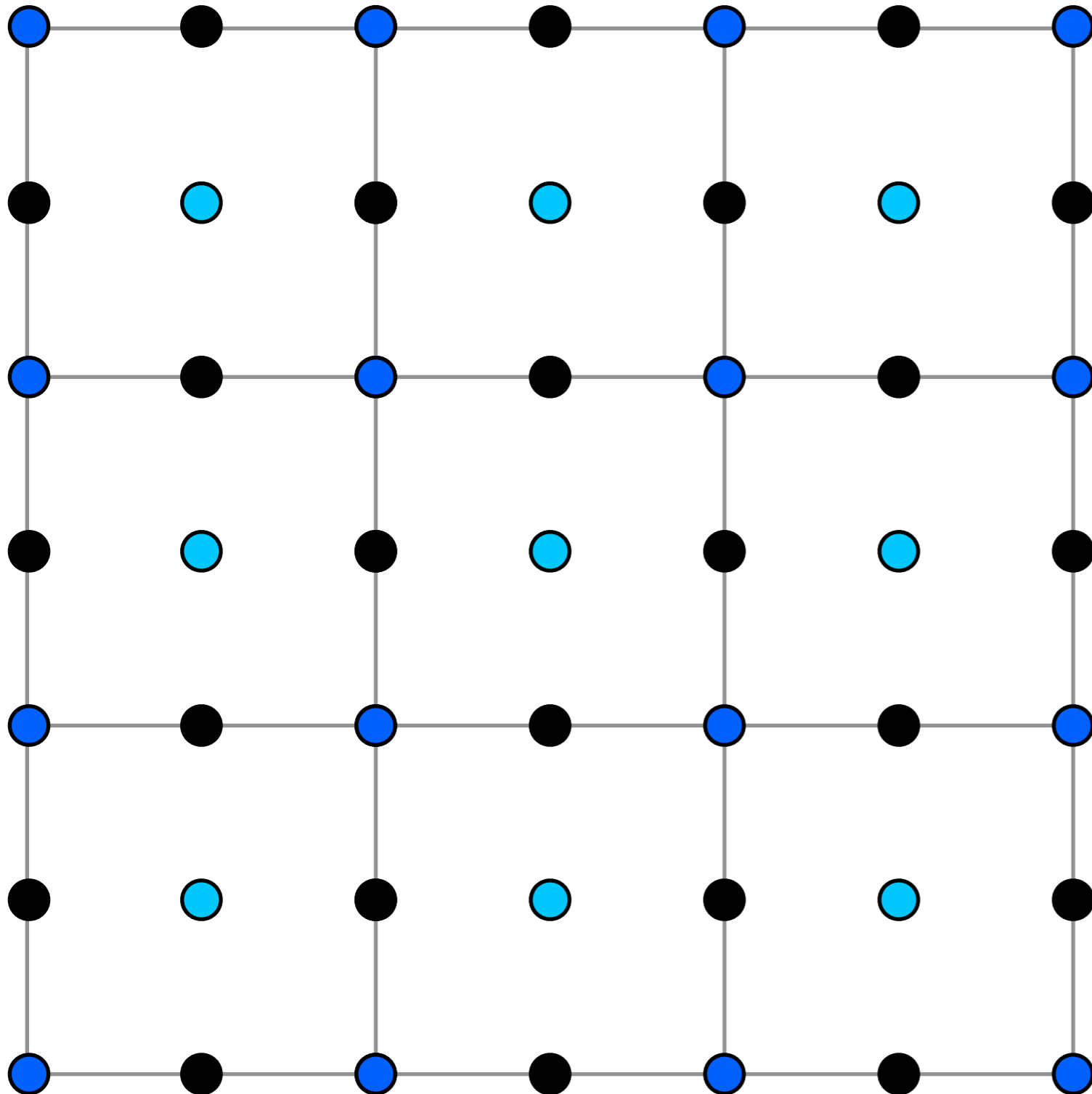
Measurement Error I



0	1	0
0	0	0
0	0	0

Measurement Error 2

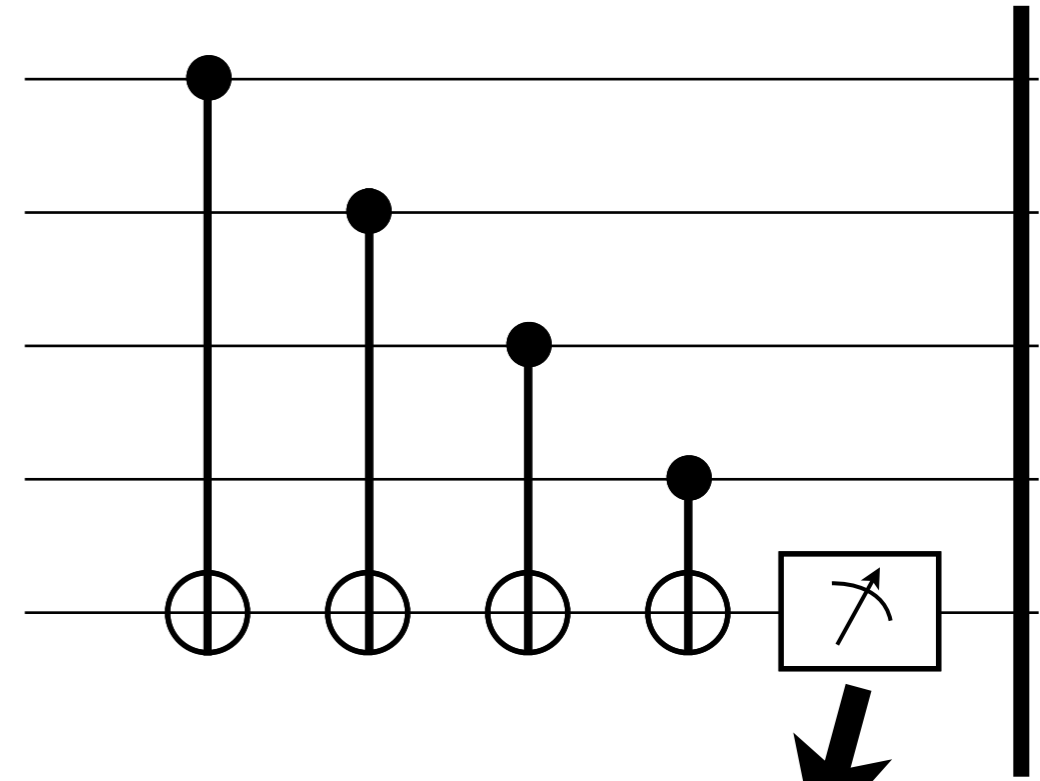
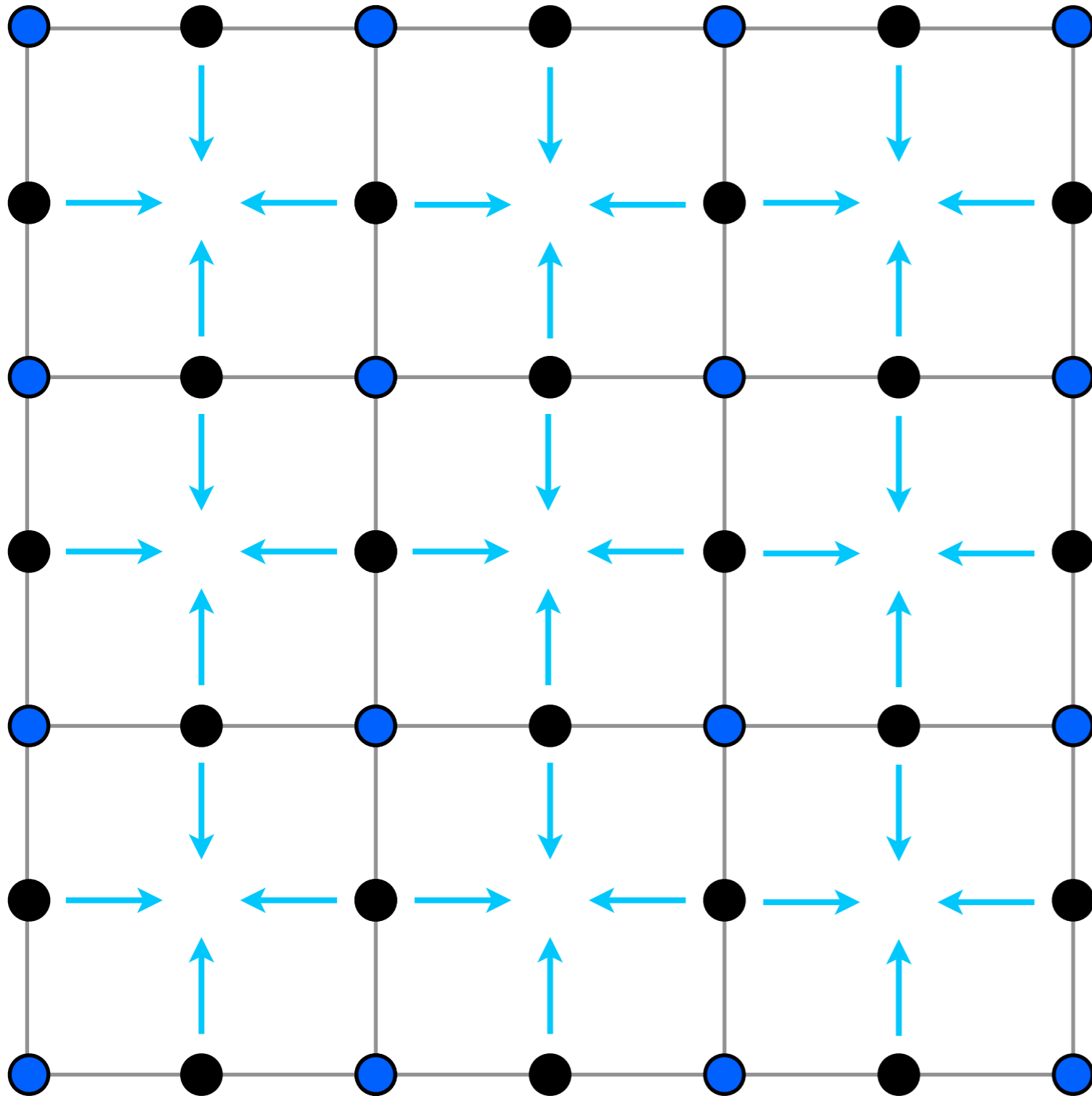
Next Step



0		0
0	0	0
0	0	0

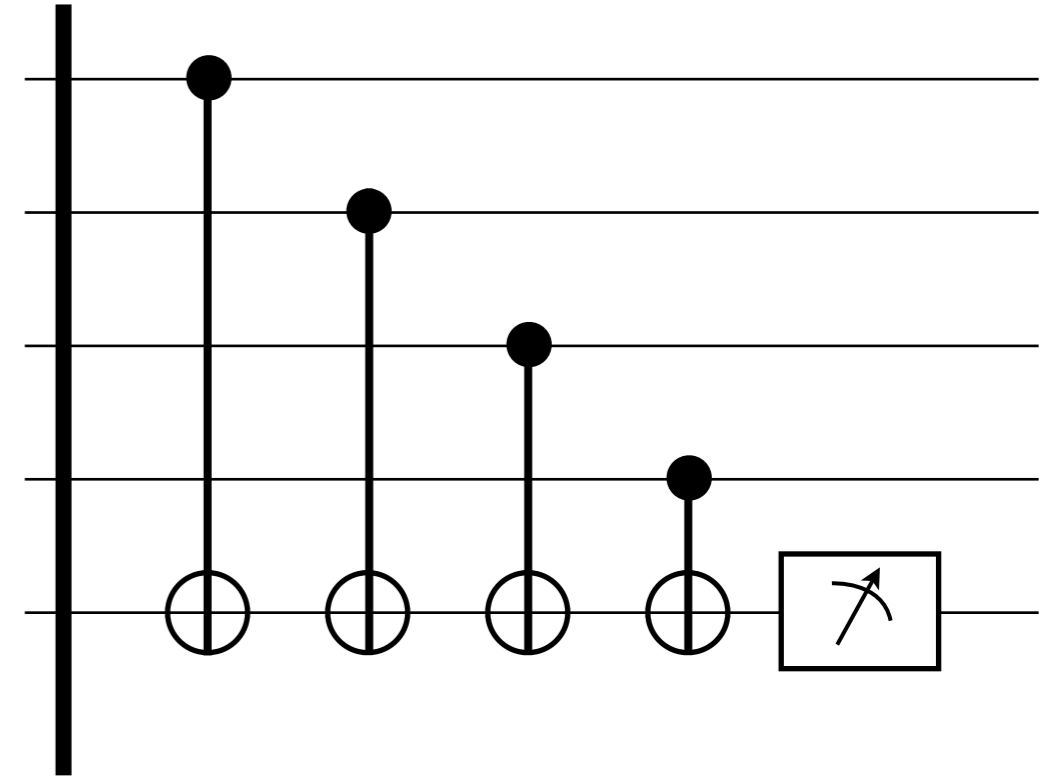
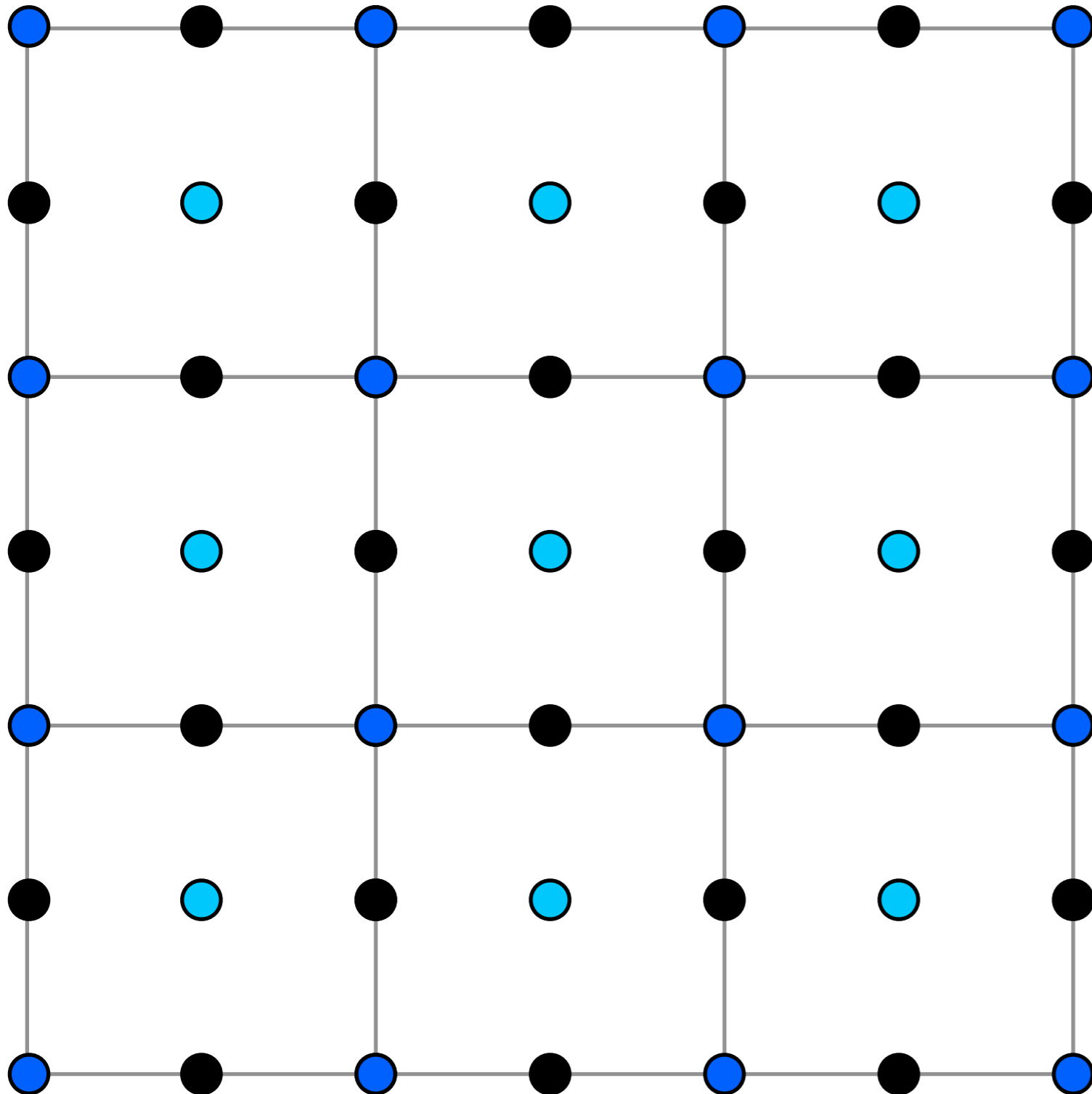
Measurement Error 2

Next Step

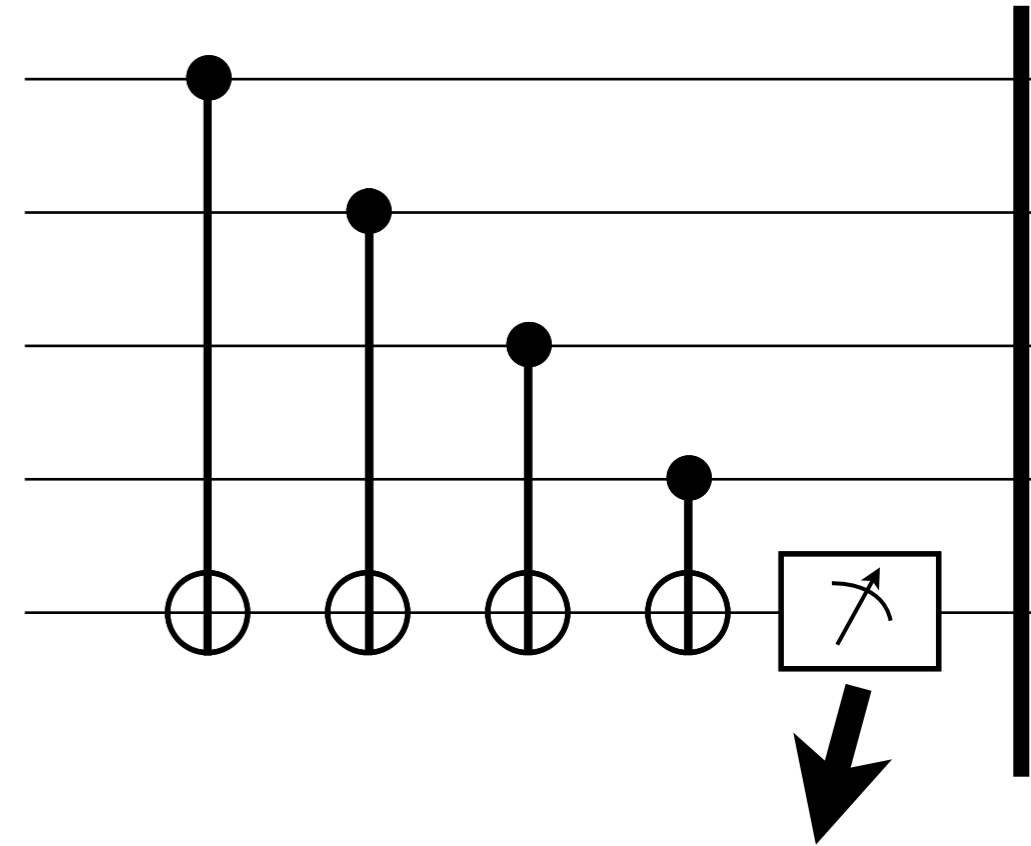
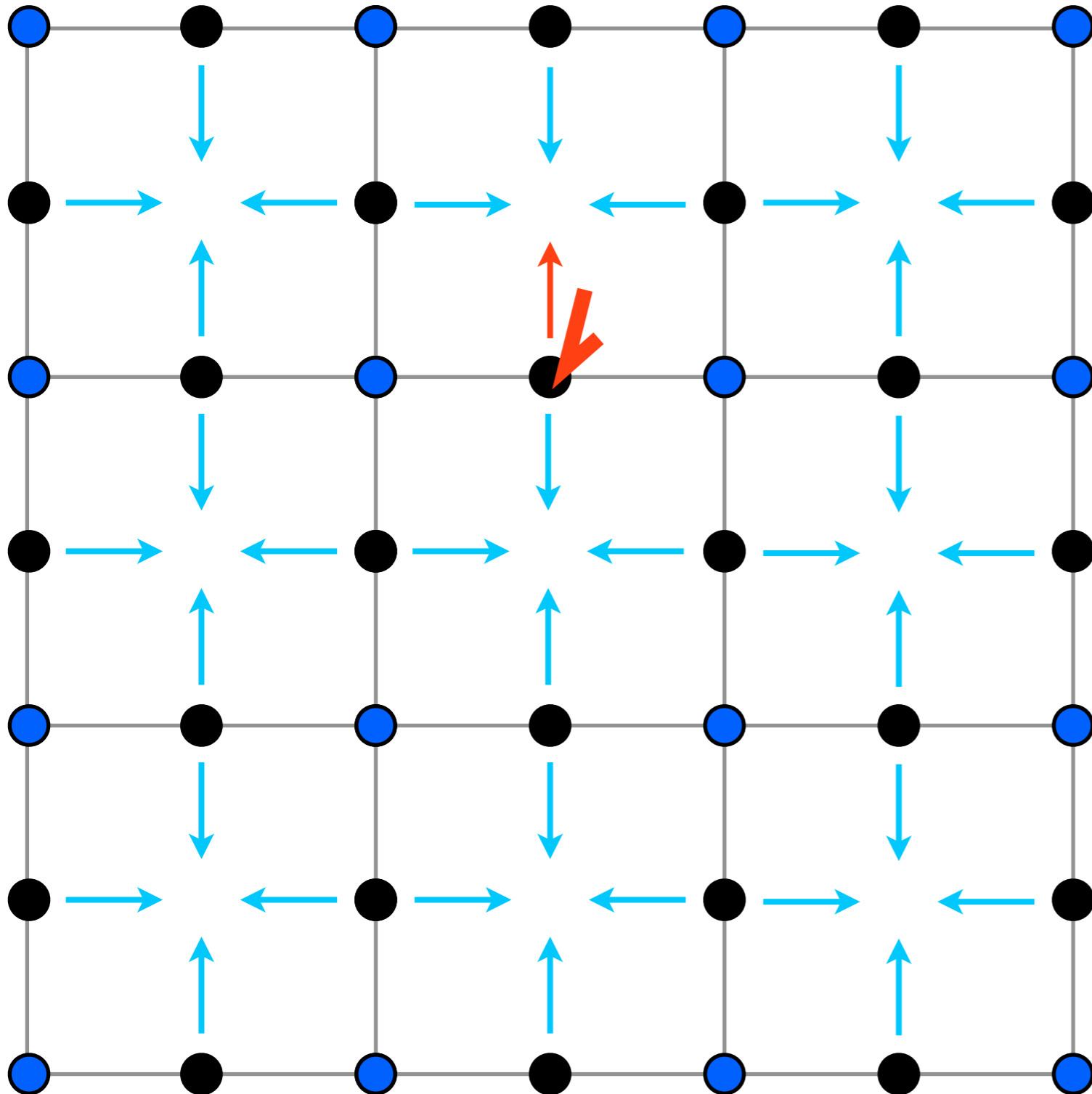


0^0	0^1	0^0
0^0	0^0	0^0
0^0	0^0	0^0

Gate Error I

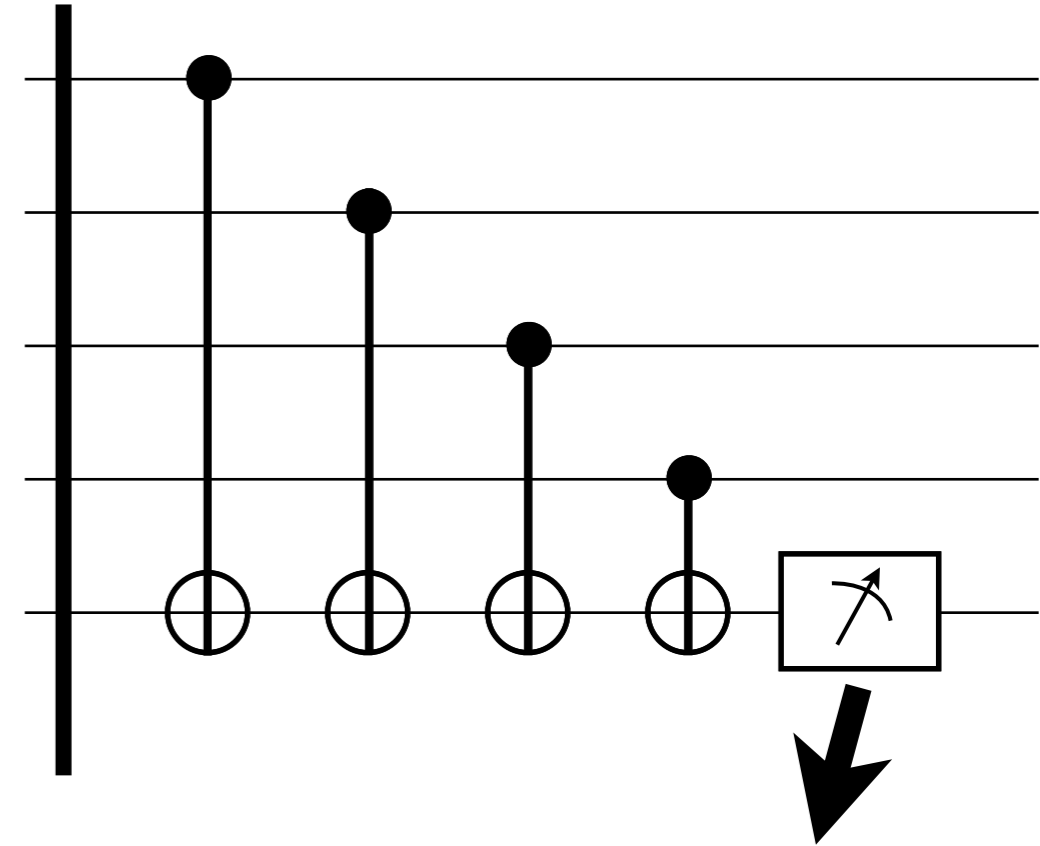
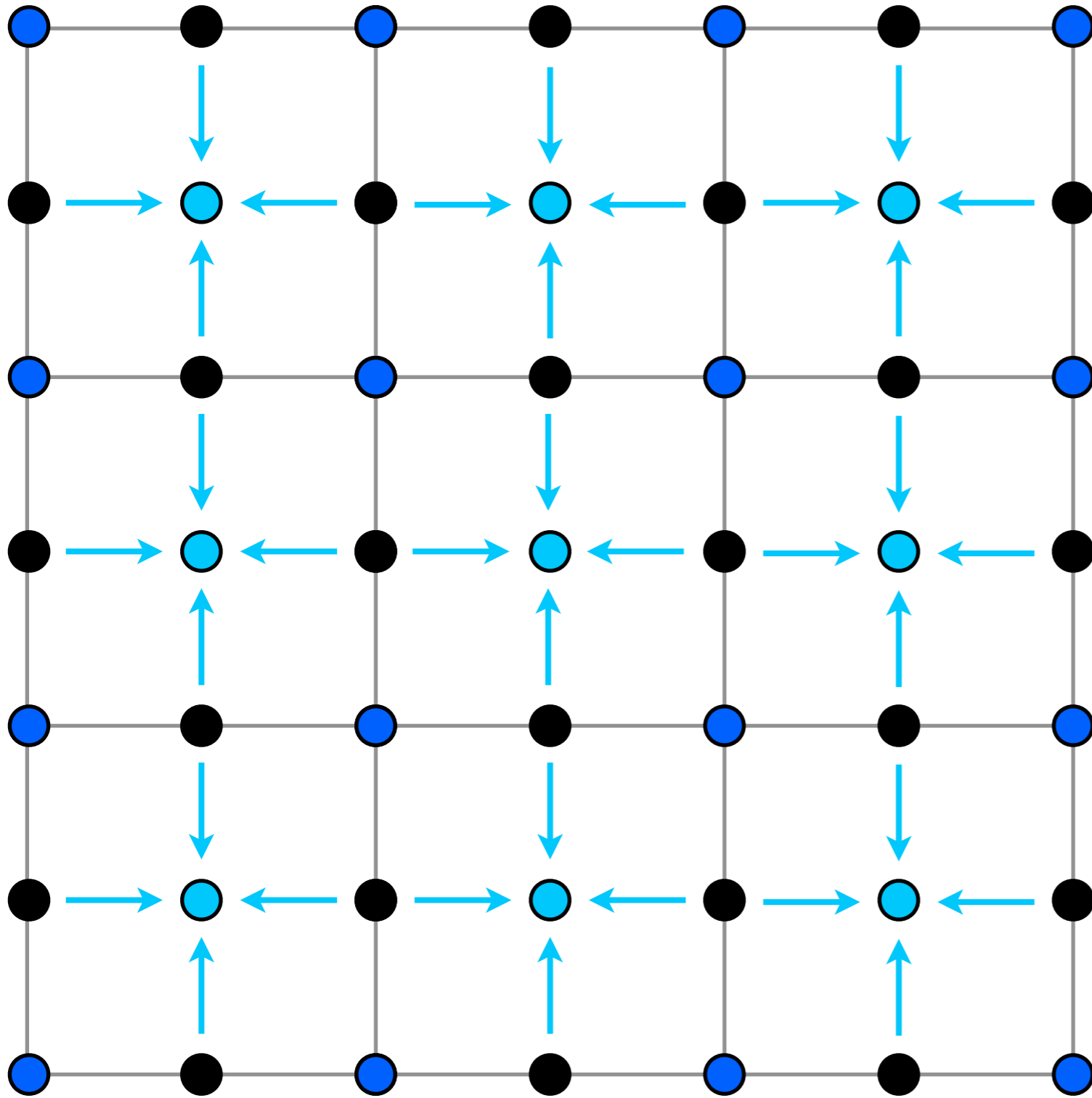


Gate Error I



0	I	0
0	0	0
0	0	0

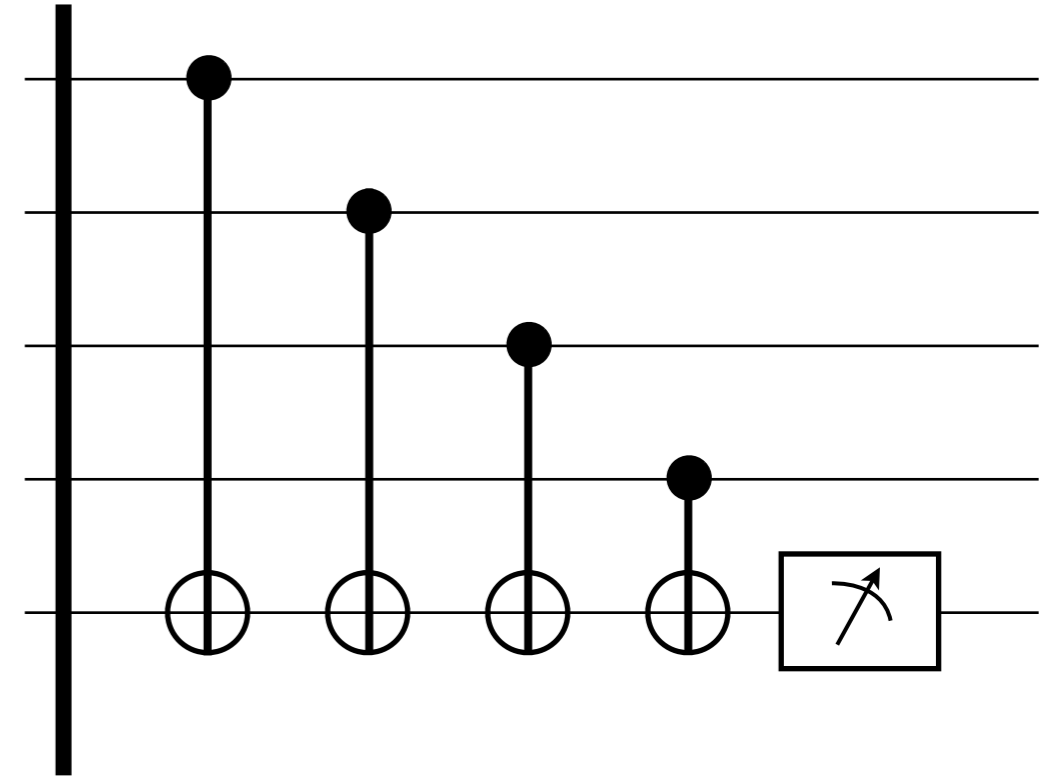
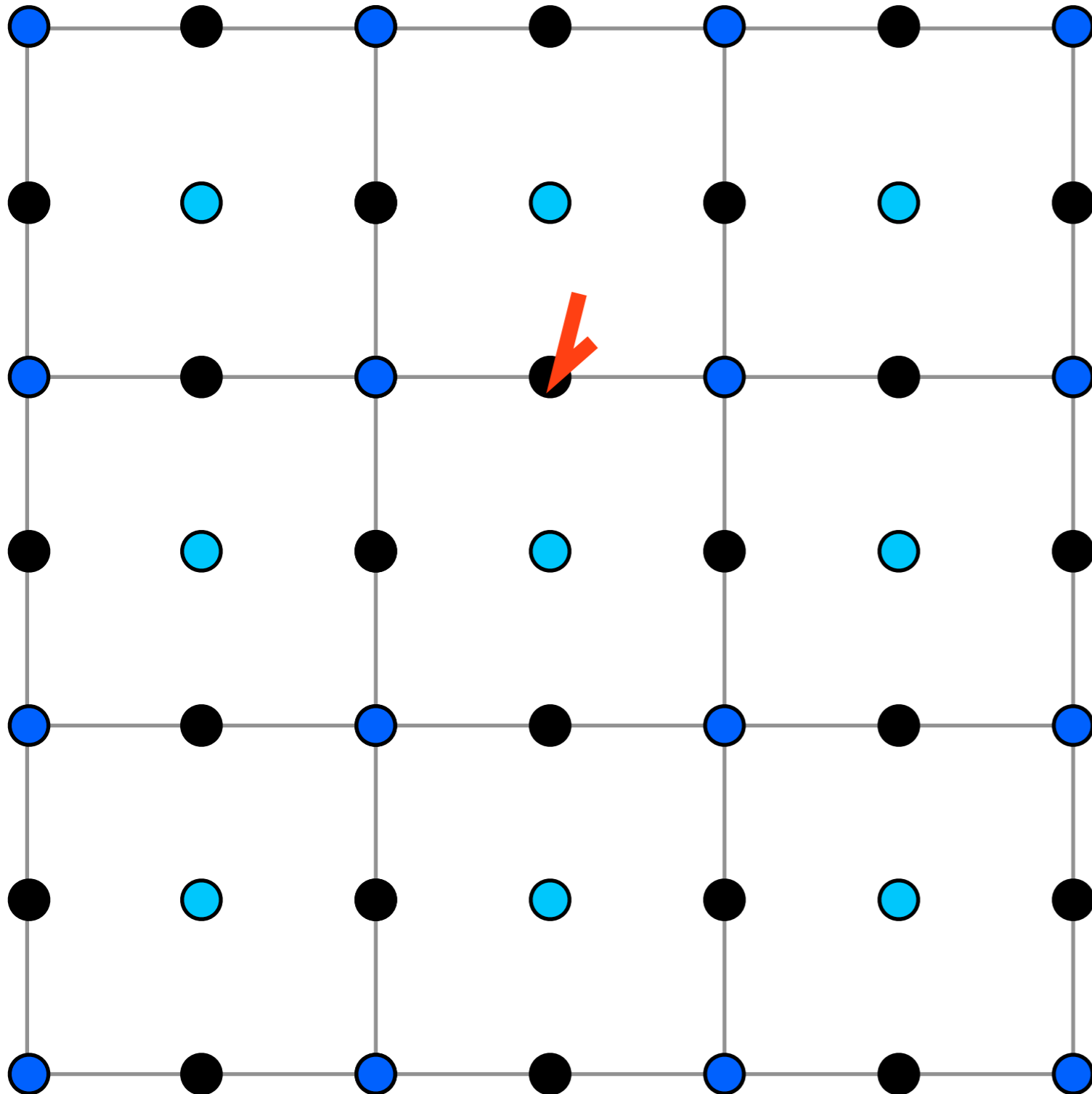
Measurement Error



0	1	0
0	0	0
0	0	0

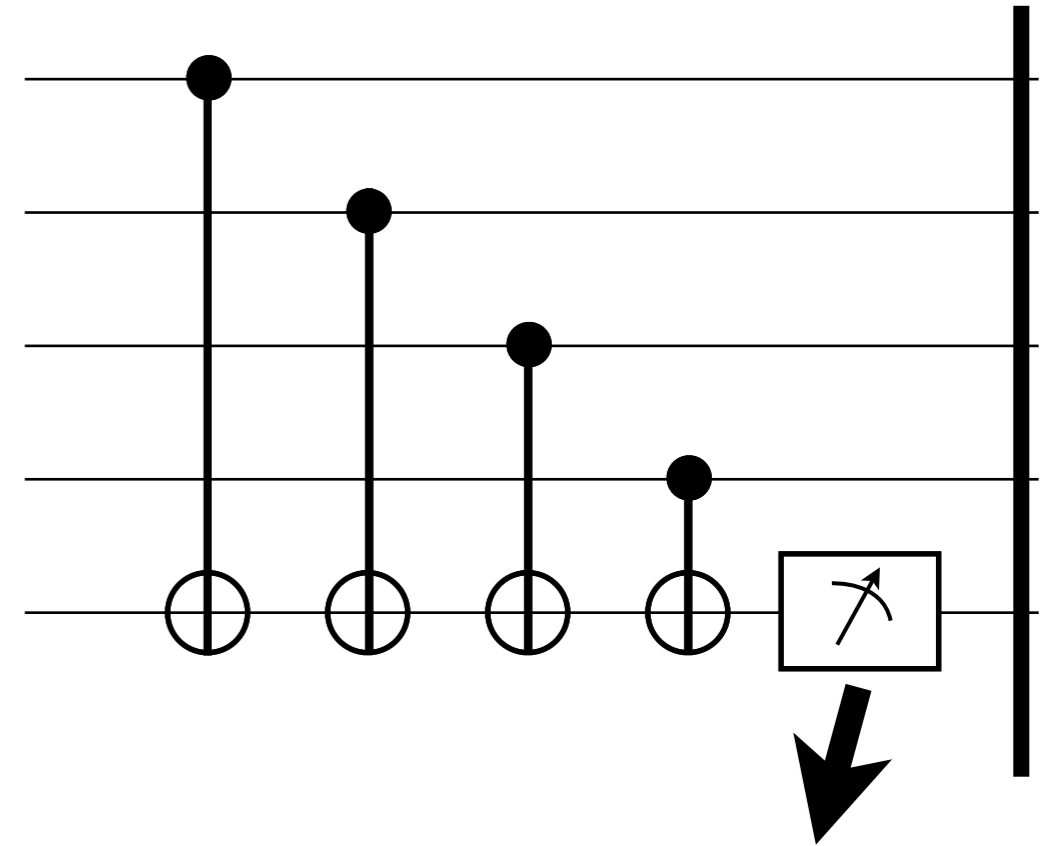
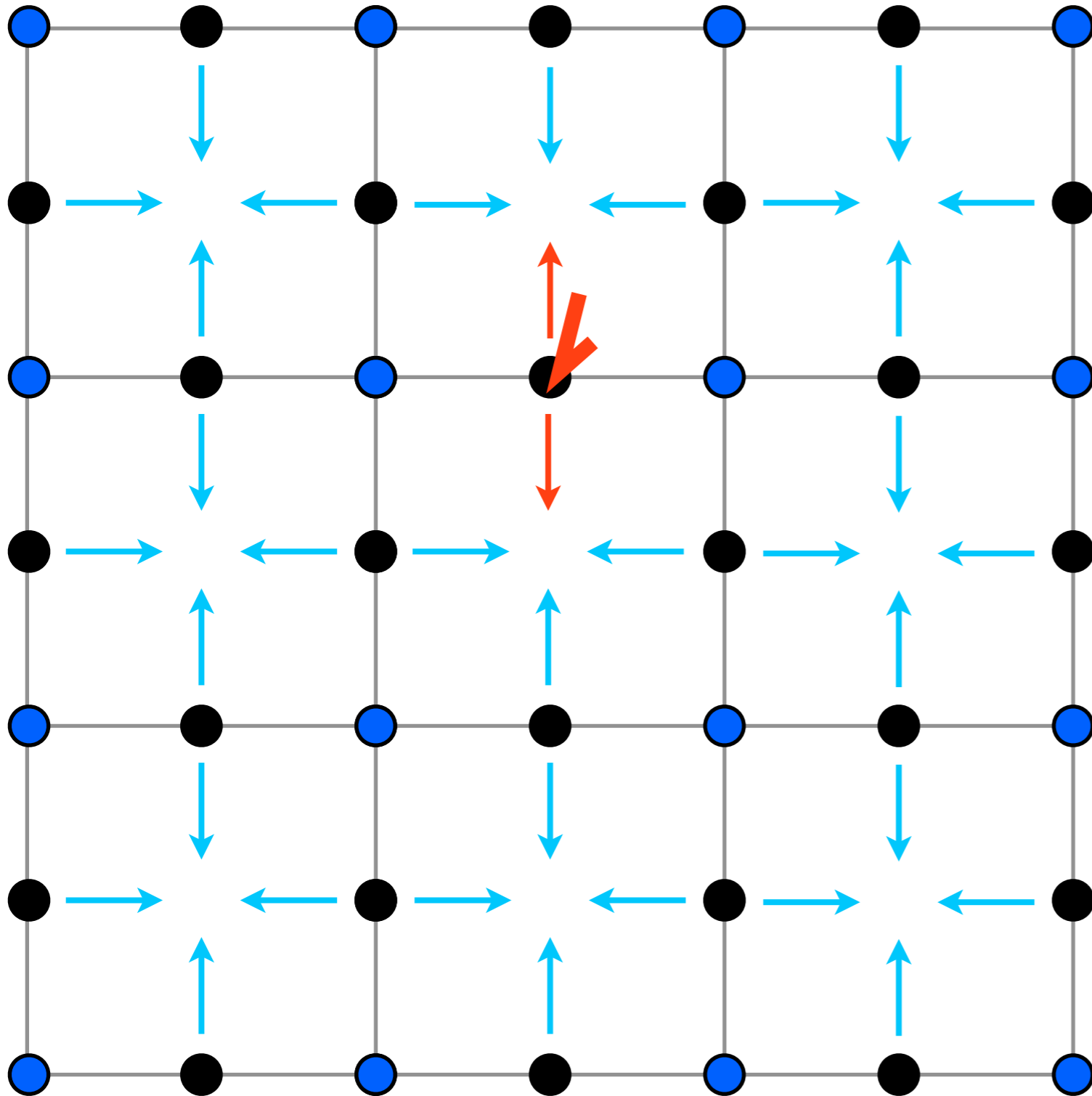
Gate Error 2

Next Step



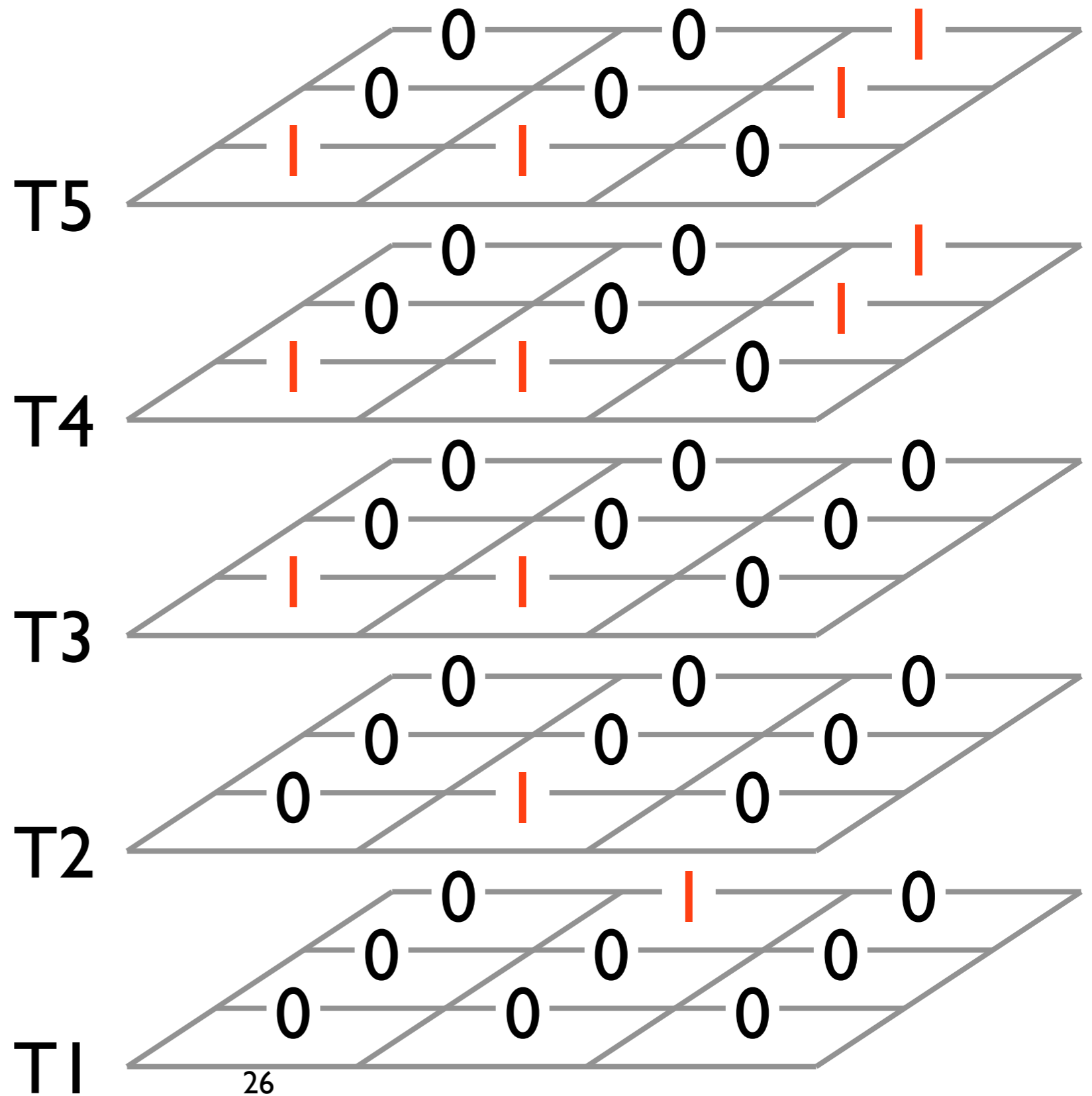
Gate Error 2

Next Step



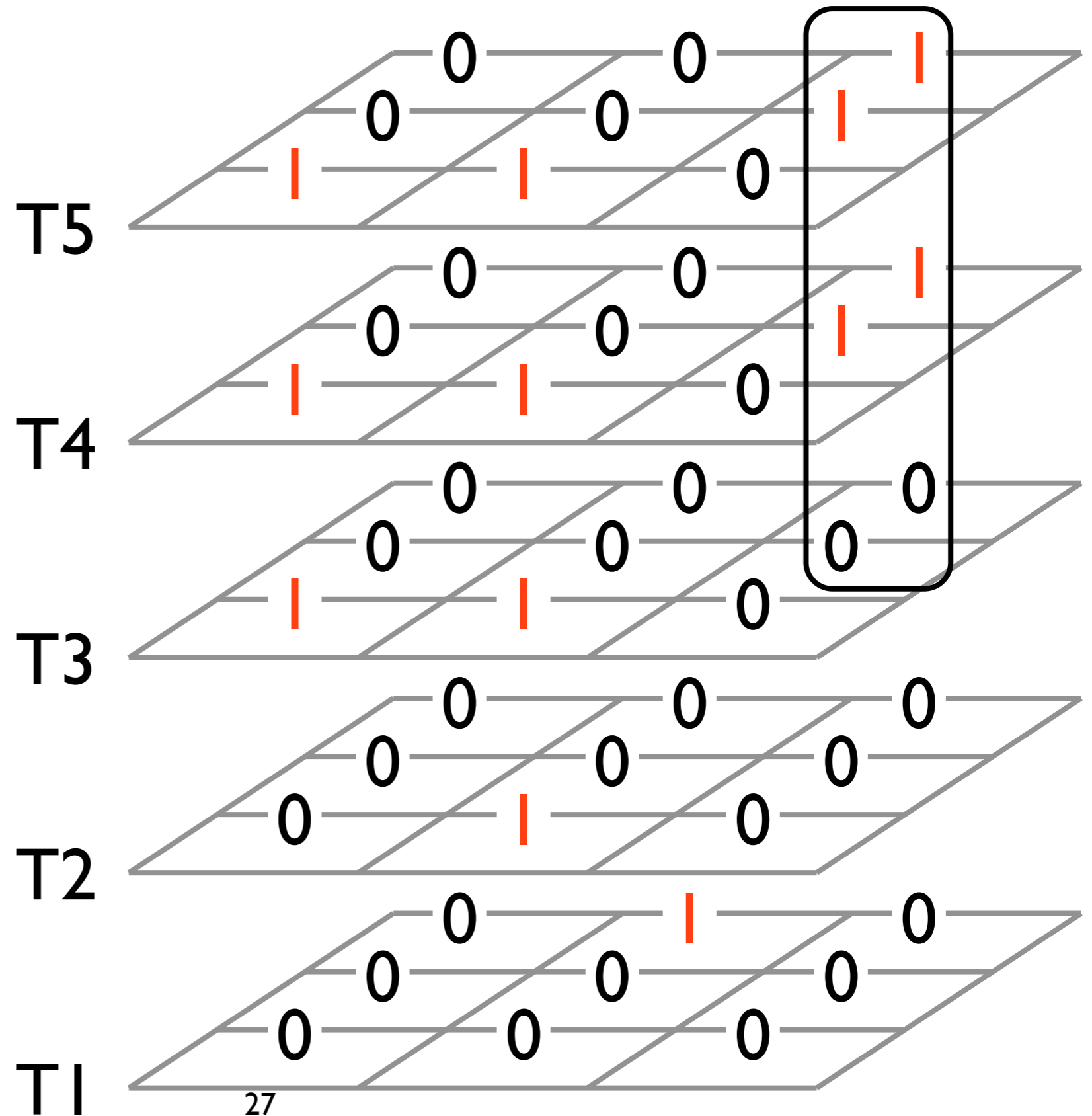
0		0
0		0
0	0	0

Error syndromes



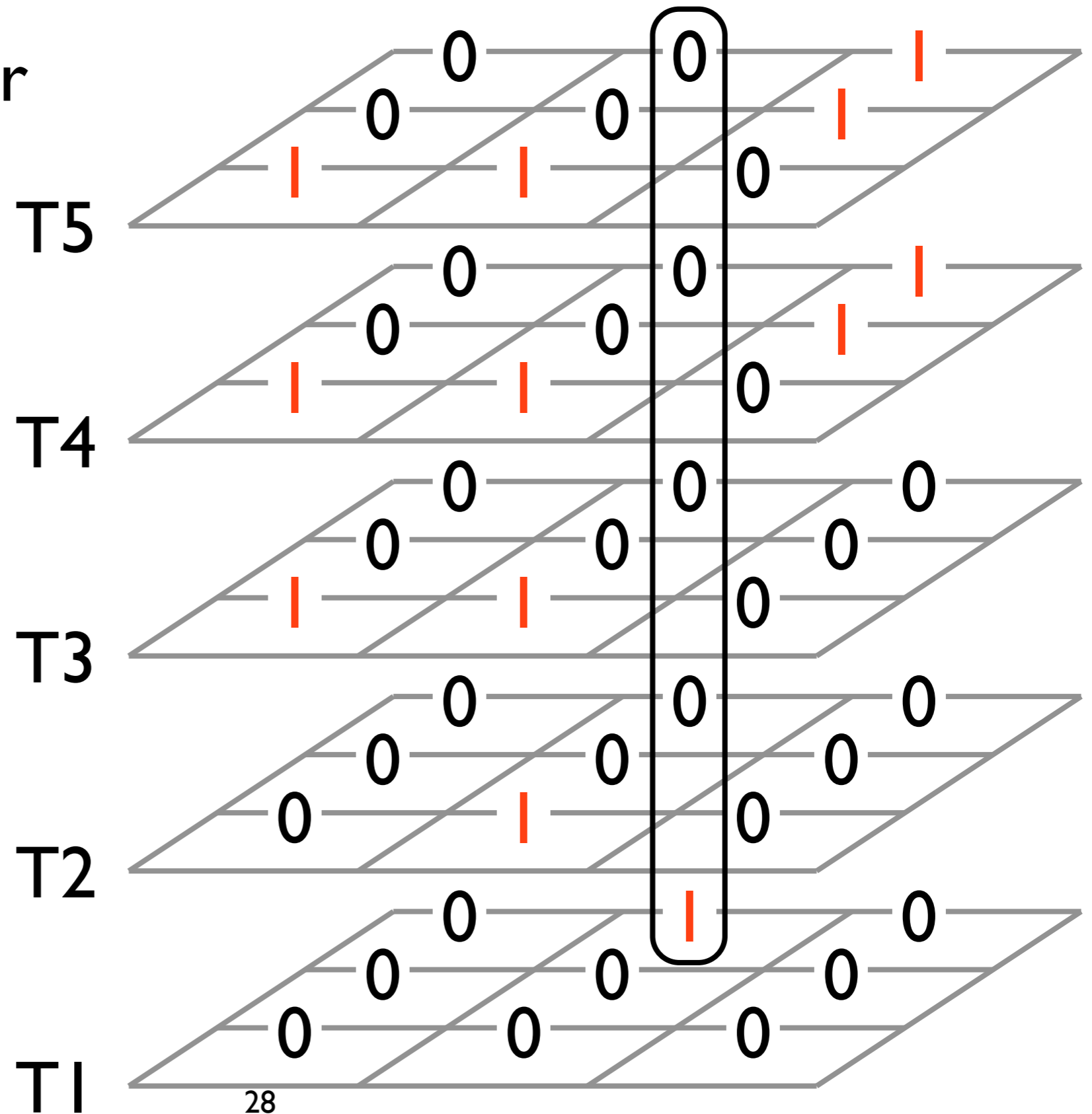
Error syndromes

memory error



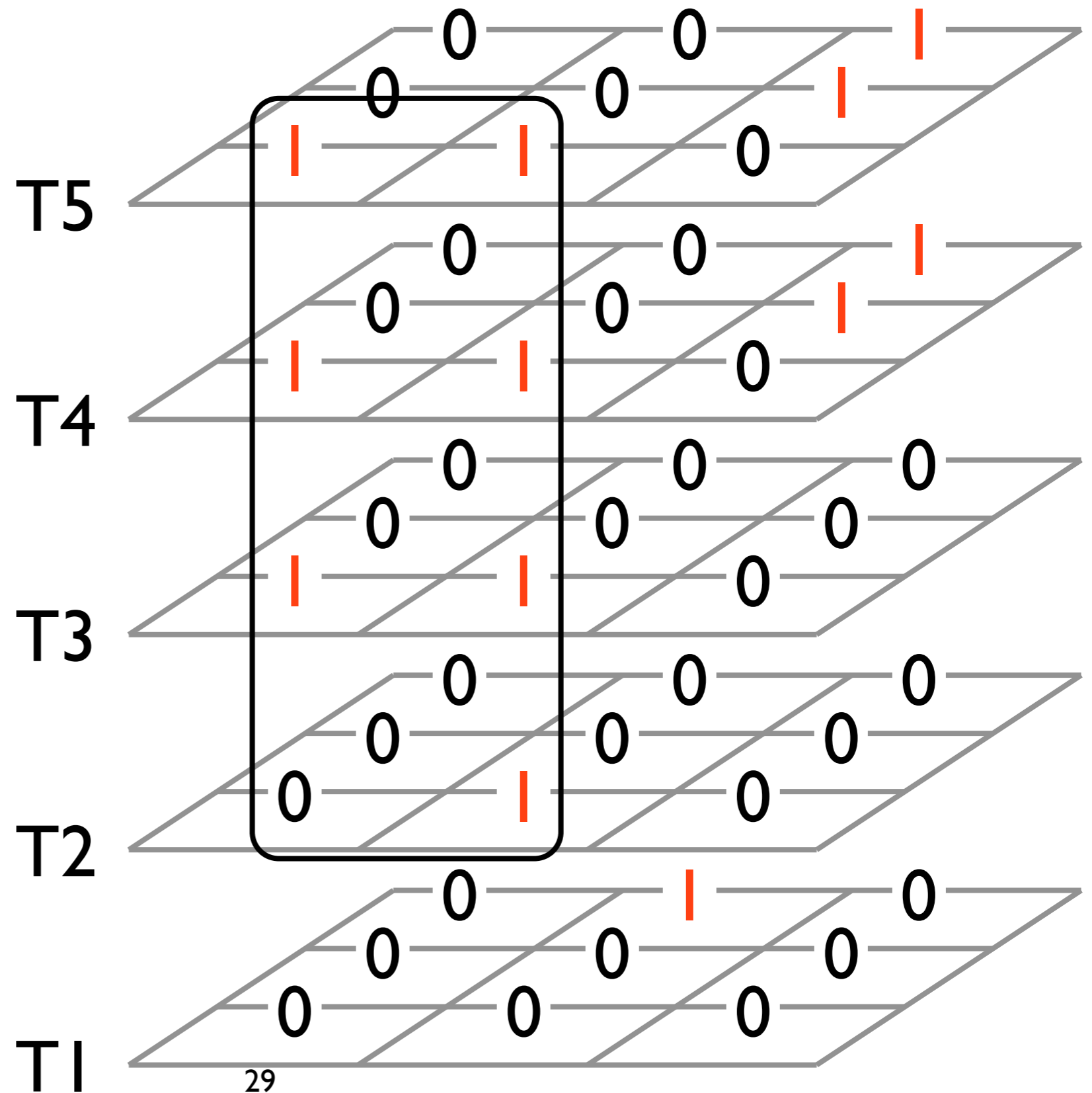
Error syndromes

measurement error



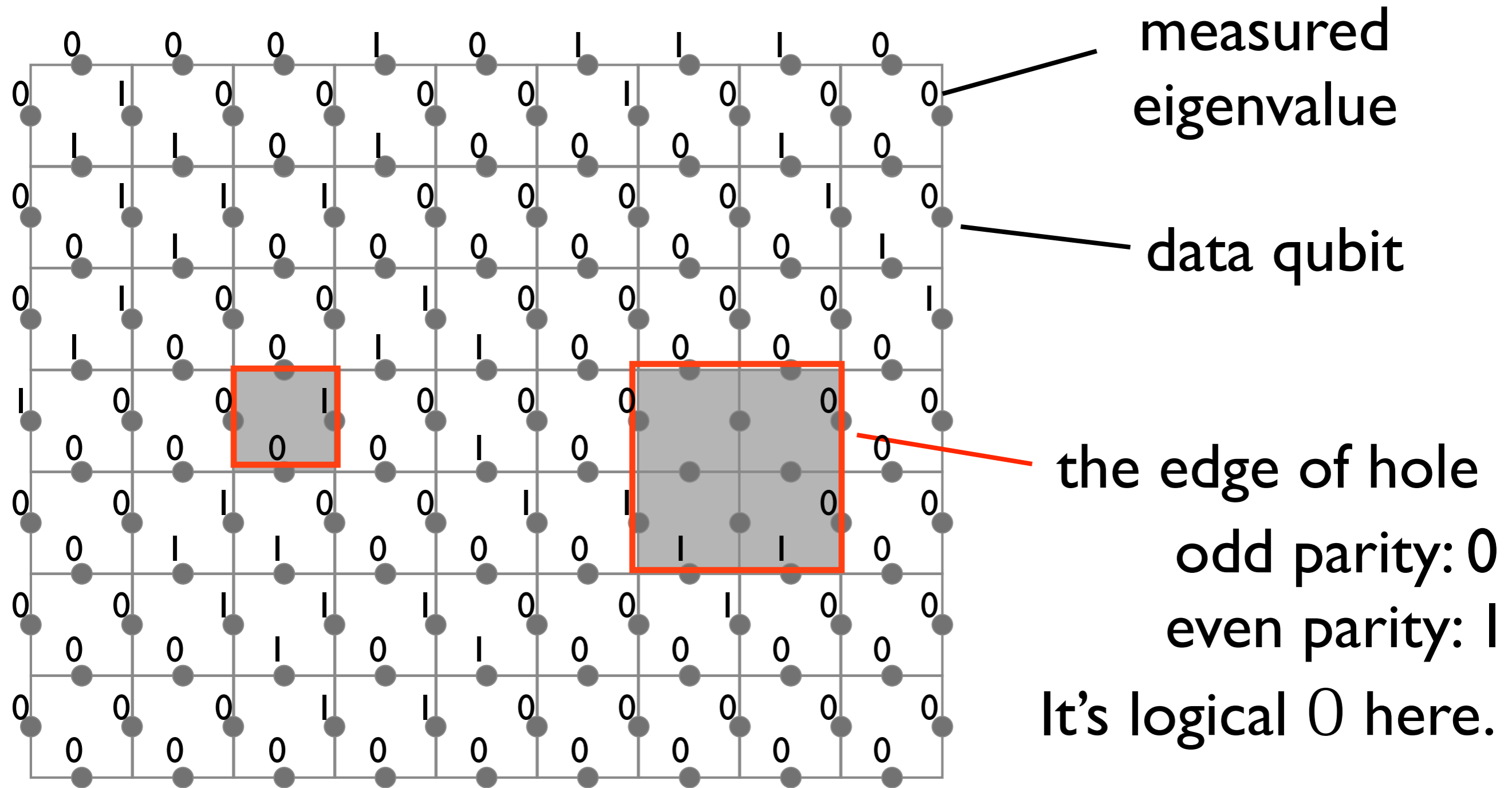
Error syndromes

gate error



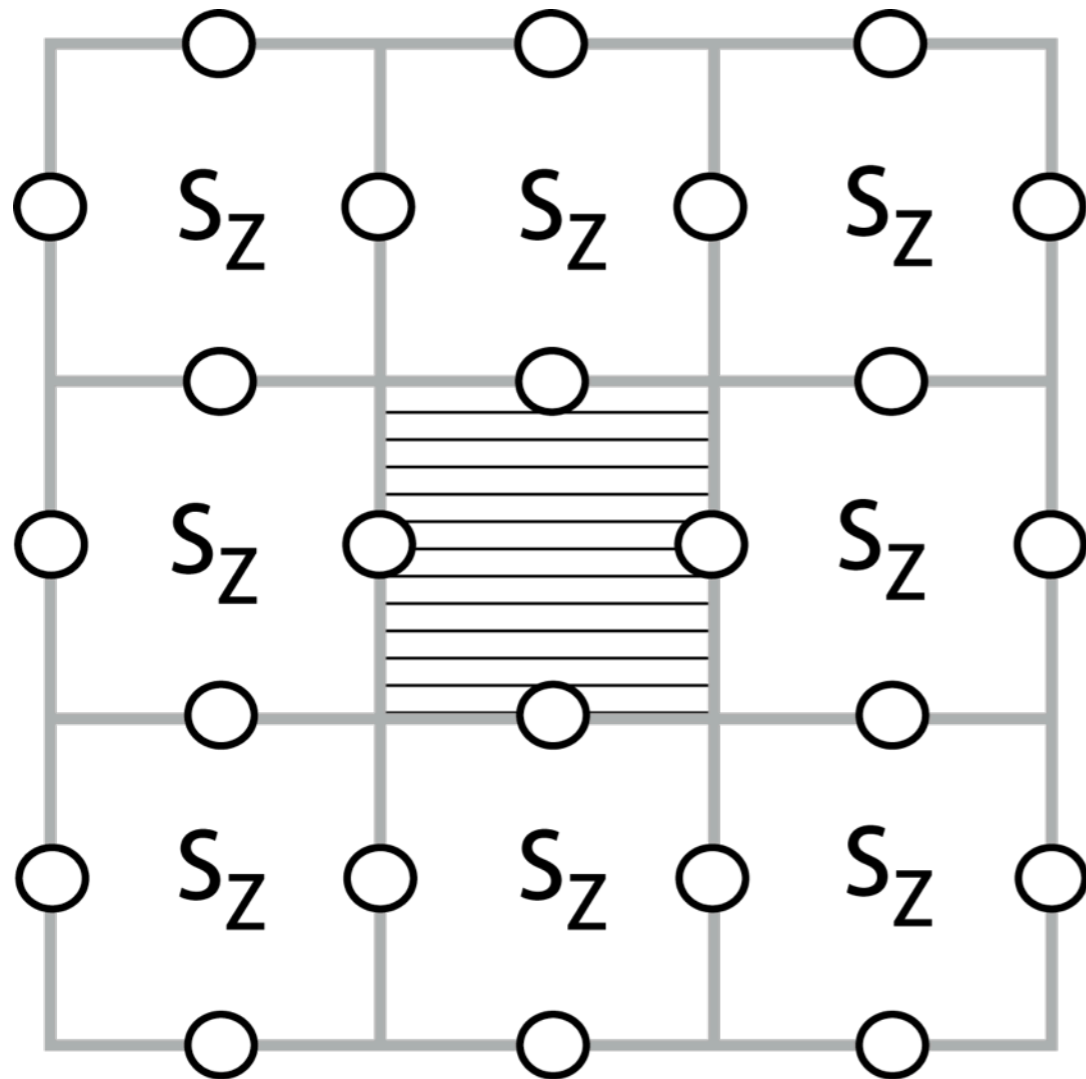
Simplest Example

~Encoding logical qubit~

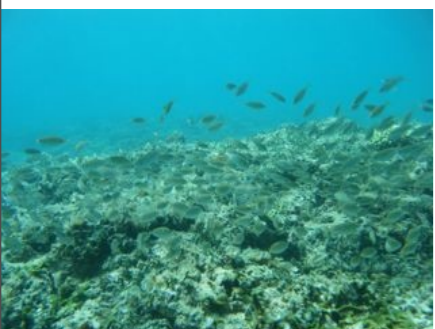




Lattice holes as logical qubits

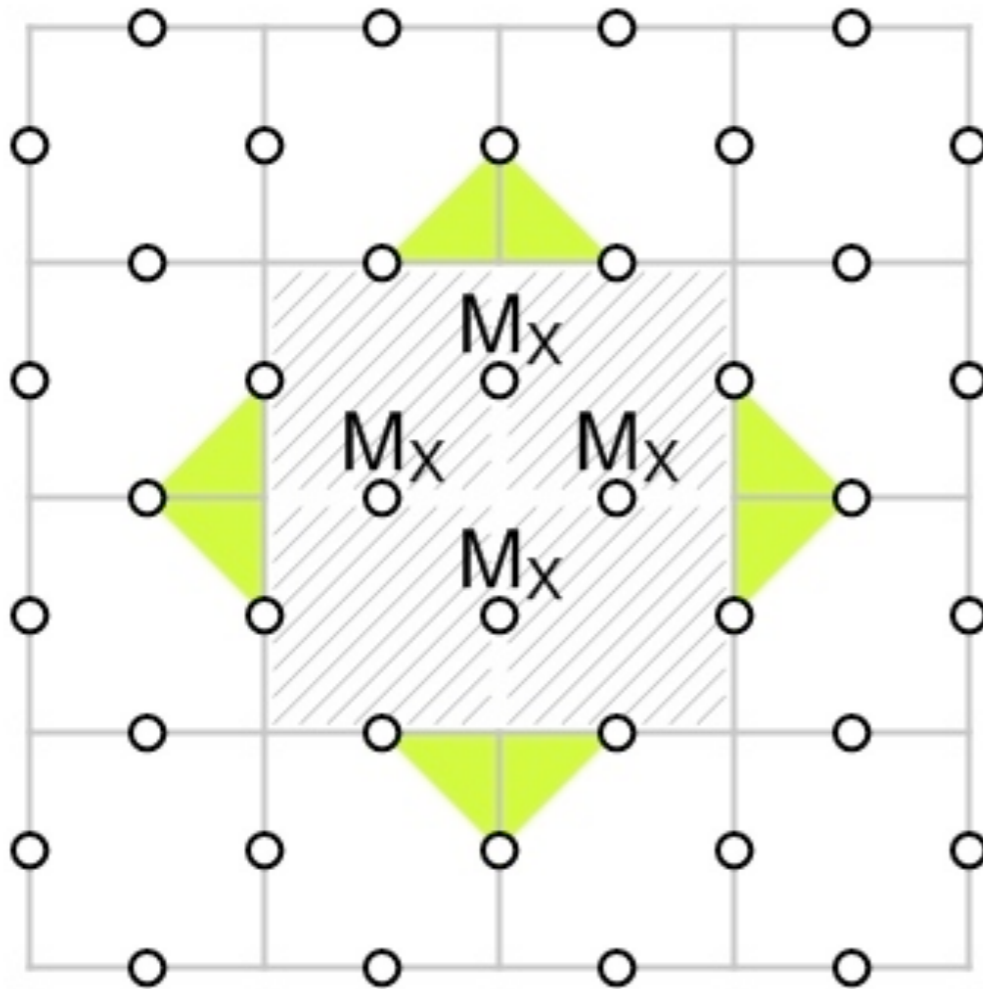


- Introduce degree of freedom by not measuring stabilizer (first, measure out any data qubits in the *interior* of the hole in X basis, to disentangle)
- Only one degree of freedom associated with arbitrary size hole
- Here, 24 lattice qubits, 9 Z stabilizers, 16 X stabilizers (but one *not independent*) = 24 independent stabilizers
- n.b.: Holes referred to as “defects” in most papers, I reserve that term for physically defective qubits

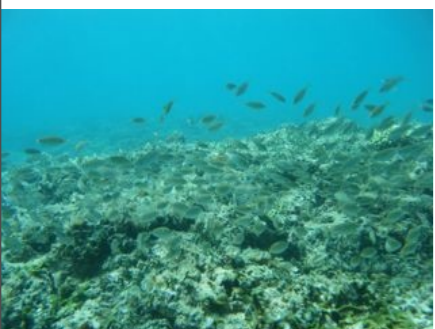




Larger holes

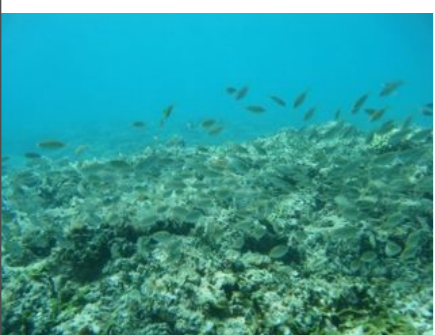
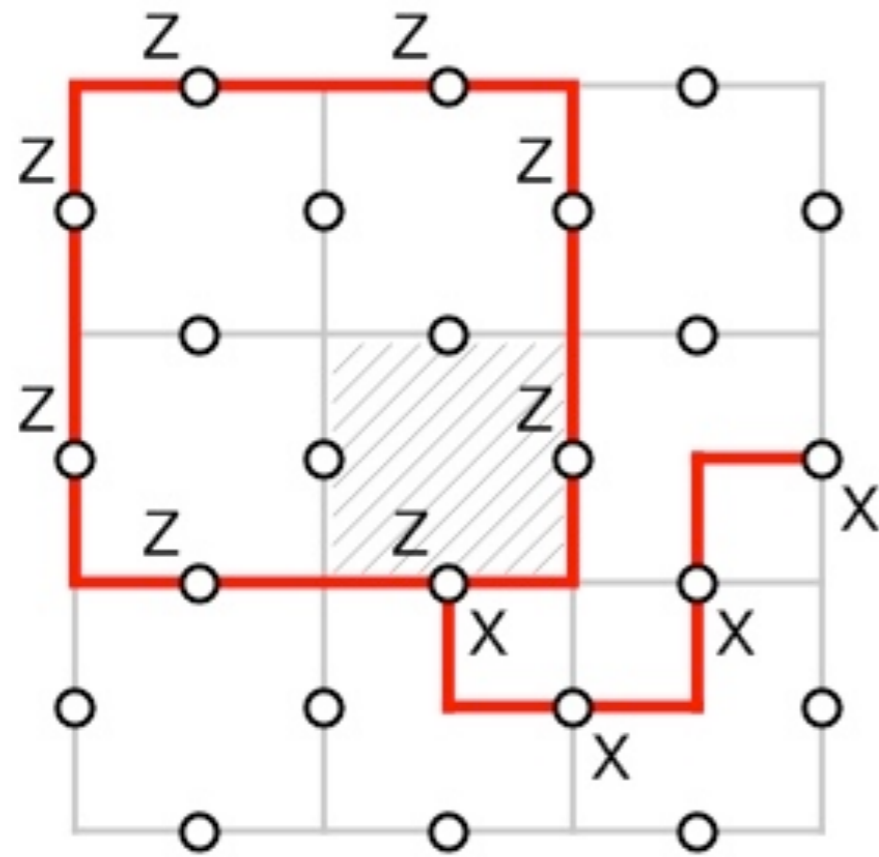


- Only one degree of freedom associated with arbitrary size hole
- Must correct three-term X stabilizers (green)
- Initialize to $|0_L\rangle$





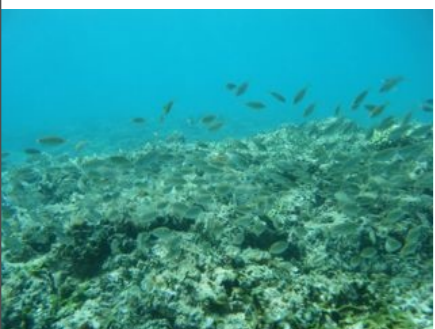
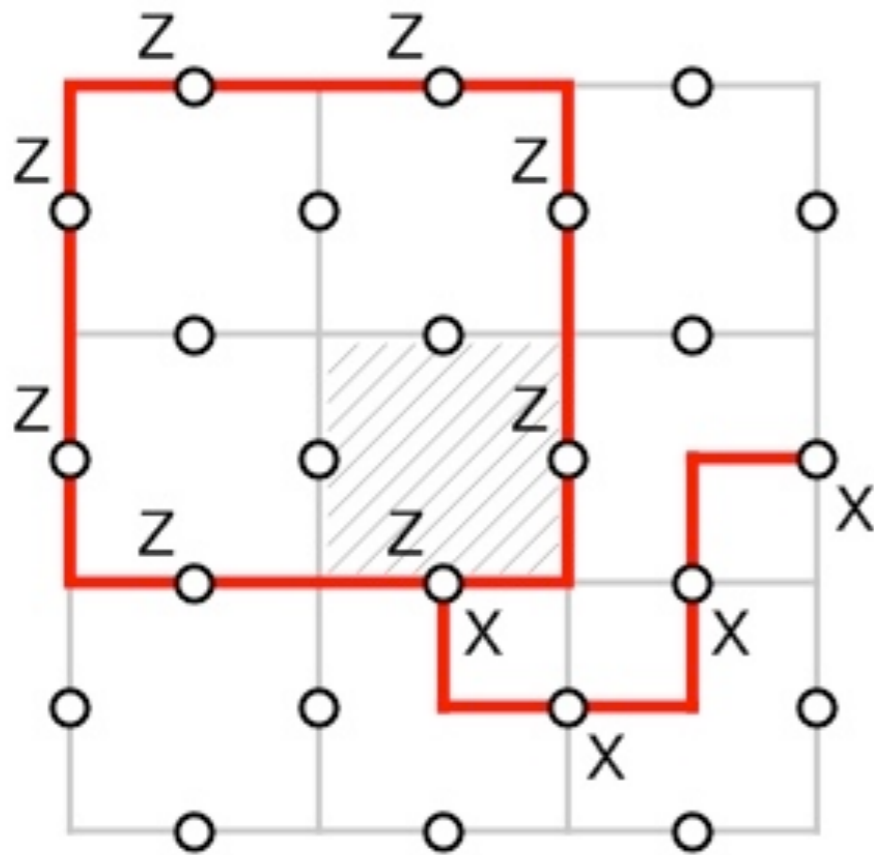
Logical X and Z gates





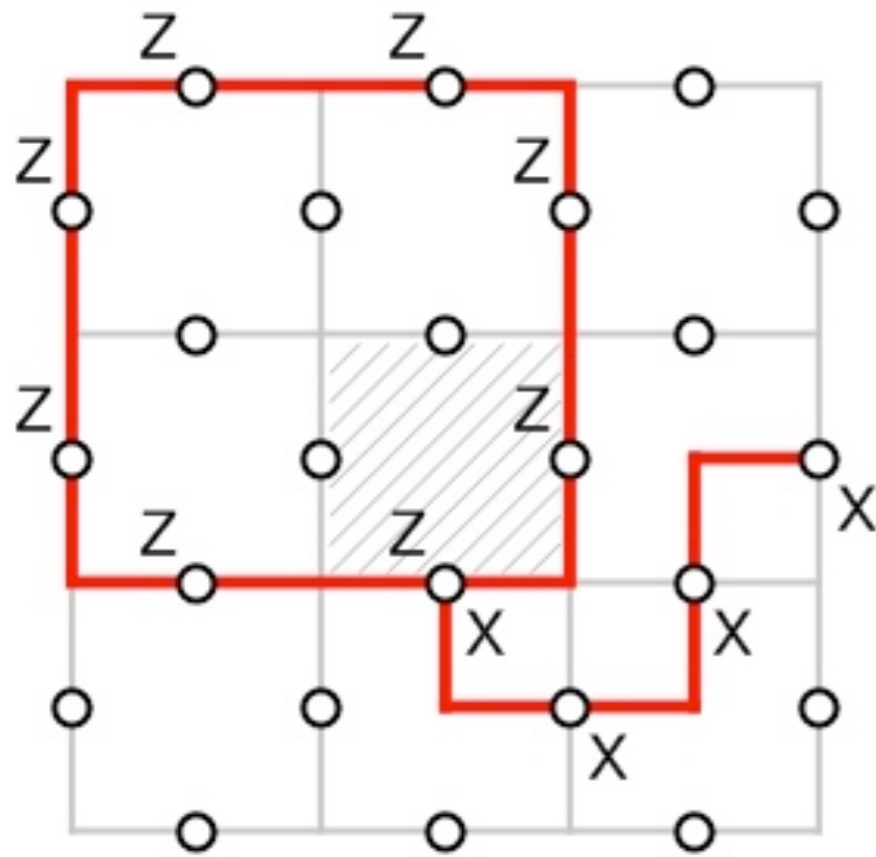
Logical X and Z gates

- Degree of freedom can be operated on by chains and rings of single-qubit X and Z *on lattice qubits*

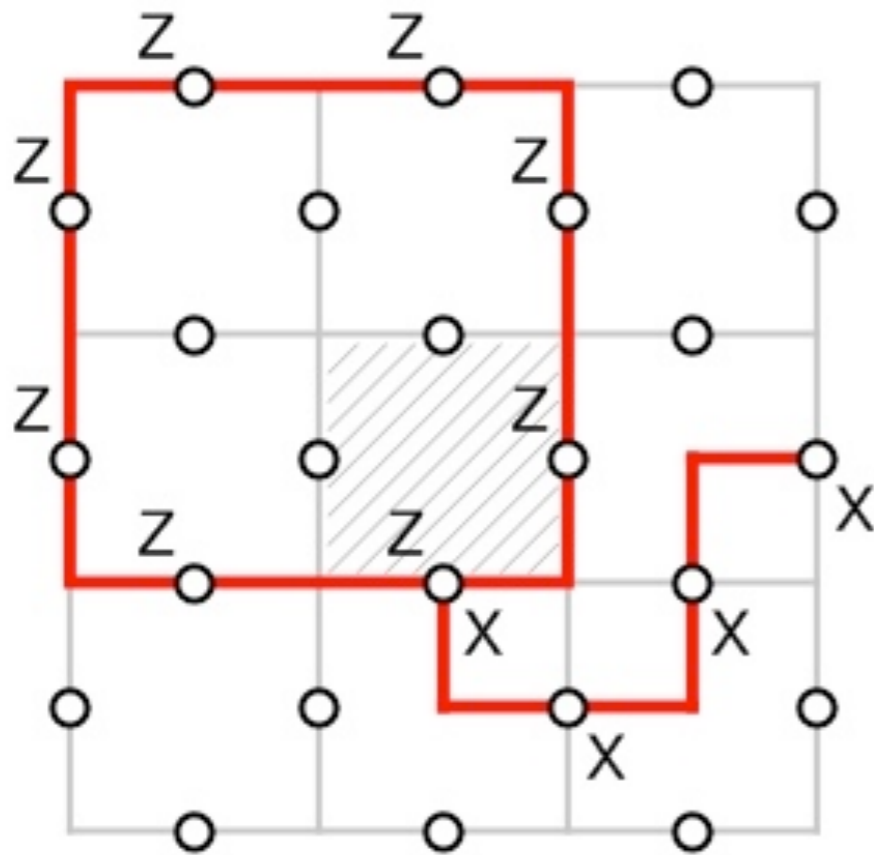


Logical X and Z gates

- Degree of freedom can be operated on by chains and rings of single-qubit X and Z *on lattice qubits*
- Chain of X flips connecting hole to boundary gives X gate on logical qubit



Logical X and Z gates

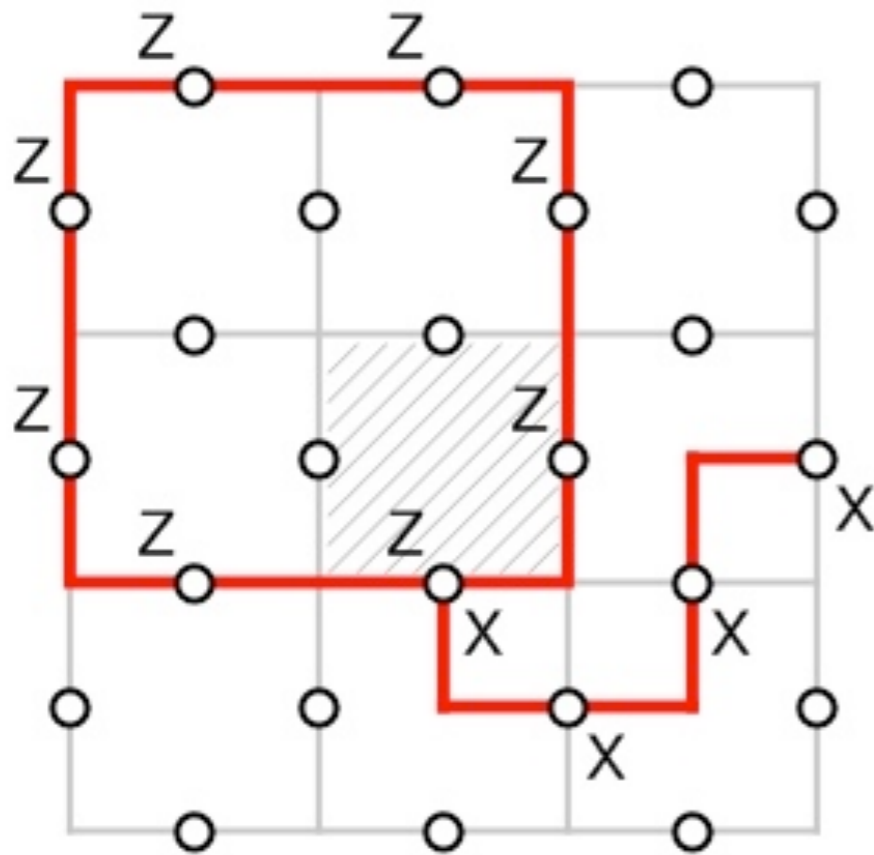


- Degree of freedom can be operated on by chains and rings of single-qubit X and Z *on lattice qubits*
- Chain of X flips connecting hole to boundary gives X gate on logical qubit
- Ring of Z flips *around* hole gives logical Z gate

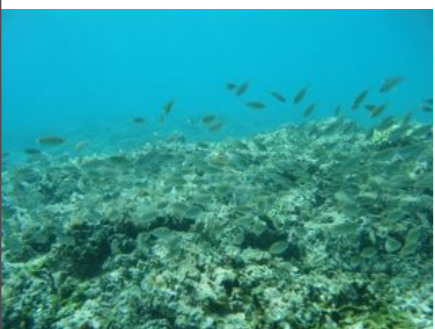




Logical X and Z gates

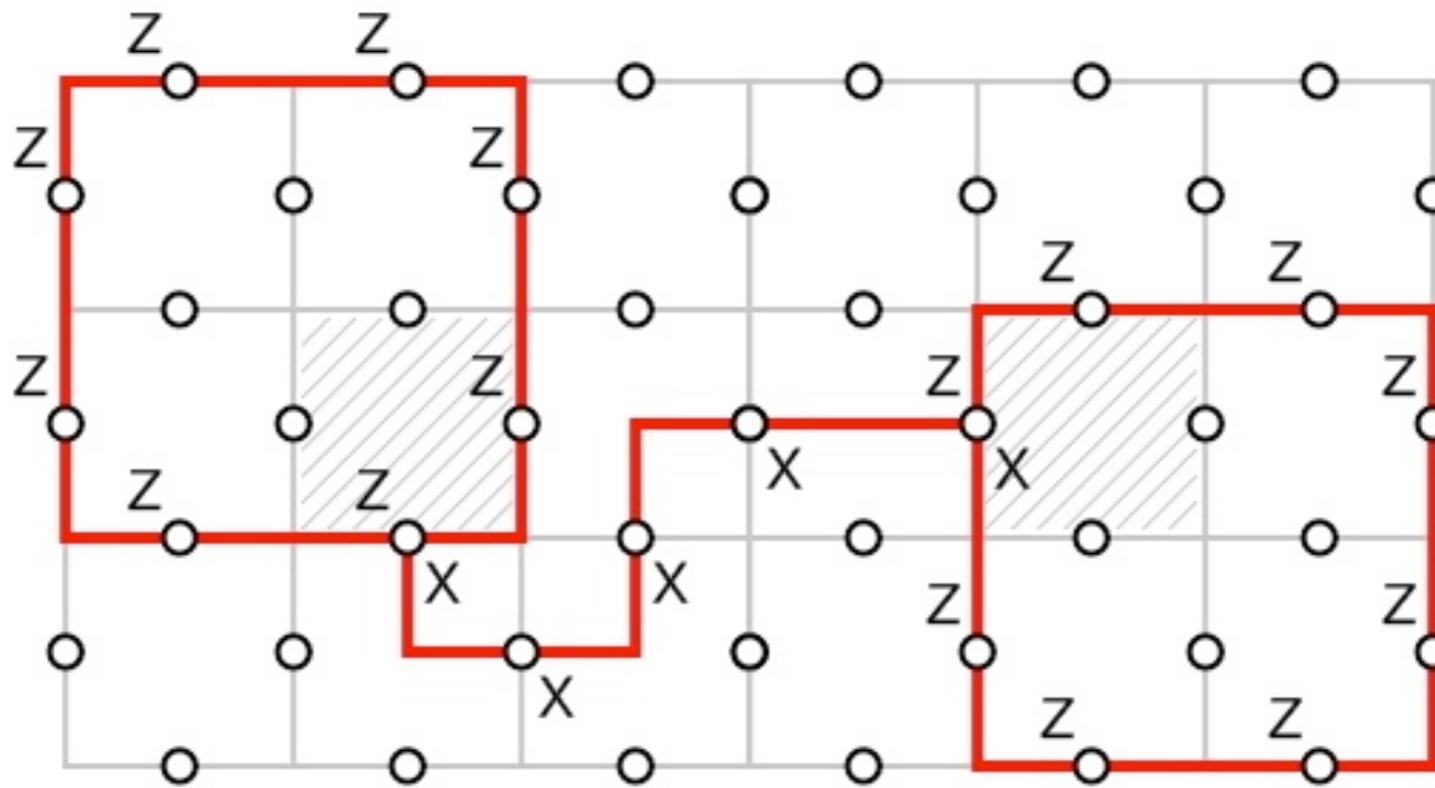


- Degree of freedom can be operated on by chains and rings of single-qubit X and Z *on lattice qubits*
- Chain of X flips connecting hole to boundary gives X gate on logical qubit
- Ring of Z flips *around* hole gives logical Z gate
- n.b.: This can be *any* ring around the hole or *any* chain connecting hole to lattice edge

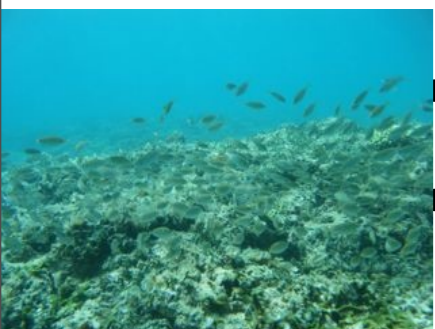




Pairing up holes

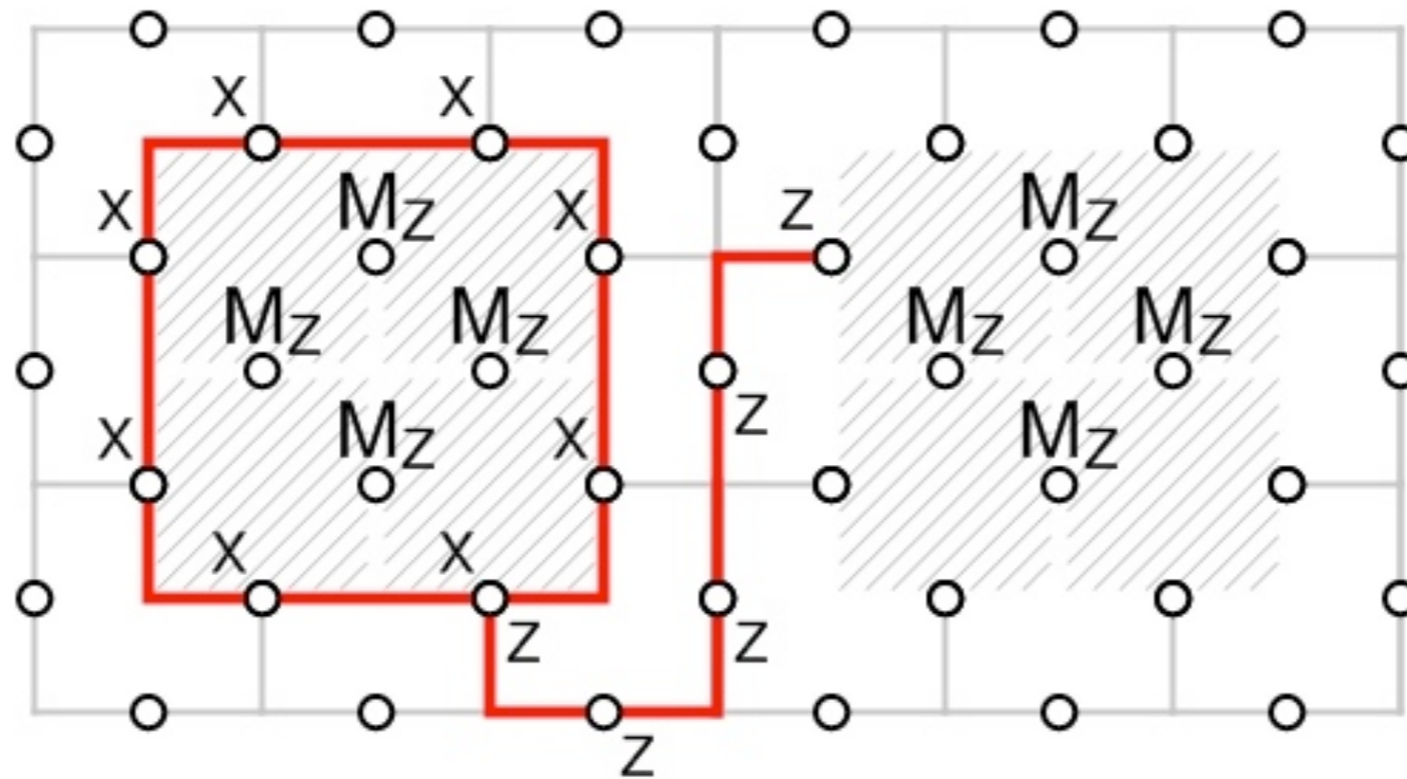


- Better to use a pair of defects to represent one qubit
 - No need to connect operators to boundary
 - Easier to move qubit around computer
- Independent adjustment of X/Z error correction strength
- Terminology: **smooth qubit**
- Logical X gate chain connects the holes
- Logical Z gate circles *one* of the holes





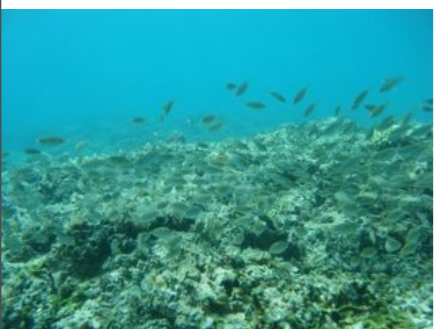
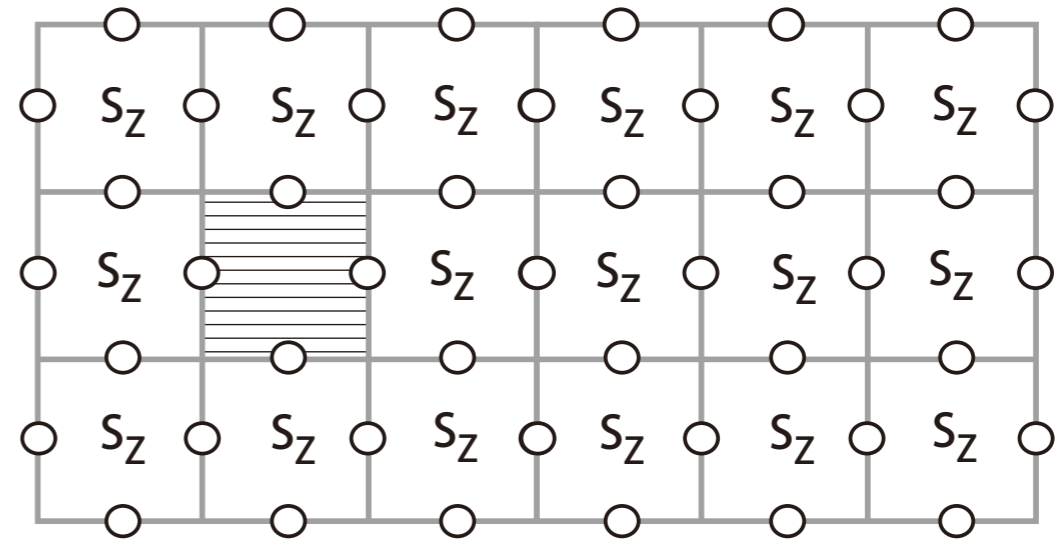
Rough qubits (Dual lattice)



- Can make a different type of qubit
- X_L and Z_L reversed
- Primal and dual lattice



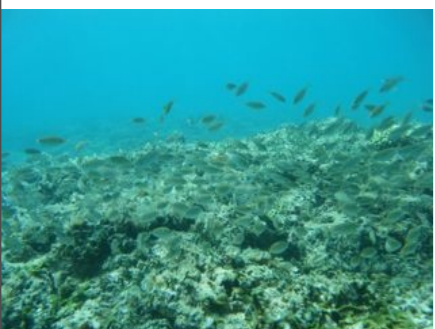
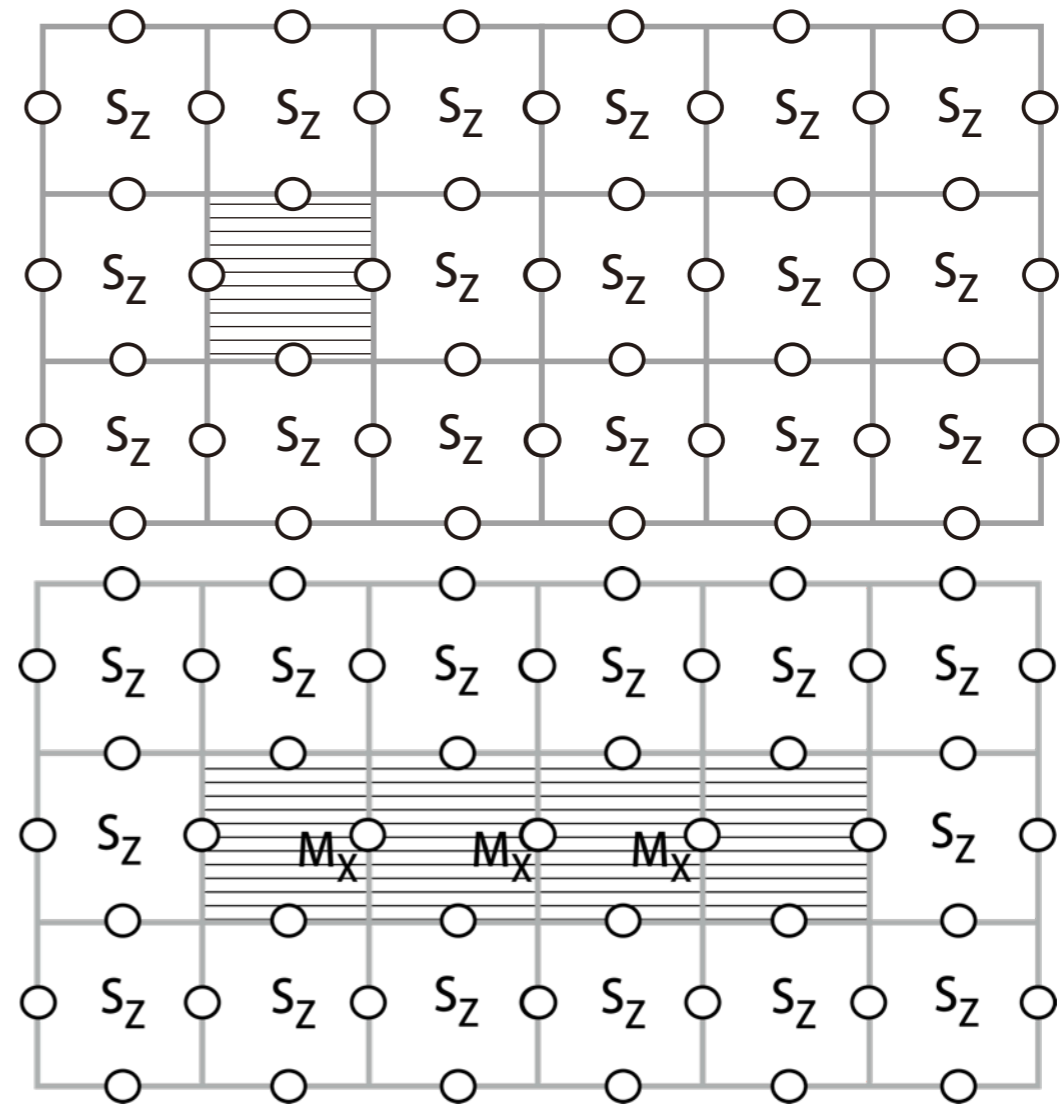
Moving holes





Moving holes

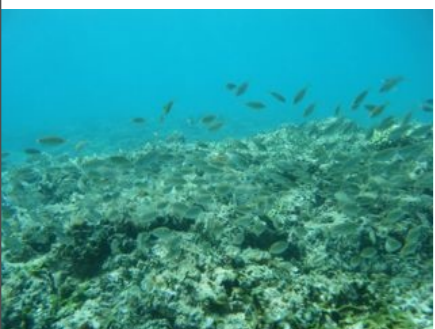
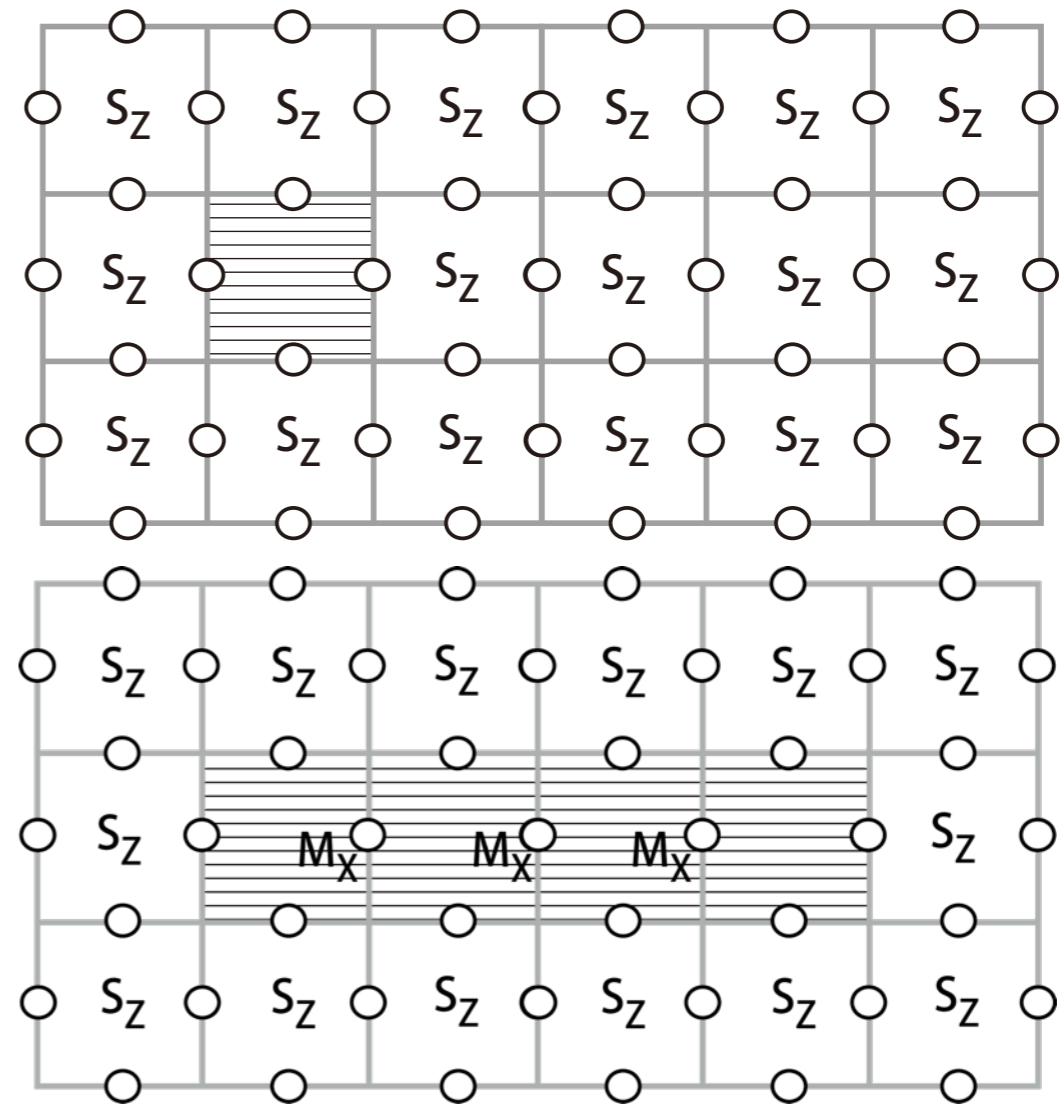
- First, expand hole using measurements and stop measuring stabilizers





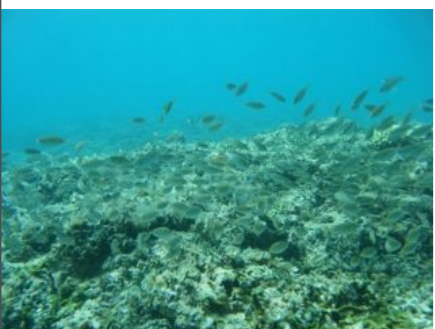
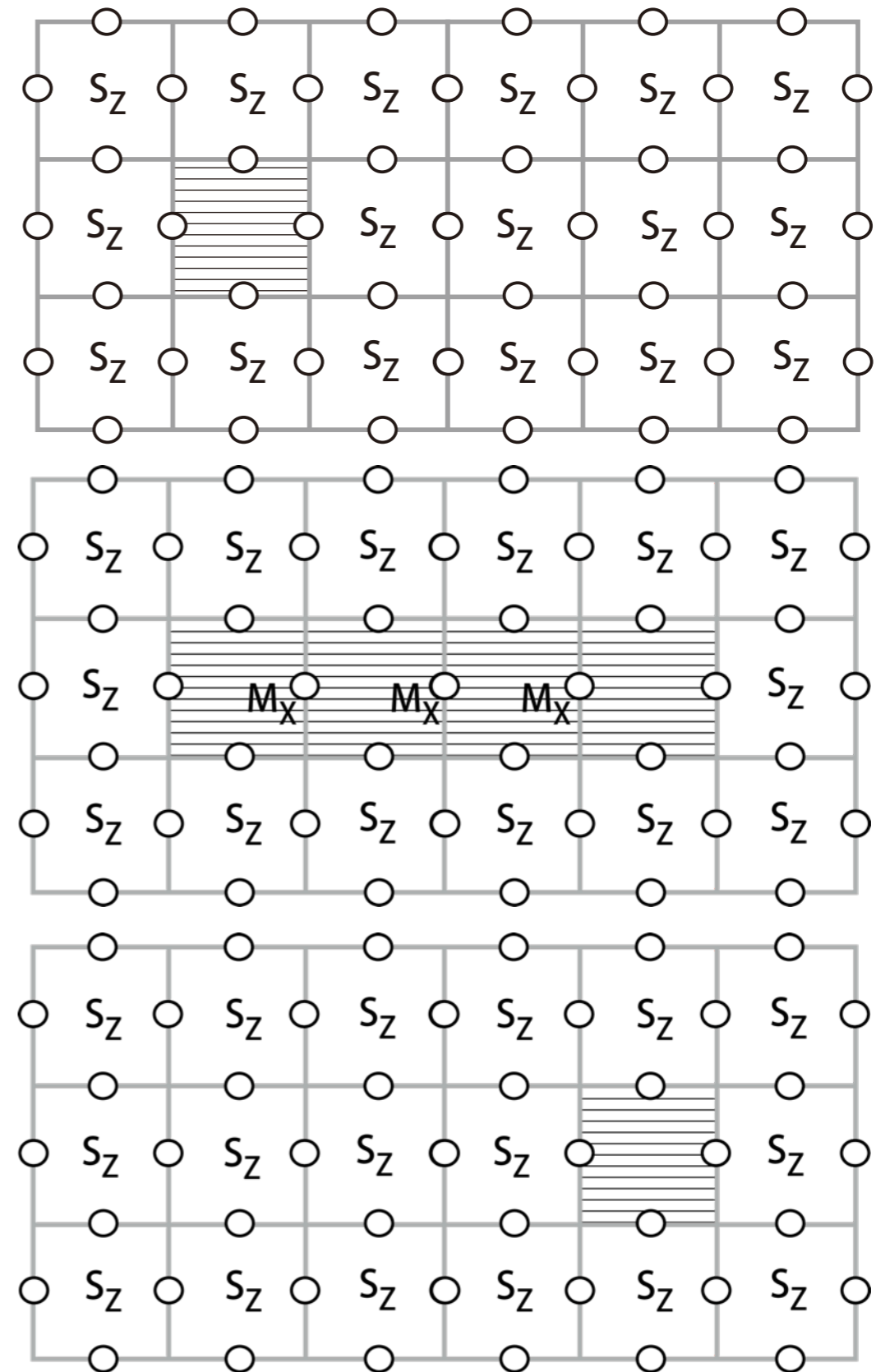
Moving holes

- First, expand hole using measurements and stop measuring stabilizers
- Fix up w/ bit flips (takes multiple rounds, depending on distance)



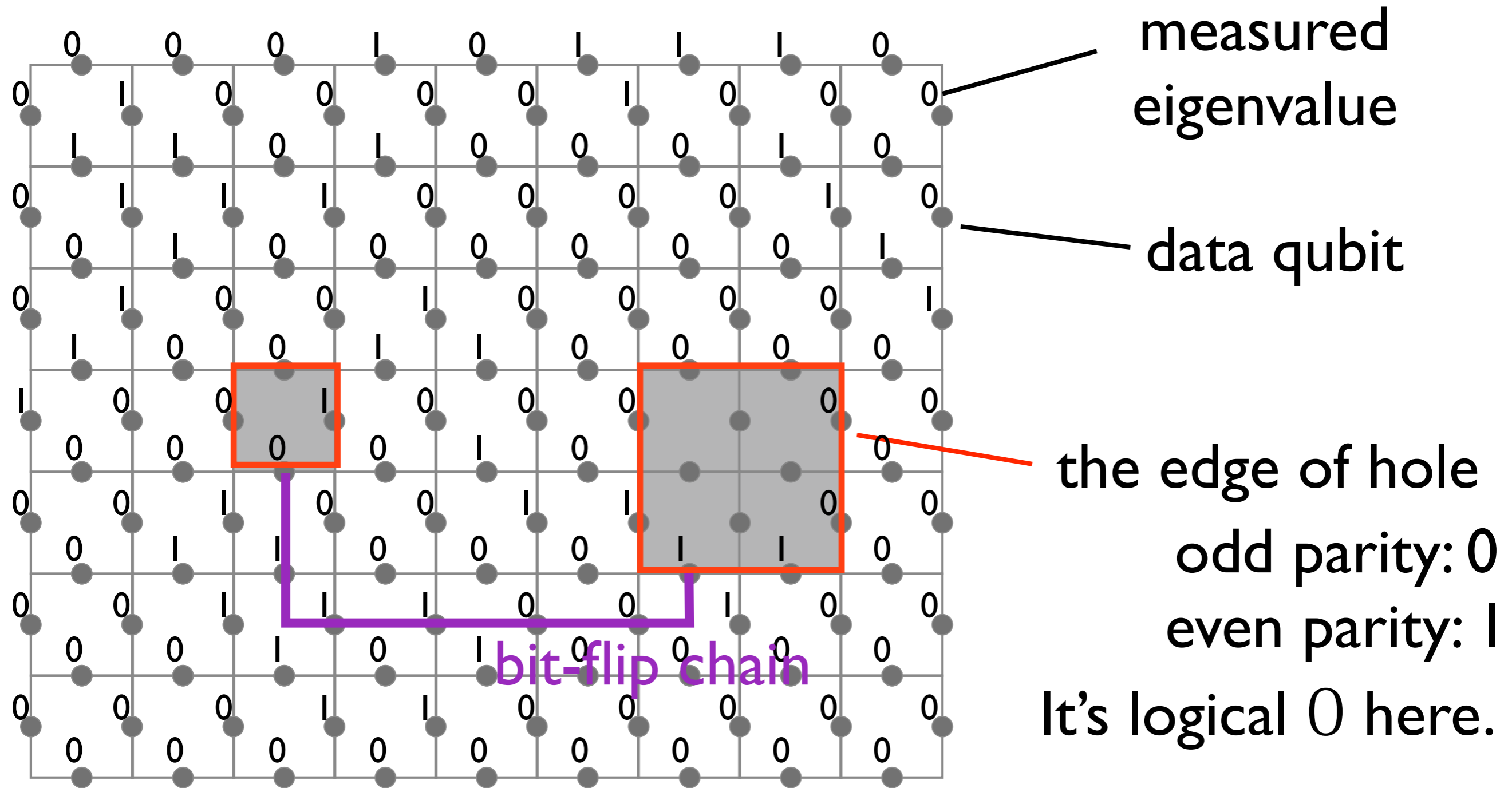
Moving holes

- First, expand hole using measurements and stop measuring stabilizers
- Fix up w/ bit flips (takes multiple rounds, depending on distance)
- Shrink hole to new location



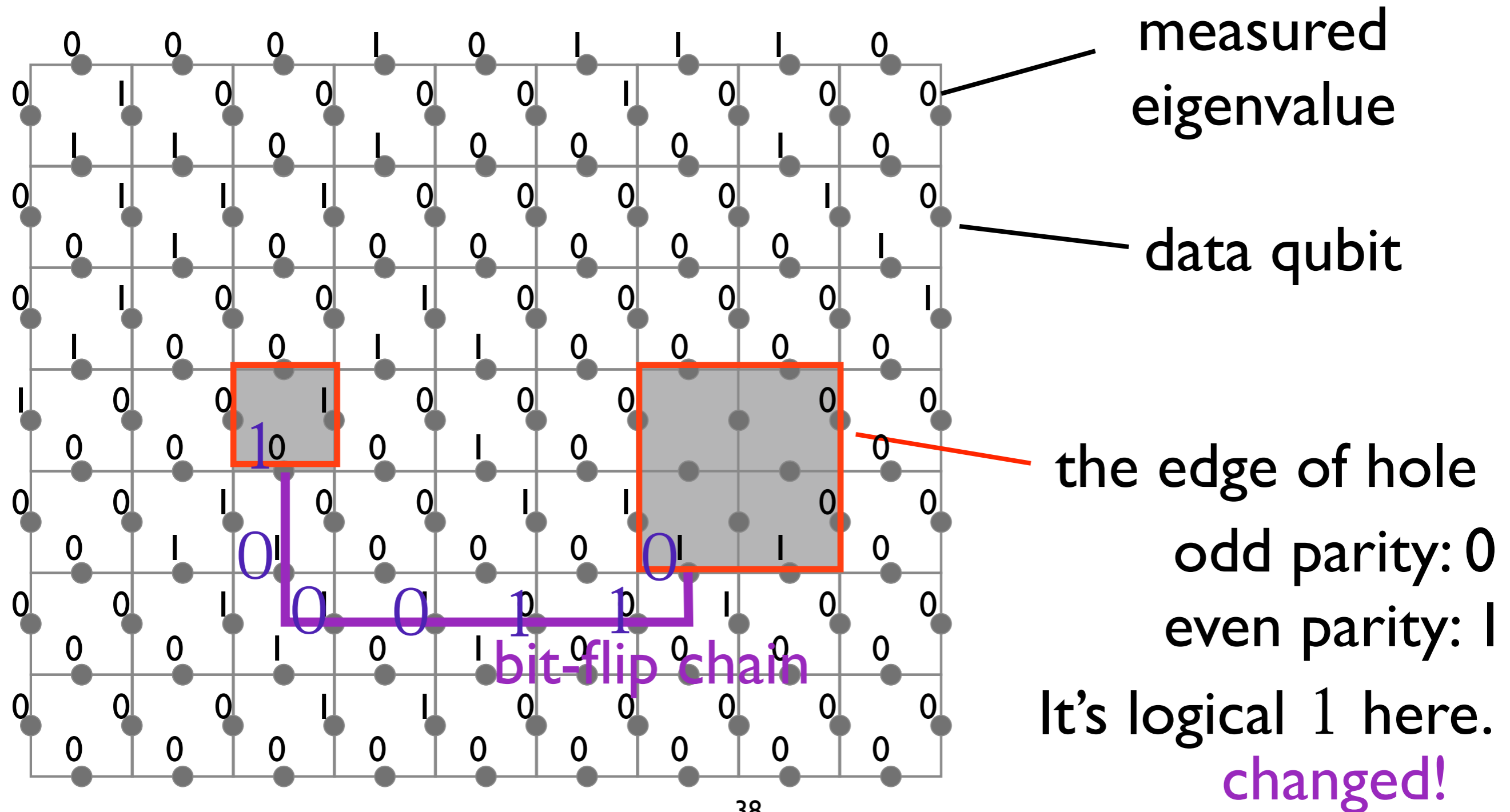
Simplest Example

~NOT gate operation~



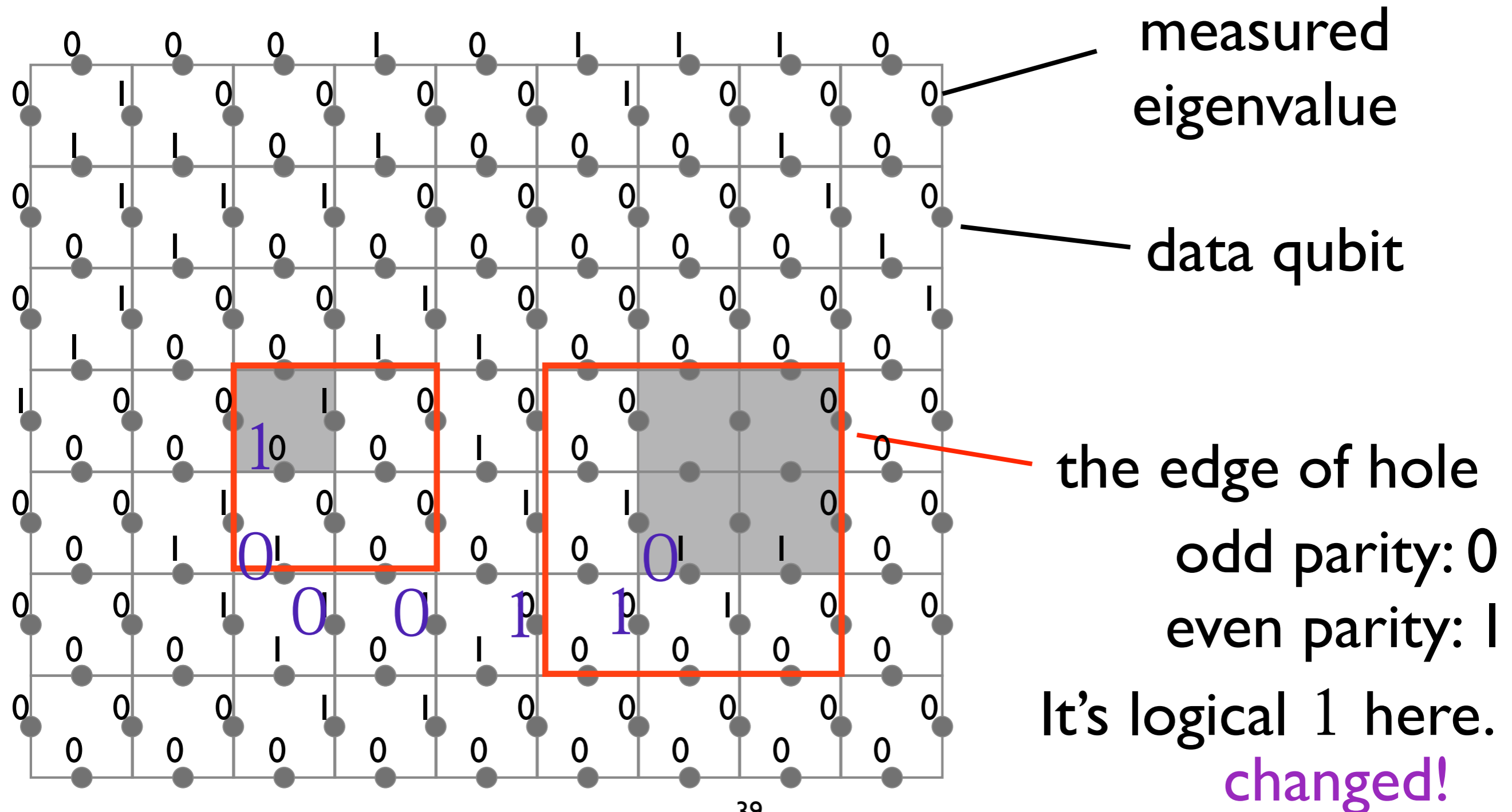
Simplest Example

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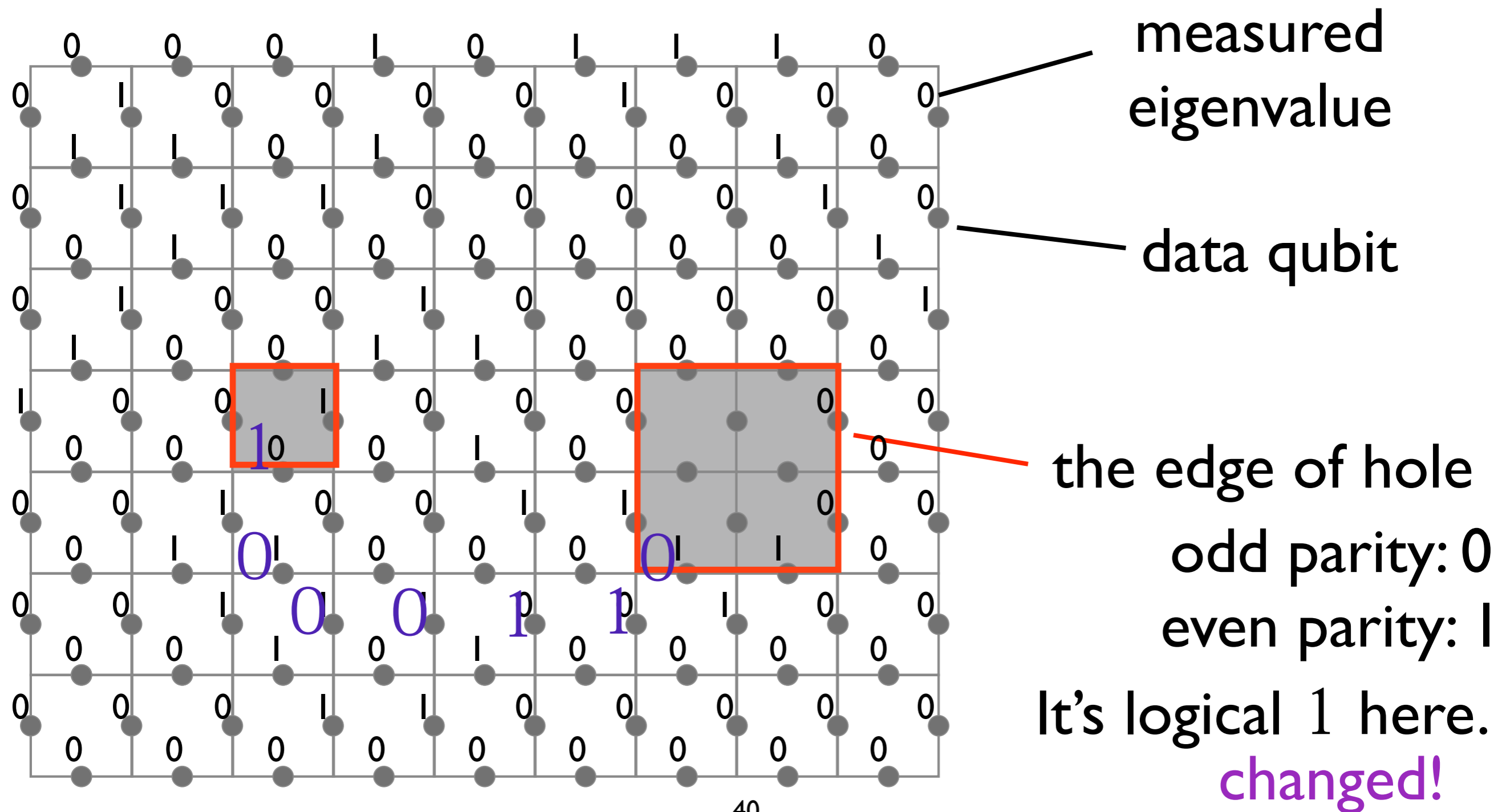
Simplest Example

~NOT gate operation~



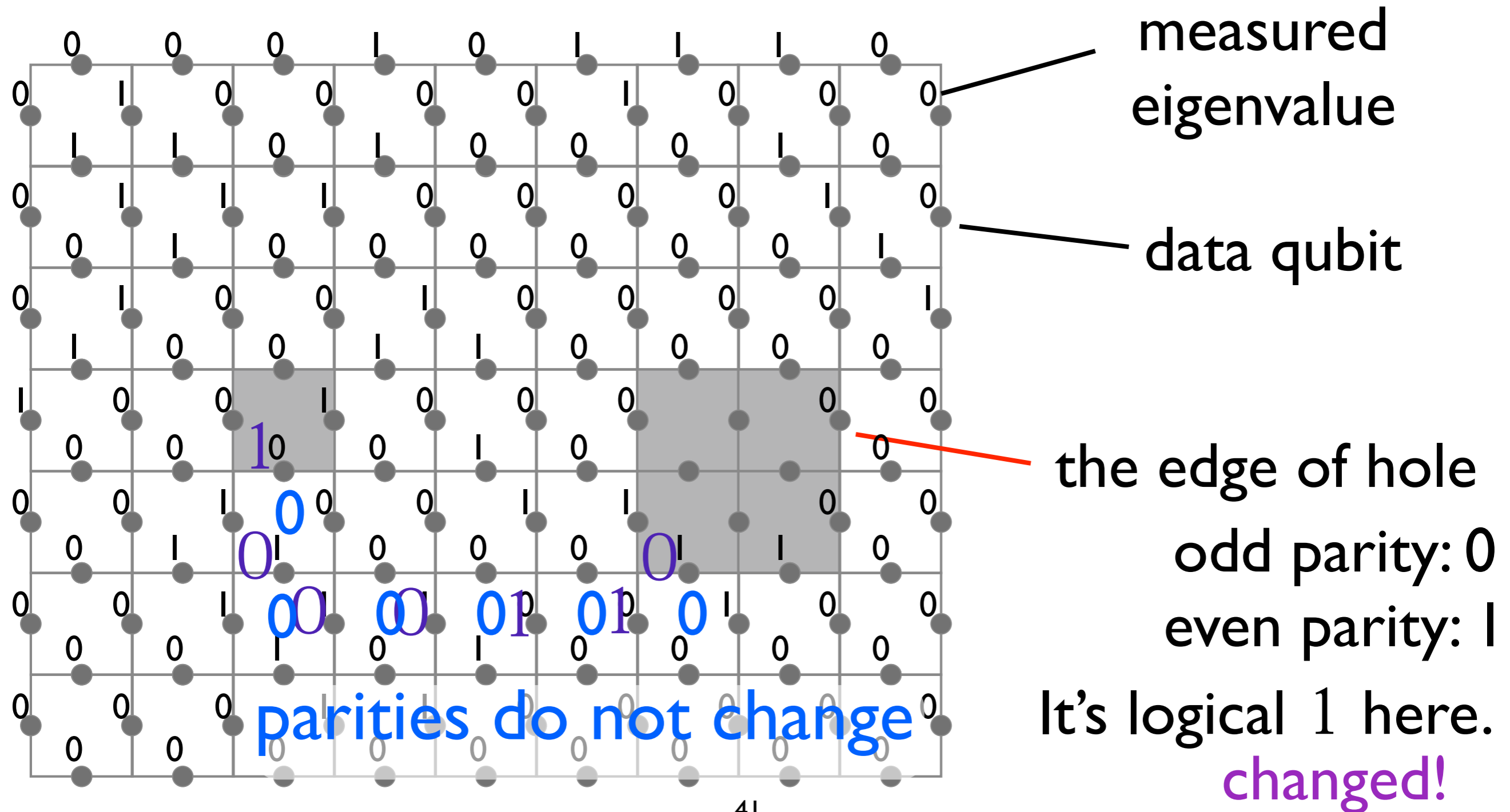
Simplest Example

~NOT gate operation~



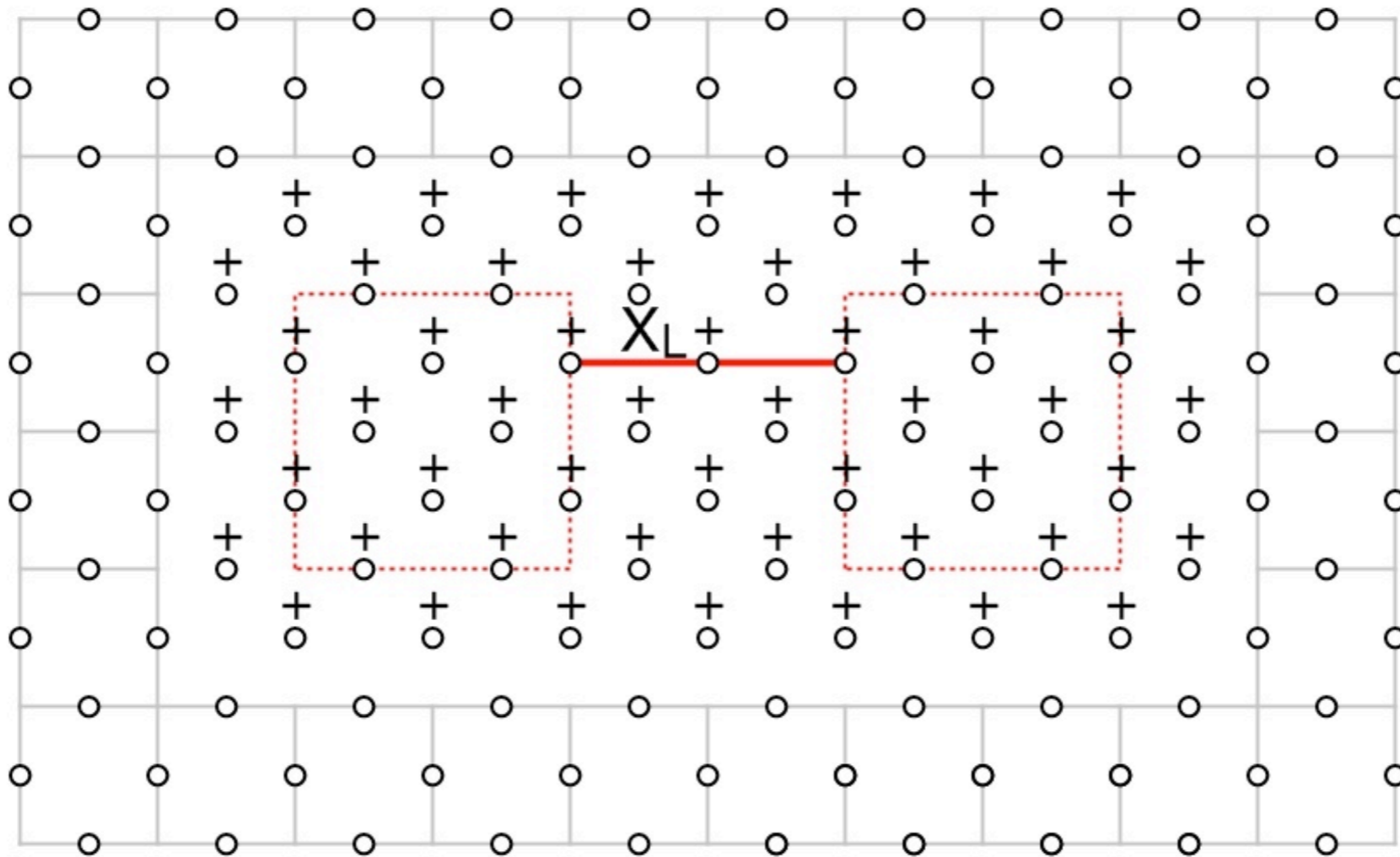
Simplest Example

~NOT gate operation~





Initializing logical qubits



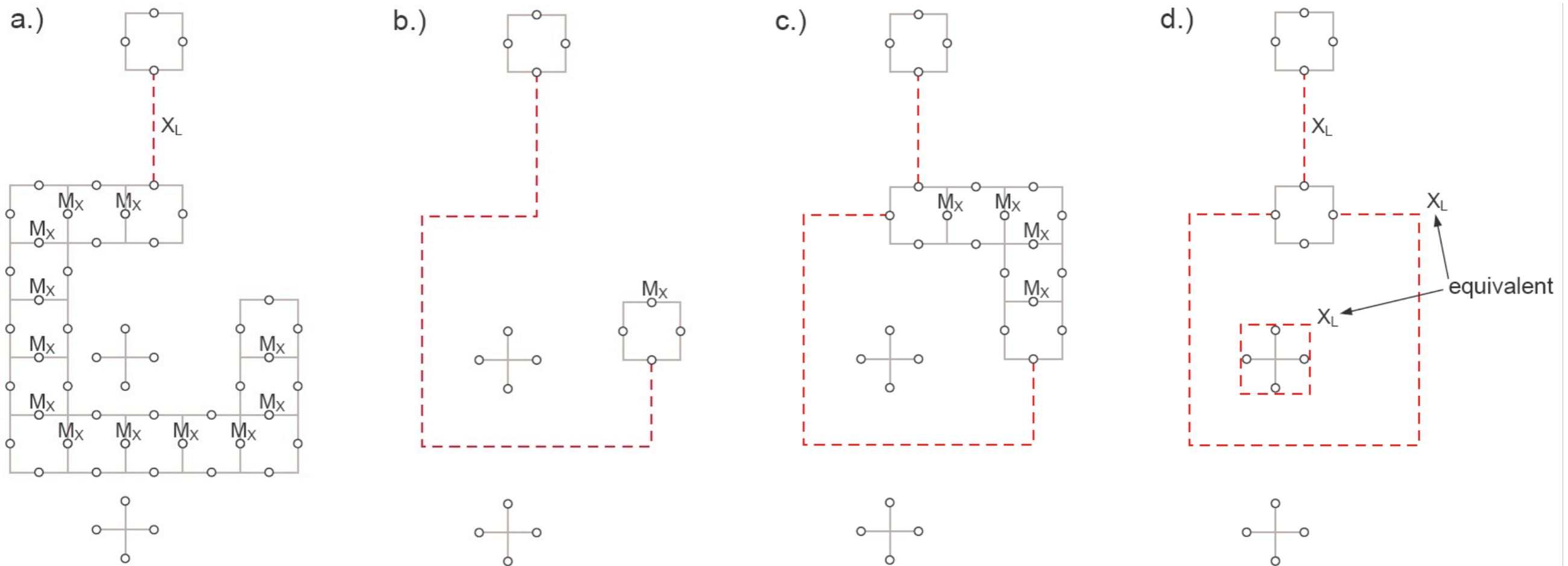
- Start with region of $|+\rangle$
- Smooth qubit in ± 1 eigenstate of X_L
- Measure Z stabilizers



Braiding a CNOT

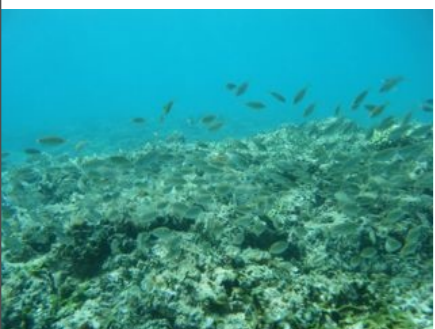
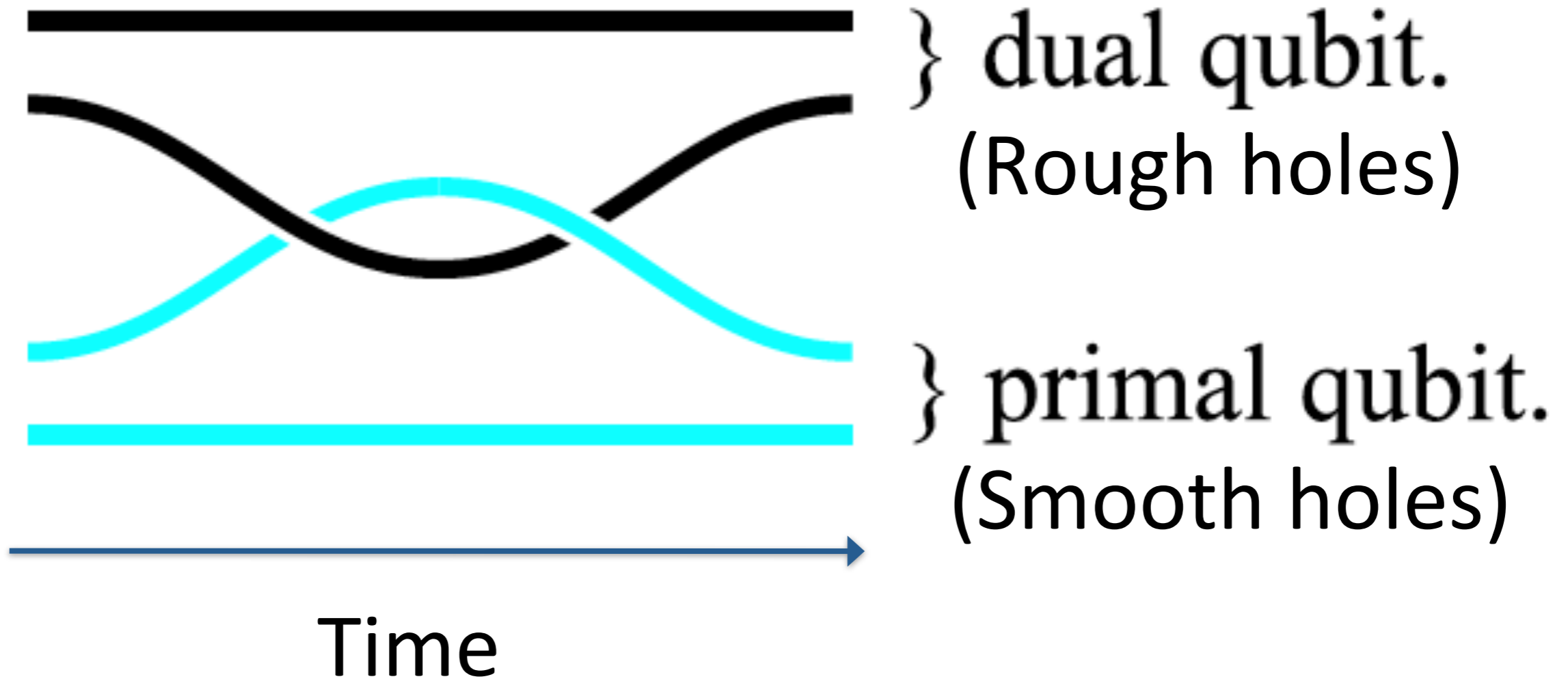
- Braiding smooth pair with rough pair gives CNOT

(WHY? Stay tuned...)





Braiding: 2D+T Picture

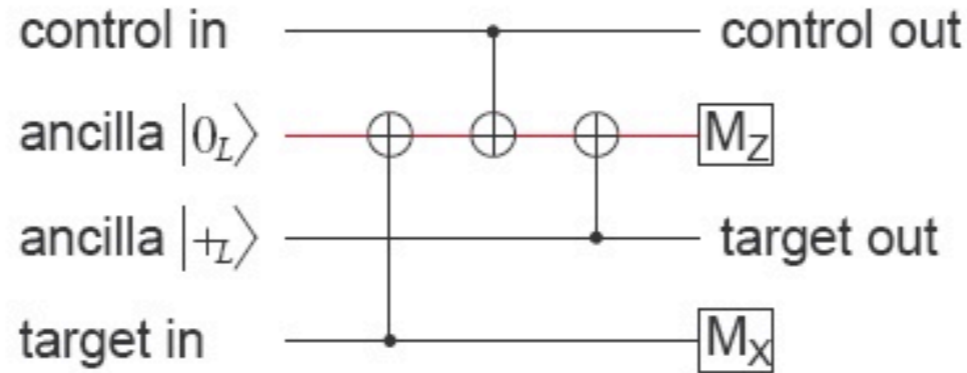


picture from Raussendorf et al.,
NJP 9, 2007

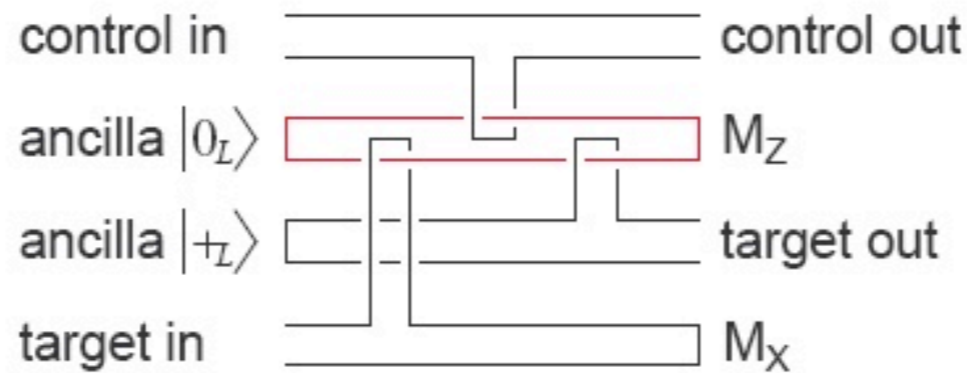


Smooth-smooth CNOT

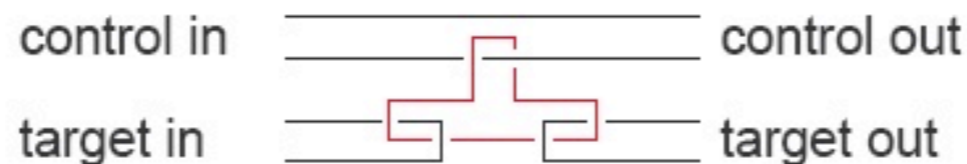
- Circuit is equivalent to CNOT followed by $(Z \otimes Z)^{M_x}$ followed by $(X_t)^{M_z}$



- Circuit can be represented as a braiding of defects

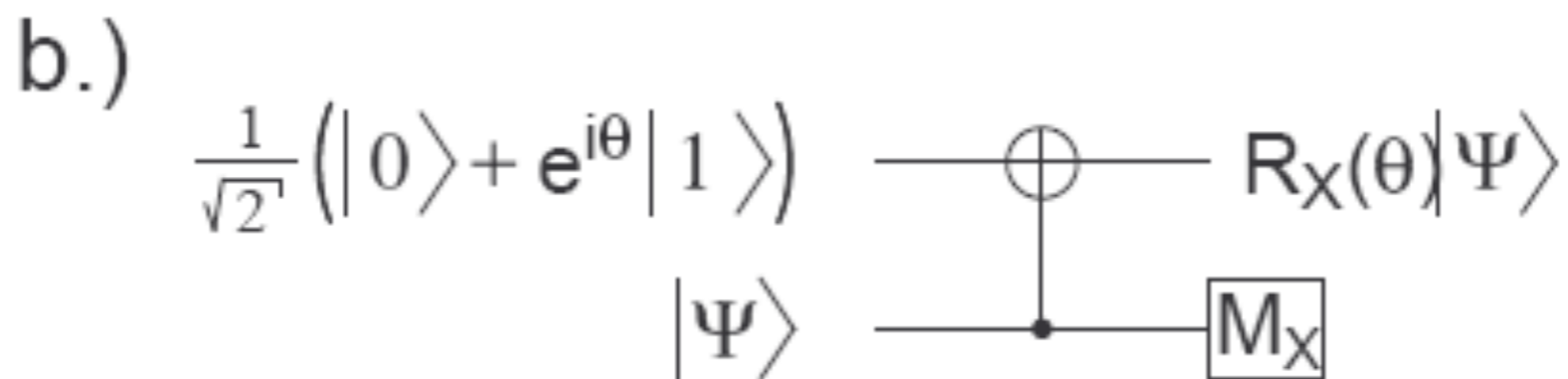
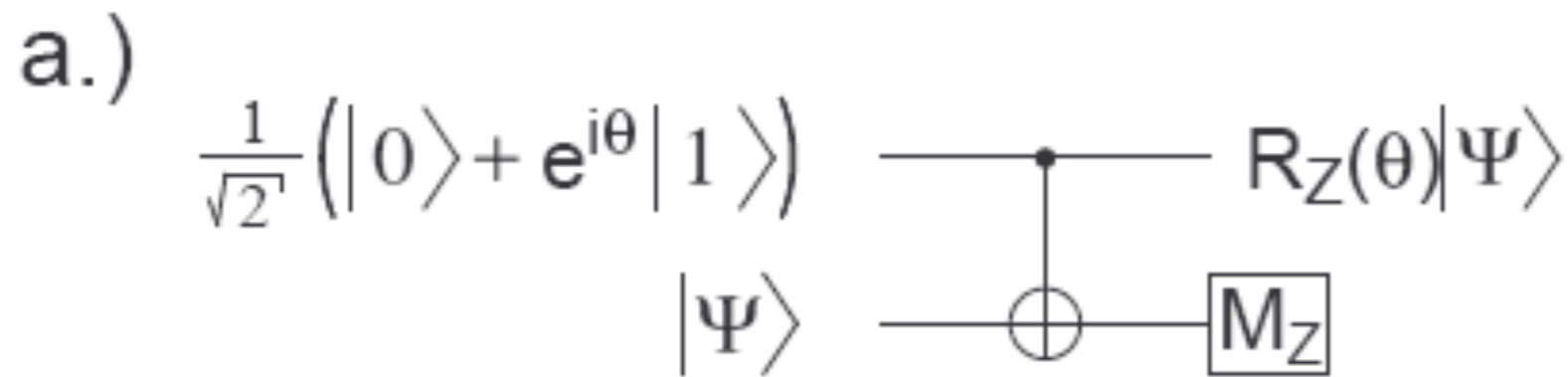


- Equivalent more compact braiding

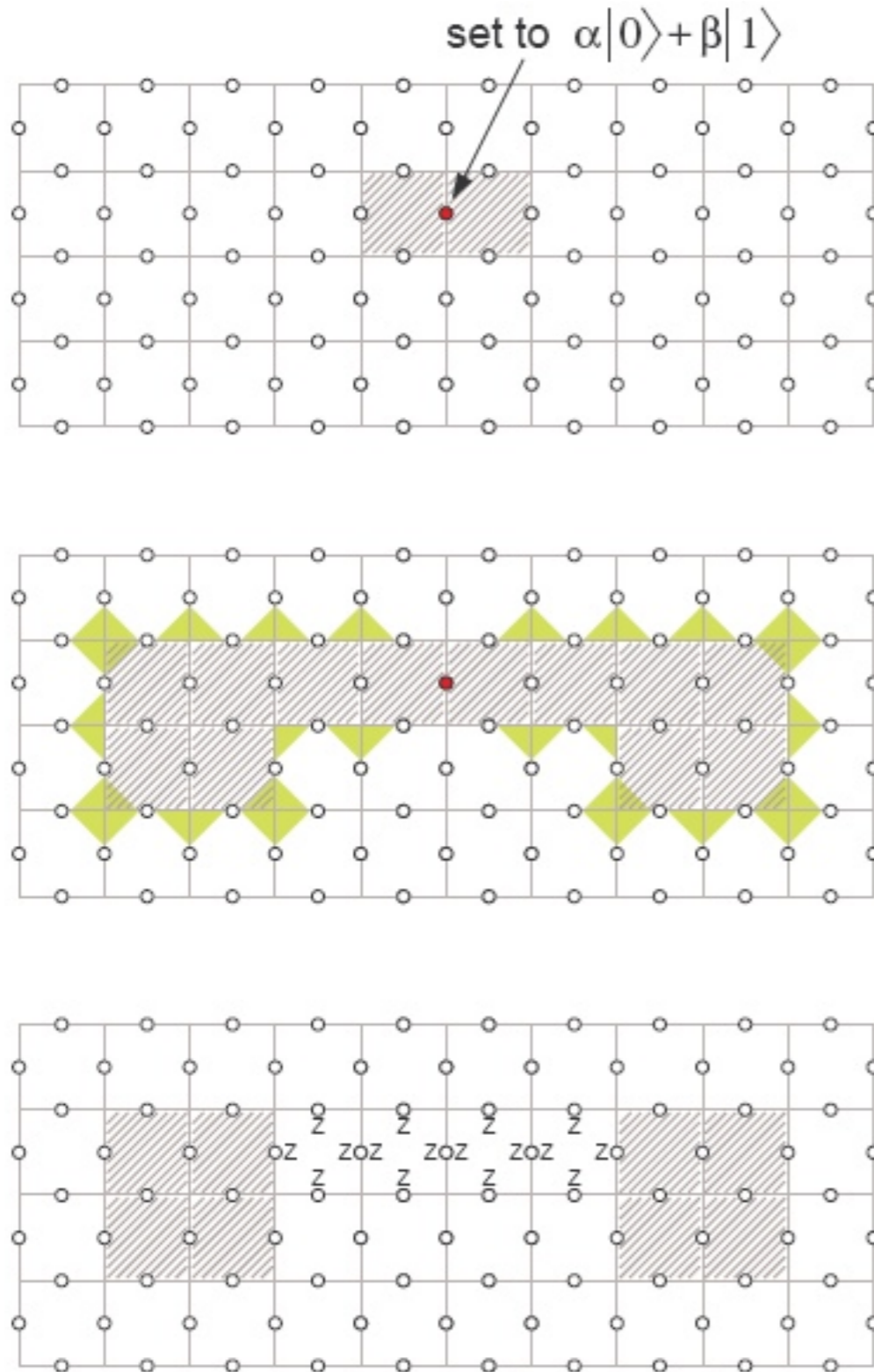




- We can prepare states $|0_L\rangle$, $|1_L\rangle$, $|+_L\rangle$ and $|-_L\rangle$
- We can perform gates X_L , Z_L , M_X , M_Z and CNOT
- Not universal without the following:



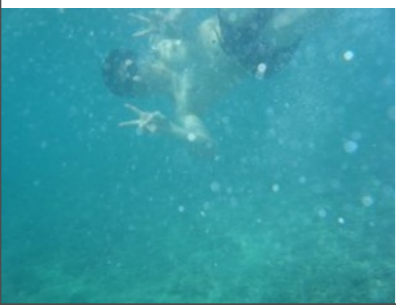
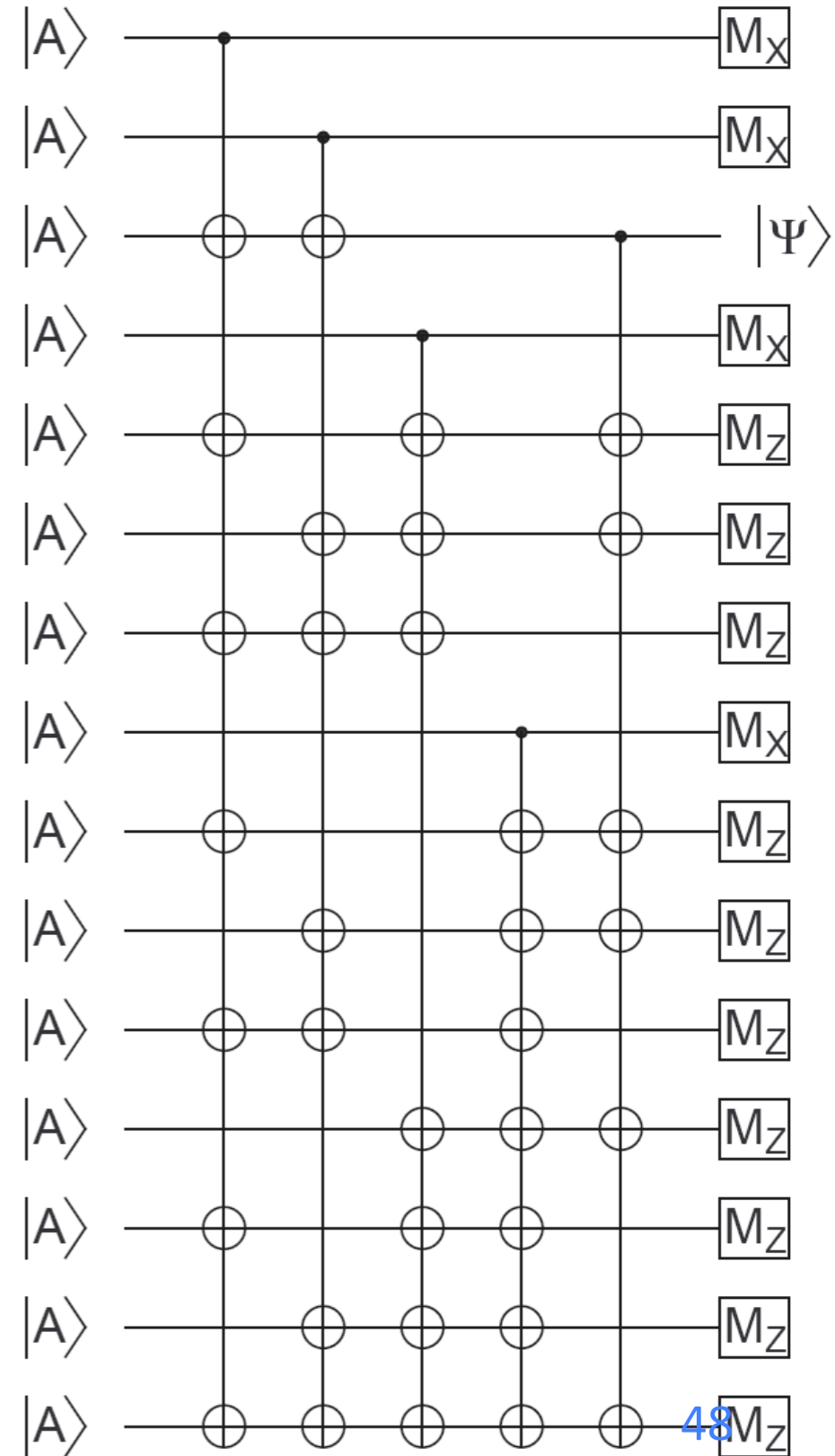
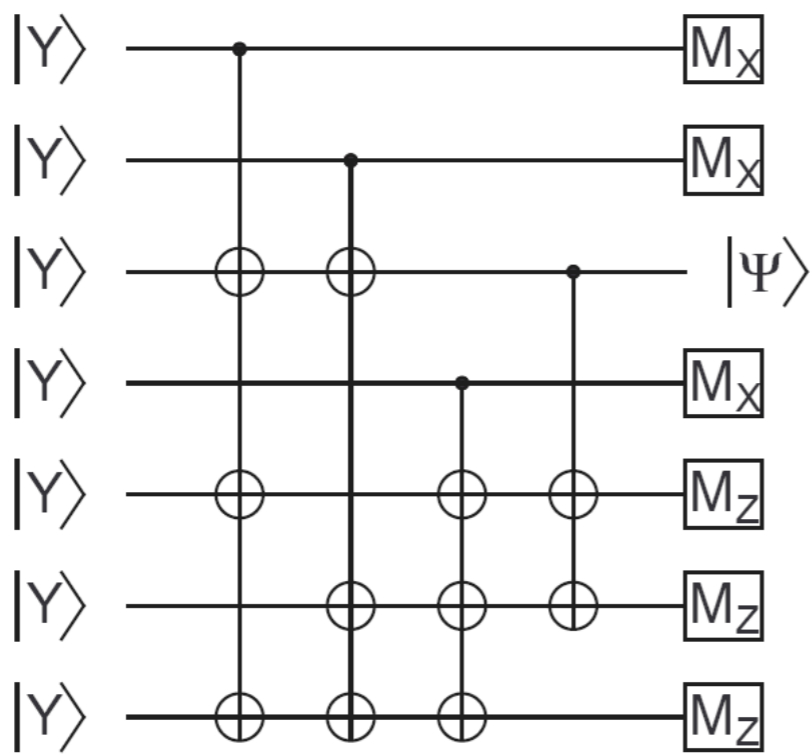
State injection



- Set a single qubit to desired state $\alpha|0\rangle + \beta|1\rangle$ (M_Z then rotate)
- Increase size of defects using X measurements
- Separate defects by measuring and correcting Z stabilizes
- Procedure is not fault-tolerant, restart if errors detected early
- Logical state will not be perfect $\alpha|0\rangle + \beta|1\rangle$, however...

State distillation: "Singular factories"

- Two very special states exist $|Y\rangle = \frac{1}{\sqrt{2}}(|0\rangle + i|1\rangle)$ and $|A\rangle = \frac{1}{\sqrt{2}}(|0\rangle + e^{i\pi/4}|1\rangle)$
- These states can be "distilled"
- Imperfect $|Y\rangle$ and $|A\rangle$ states approach perfect versions exponentially quickly
- Probability of success asymptotically close to 1, though some byproduct operators needed





Universality

- Adding these states:

$$|Y\rangle := (|0\rangle + i|1\rangle) / \sqrt{2}$$

$$|A\rangle := (|0\rangle + e^{i\pi/4}|1\rangle) / \sqrt{2}$$

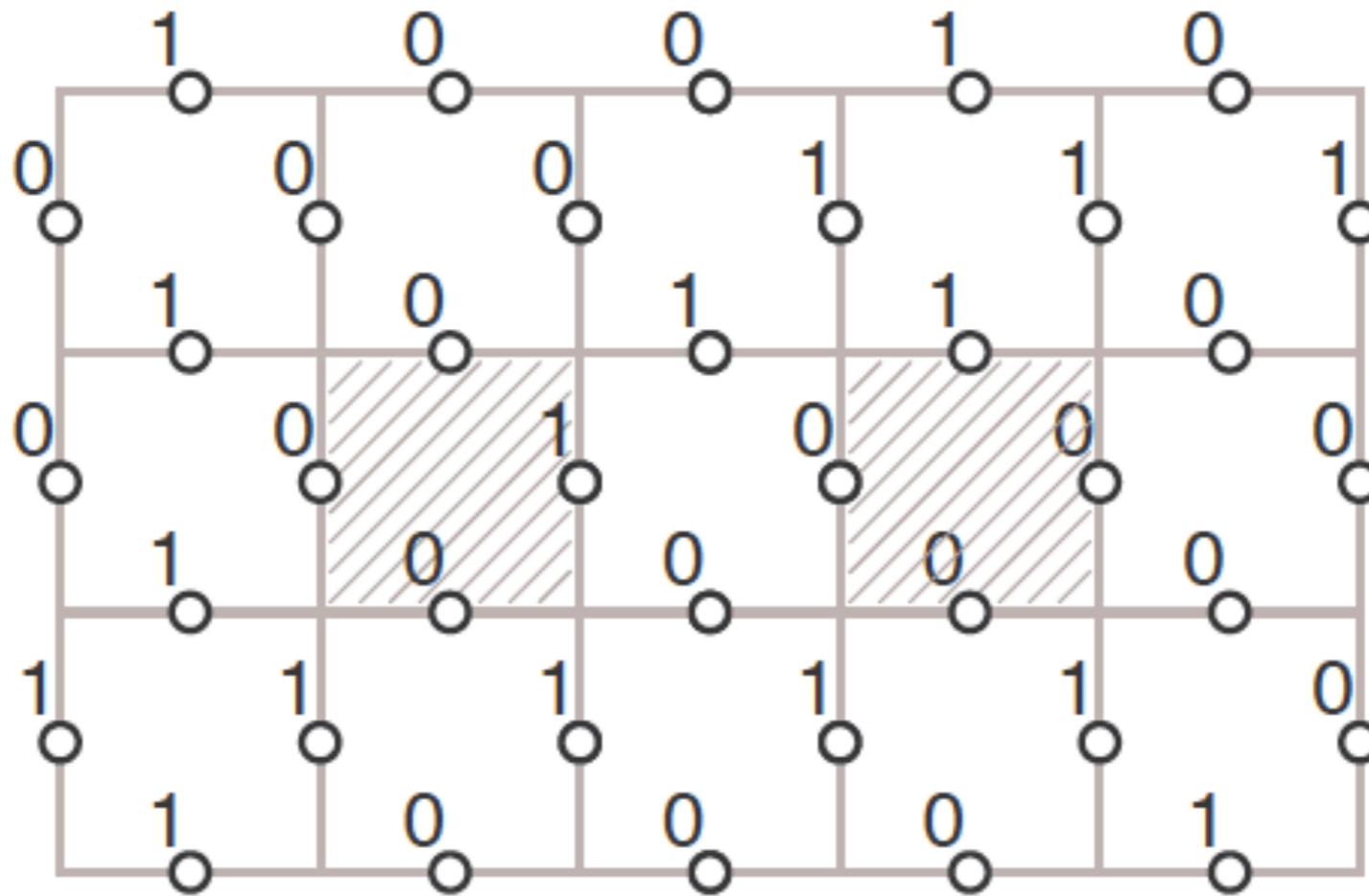
gives us *probabilistic, heralded* universality

- Logical Toffoli gate consumes
7 $|Y\rangle$ and 4.5 $|A\rangle$ states (average)
- Distillation of enough states for *one* Toffoli gate
~1800 singular qubits + braidings
(certain assumptions, not detailed here)





Measuring a Qubit



- Measure region in Z (or X) basis
- Every ring around either defect odd parity, state is 1
- Majority vote when not all same



a : Advancing



51

CALTECH





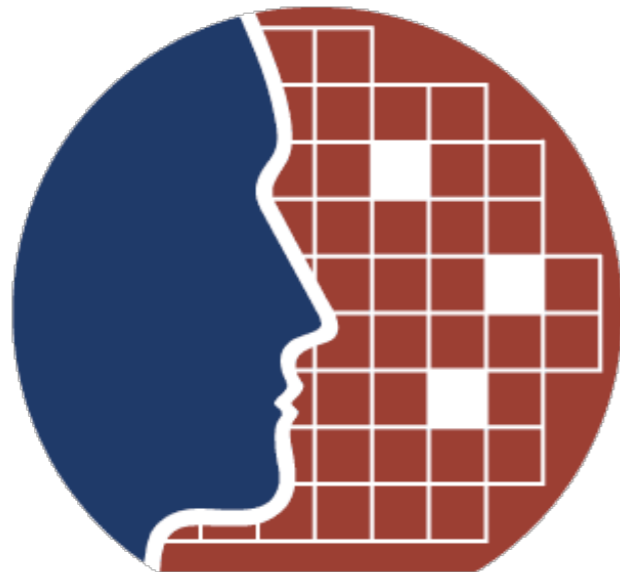
自己紹介



Information Sciences Institute
USC Viterbi School of Engineering

NOKIA

Connecting People



MOSIS

Quantum.

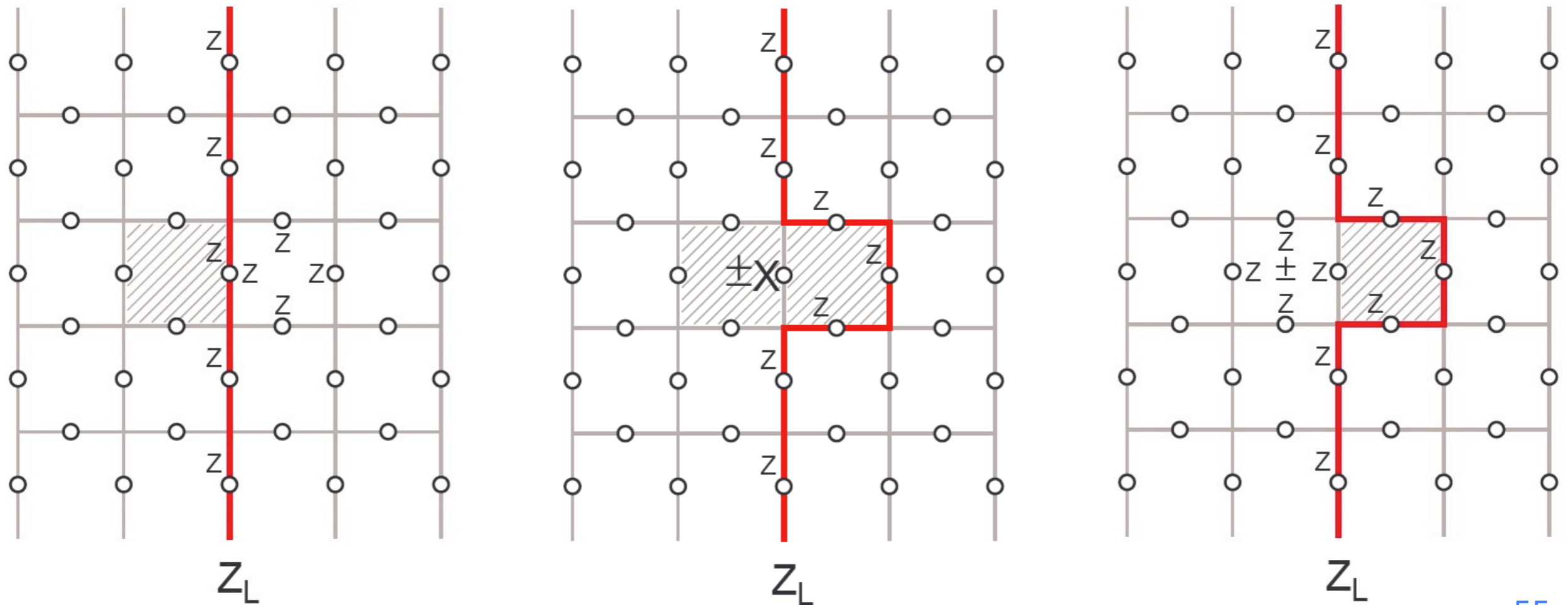
Theory

- Why braiding works
- Code distance
- Estimating logical error rate



CNOT: moving defects

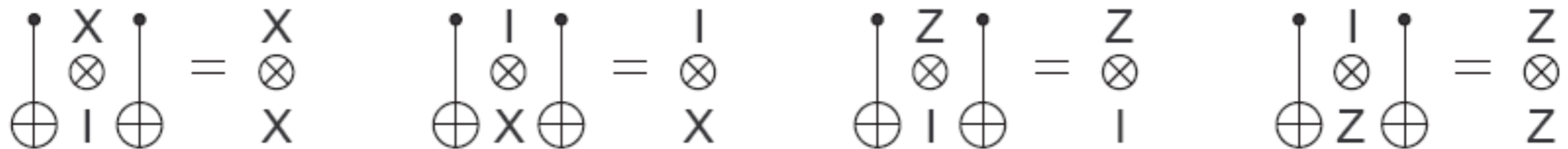
- Start with surface $|\Psi\rangle$ satisfying $Z_L|\Psi\rangle = |\Psi\rangle$ and Z_{face}
- Measure center qubit in the X basis to produce $|\Psi'\rangle$
- Surface $|\Psi'\rangle$ satisfies $Z_L Z_{\text{face}}|\Psi'\rangle = |\Psi'\rangle$ and $\pm X|\Psi'\rangle = |\Psi'\rangle$



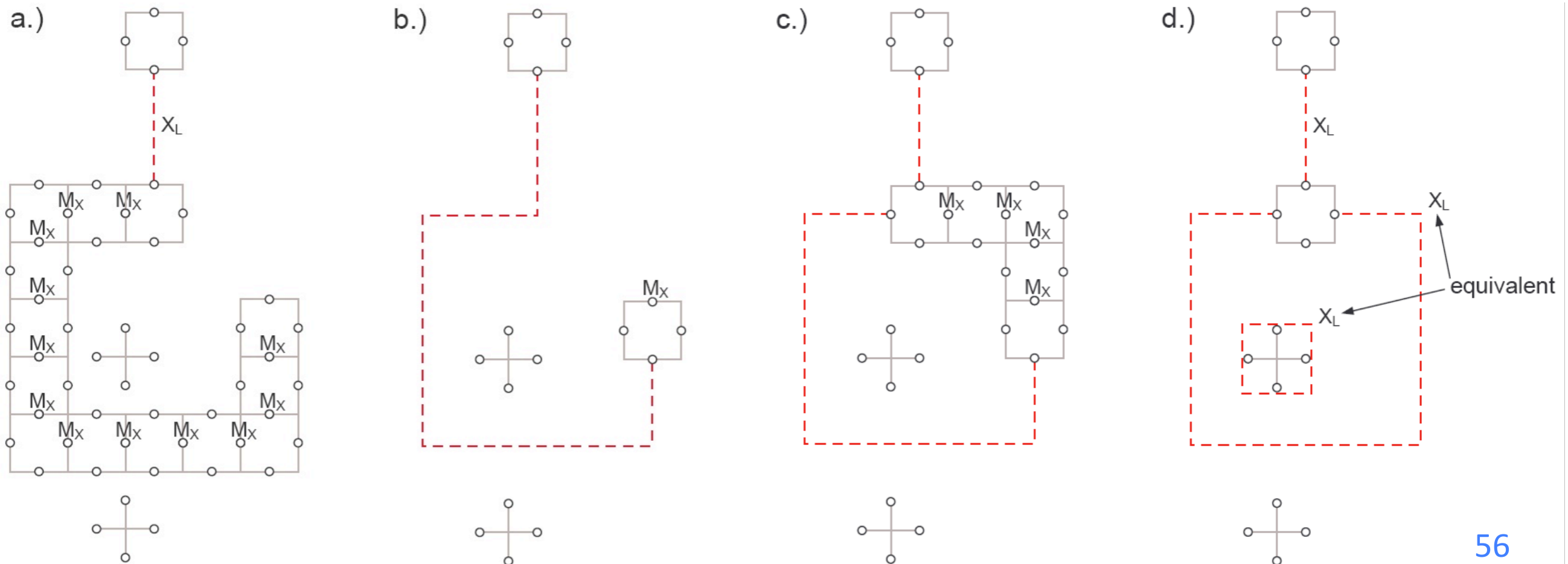
WIDE CNOT: $X \otimes I \rightarrow X \otimes X$



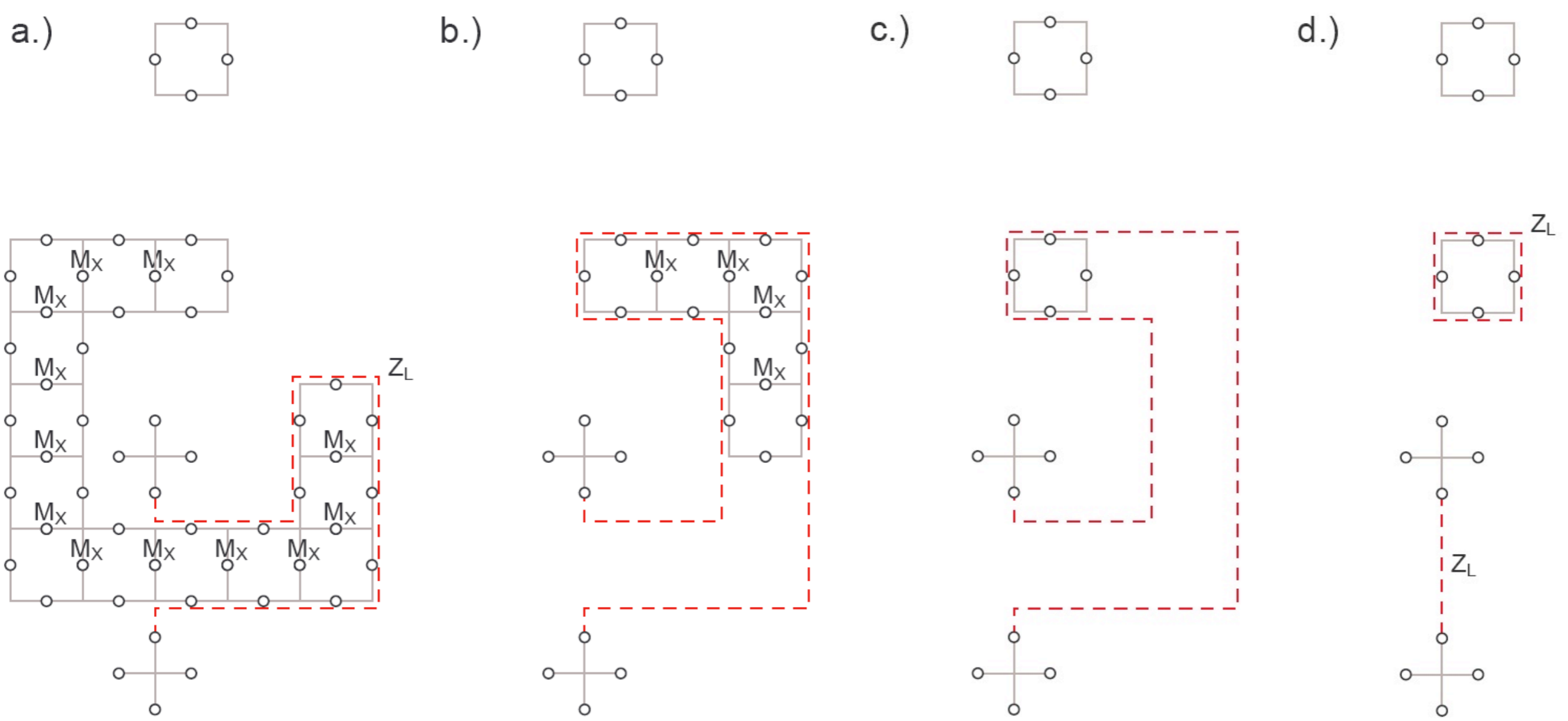
- If $M|\psi\rangle = |\psi\rangle$, then $U|\psi\rangle = UMU^\dagger U|\psi\rangle \Rightarrow M \rightarrow UMU^\dagger$
- CNOT manipulates stabilizers in the following manner:



- Need to show we can do this by braiding defects



CNOT: $I \otimes Z \rightarrow Z \otimes Z$

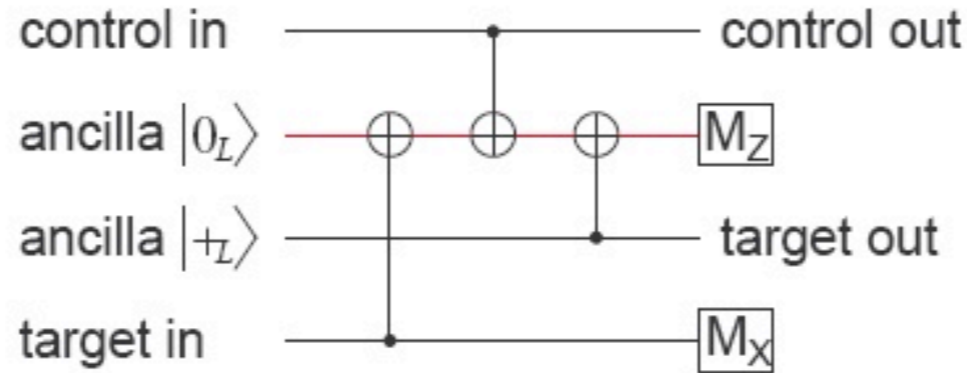


- Mappings $I \otimes X \rightarrow I \otimes X$ and $Z \otimes I \rightarrow Z \otimes I$ also easy to show
- Defects can be interacted over arbitrary distances in almost constant time
- Still need CNOT between two smooth qubits

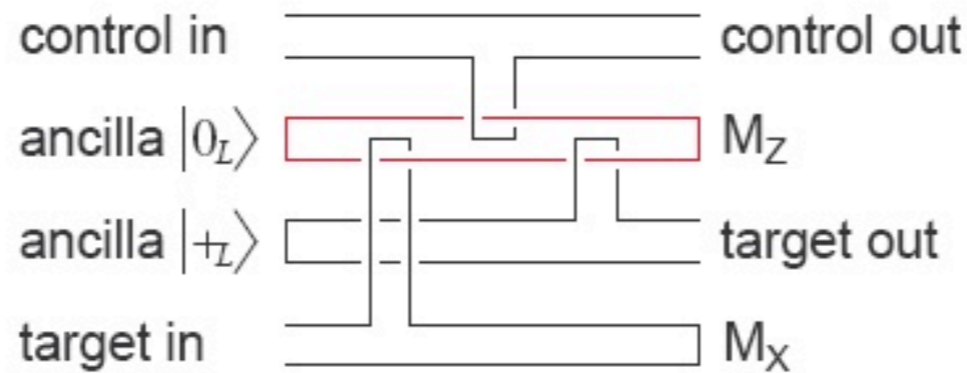


Smooth-smooth CNOT

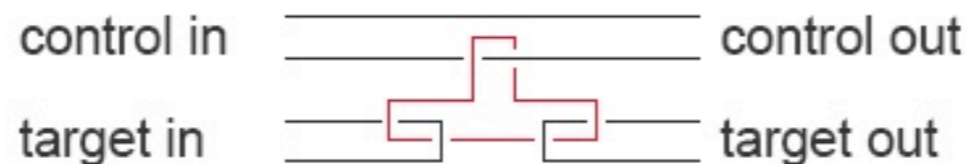
- Circuit is equivalent to CNOT followed by $(Z \otimes Z)^{M_x}$ followed by $(X_t)^{M_z}$



- Circuit can be represented as a braiding of defects



- Equivalent more compact braiding

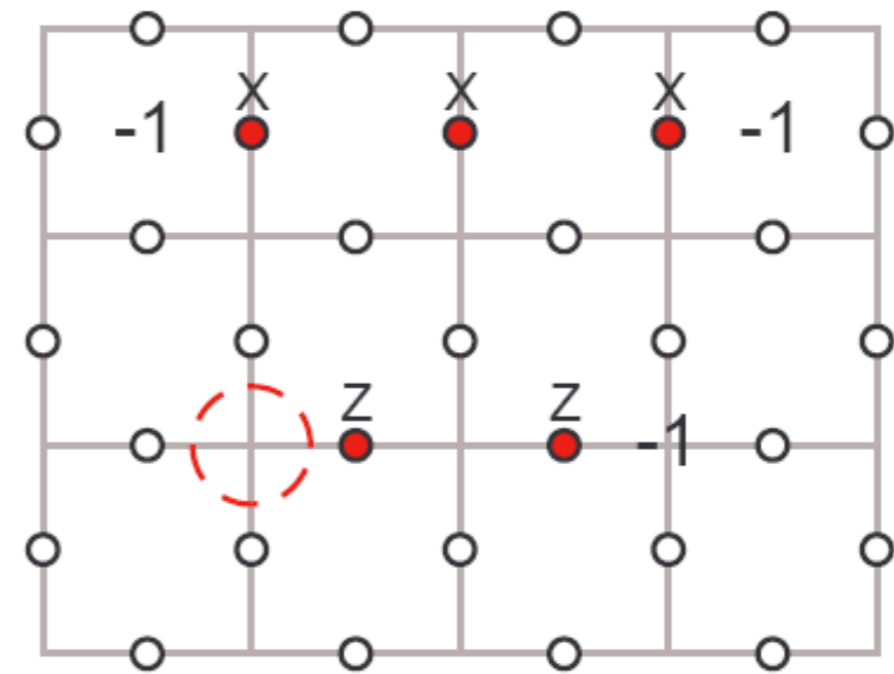
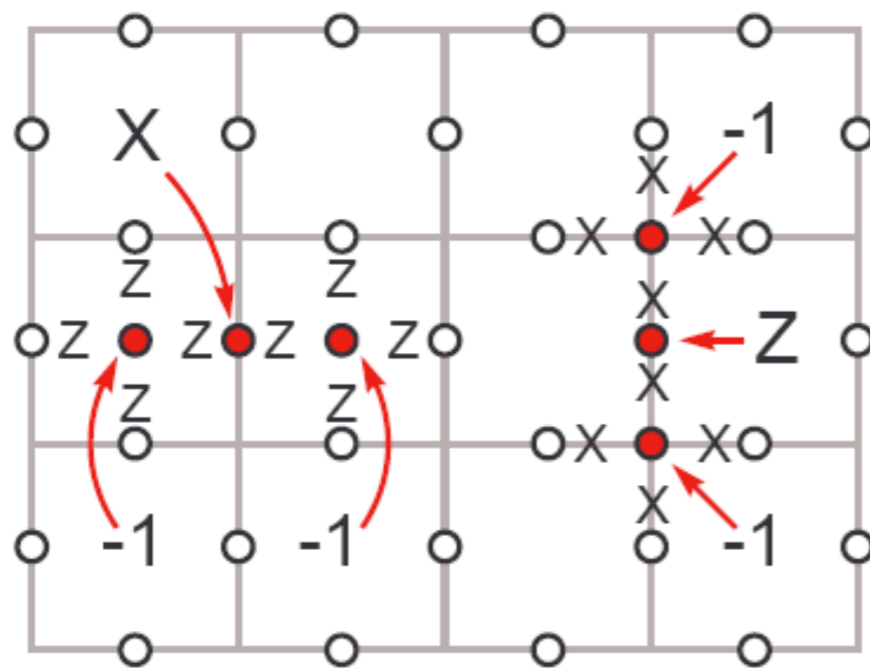


Error correction

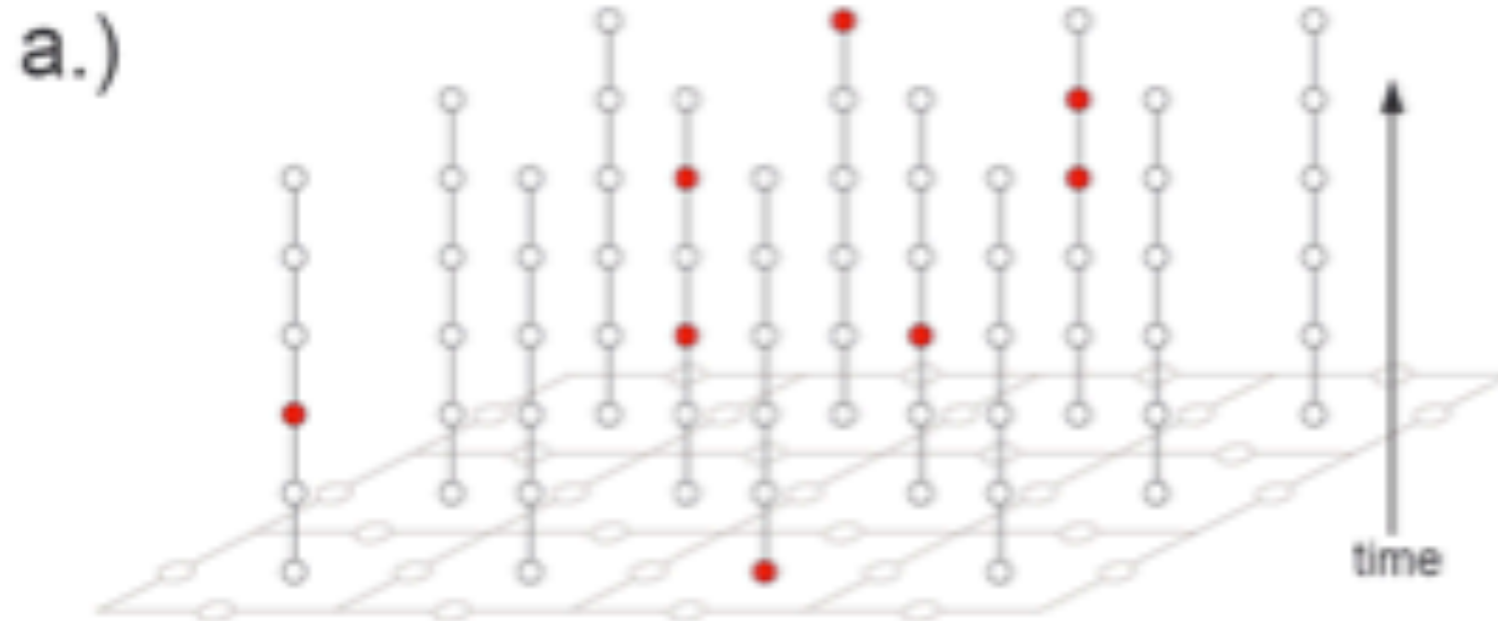
- All error correction based on $XZ = -ZX$, eg:

$$Z_1 Z_2 Z_3 Z_4 |\psi\rangle = |\psi\rangle$$

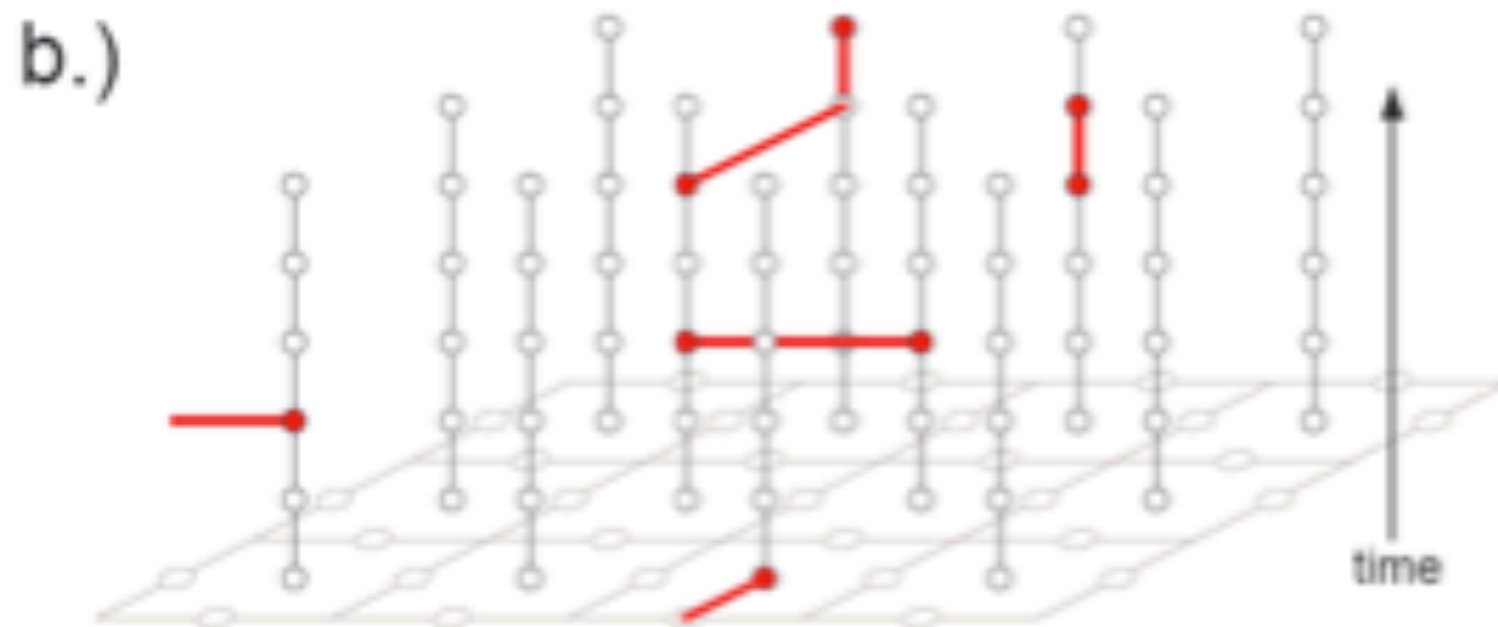
$$Z_1 Z_2 Z_3 Z_4 X_1 |\psi\rangle = -X_1 Z_1 Z_2 Z_3 Z_4 |\psi\rangle = -X_1 |\psi\rangle$$



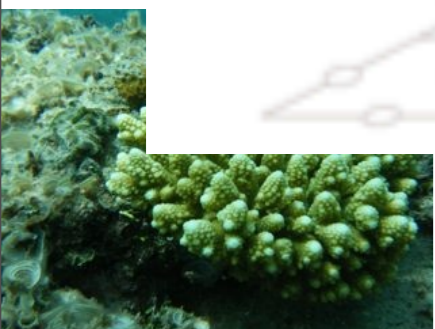
Need to carefully handle false syndromes and error chains



- Record time and position of changed syndromes

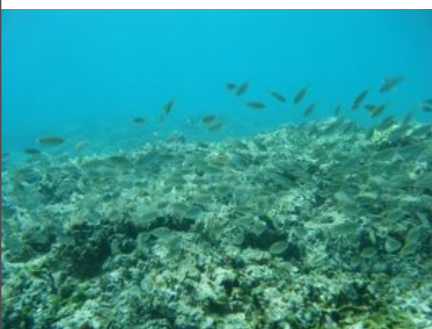
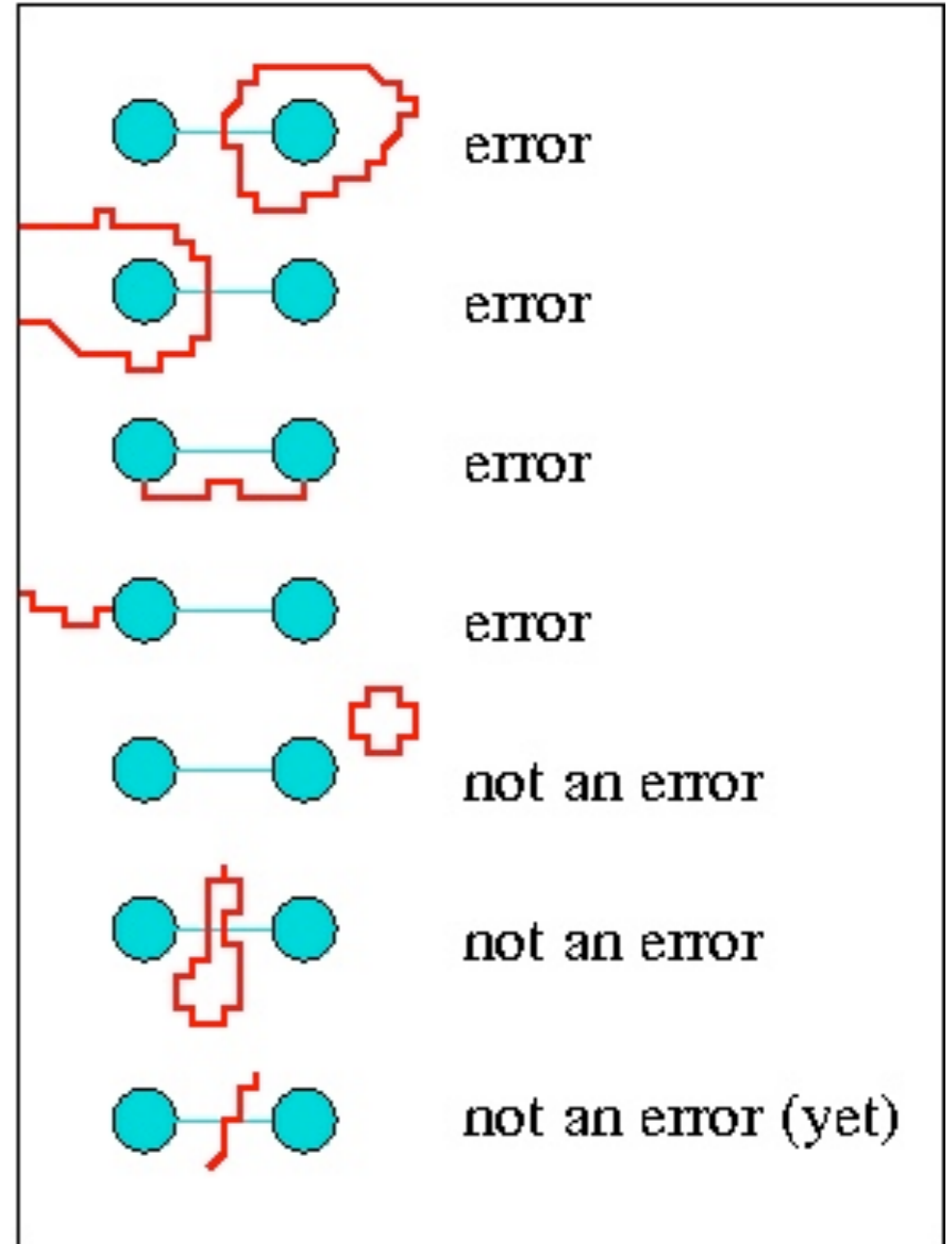


- Match closest pairs
- Apply corrective operations to spacelike edges
- Works *very* well, threshold error rate $\sim 1\%$



Error Chains and Logical Gates

- Learn only endpoints of error chains
- Guess a path
- Usually (hopefully) results in meaningless loop
- Complete error chain (natural or mis-correct) results in logical X or Z gate





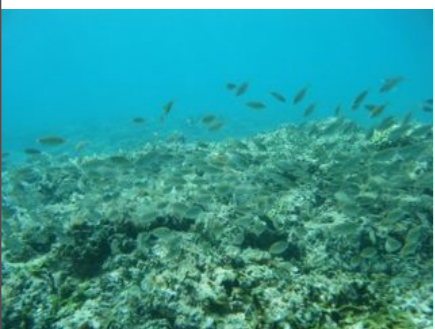
What is error correction?





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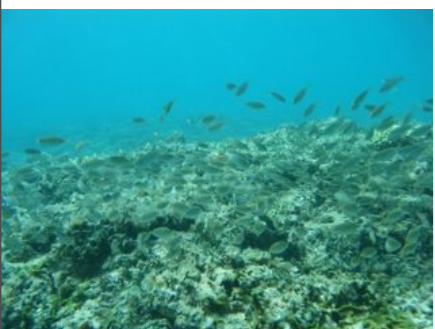
- To protect against errors, encode a few logical states in larger physical Hilbert space





What is error correction?

- To protect against errors, encode a few logical states in larger physical Hilbert space
- Physical errors take the system *out of the code space*
 - → Goal of correction is to get back into code space





What is error correction?

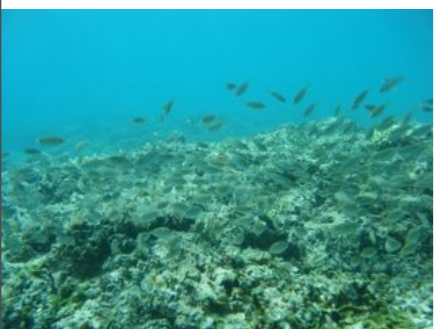
- To protect against errors, encode a few logical states in larger physical Hilbert space
- Physical errors take the system *out of the code space*
 - → Goal of correction is to get back into code space
- Some measure of “distance” from a legitimate code word (logical state)
 - → Usually move to “closest” logical state





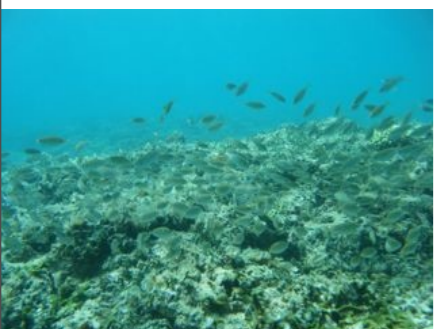
What is error correction?

- To protect against errors, encode a few logical states in larger physical Hilbert space
- Physical errors take the system *out of the code space*
 - → Goal of correction is to get back into code space
- Some measure of “distance” from a legitimate code word (logical state)
 - → Usually move to “closest” logical state
- More than halfway, you mis-correct
 - → Mis-correct equivalent to accidentally executing a logical gate





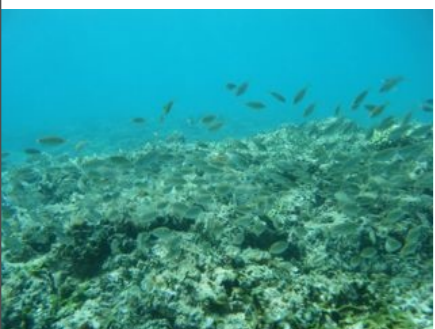
Threshold





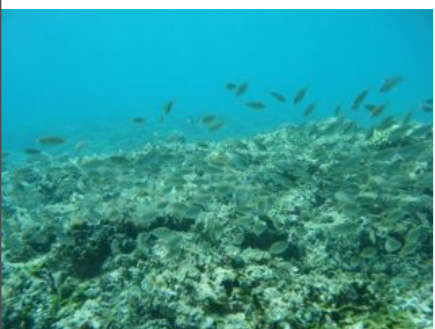
Threshold

- Executing error correction is itself a physical procedure



Threshold

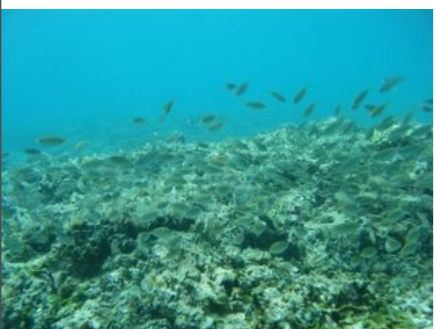
- Executing error correction is itself a physical procedure
- EC can introduce errors





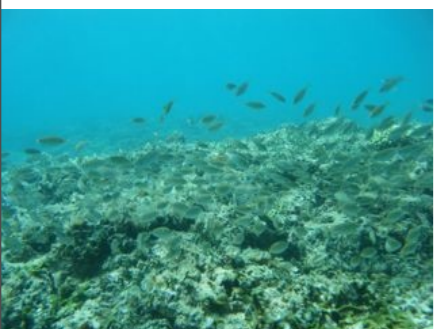
Threshold

- Executing error correction is itself a physical procedure
- EC can introduce errors
- *Threshold* is error level at which probability of logical error *declines* if you apply EC



Threshold

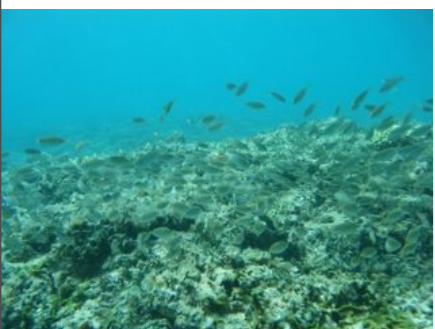
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- EC can introduce errors
- *Threshold* is error level at which probability of logical error *declines* if you apply EC
- Most threshold analyses assume gate, memory, measurement errors same probability





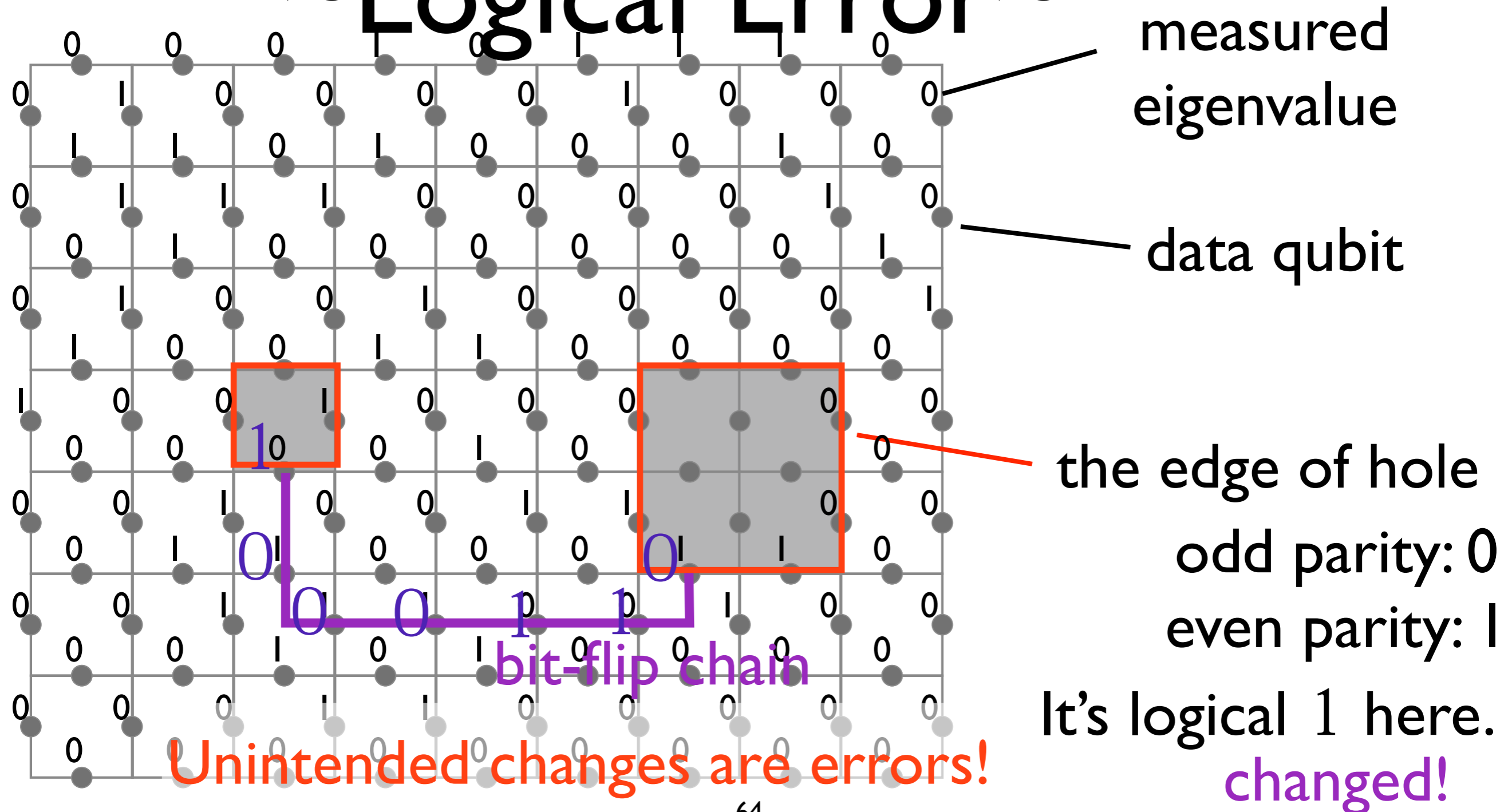
Threshold

- Executing error correction is itself a physical procedure
- EC can introduce errors
- *Threshold* is error level at which probability of logical error *declines* if you apply EC
- Most threshold analyses assume gate, memory, measurement errors same probability
- Generally must *beat* threshold by 1-2 orders of magnitude



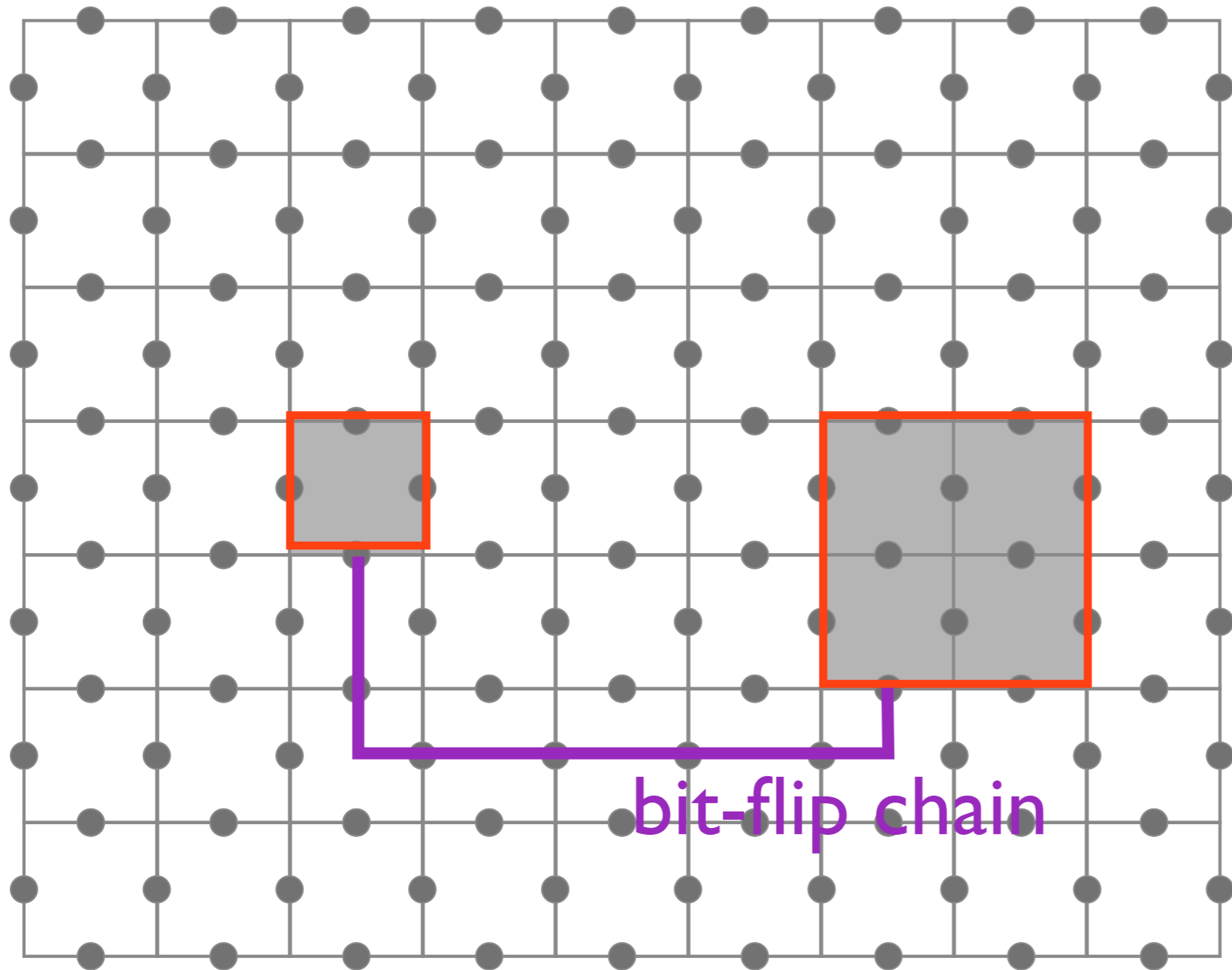
Simplest Example

~ Logical Error ~



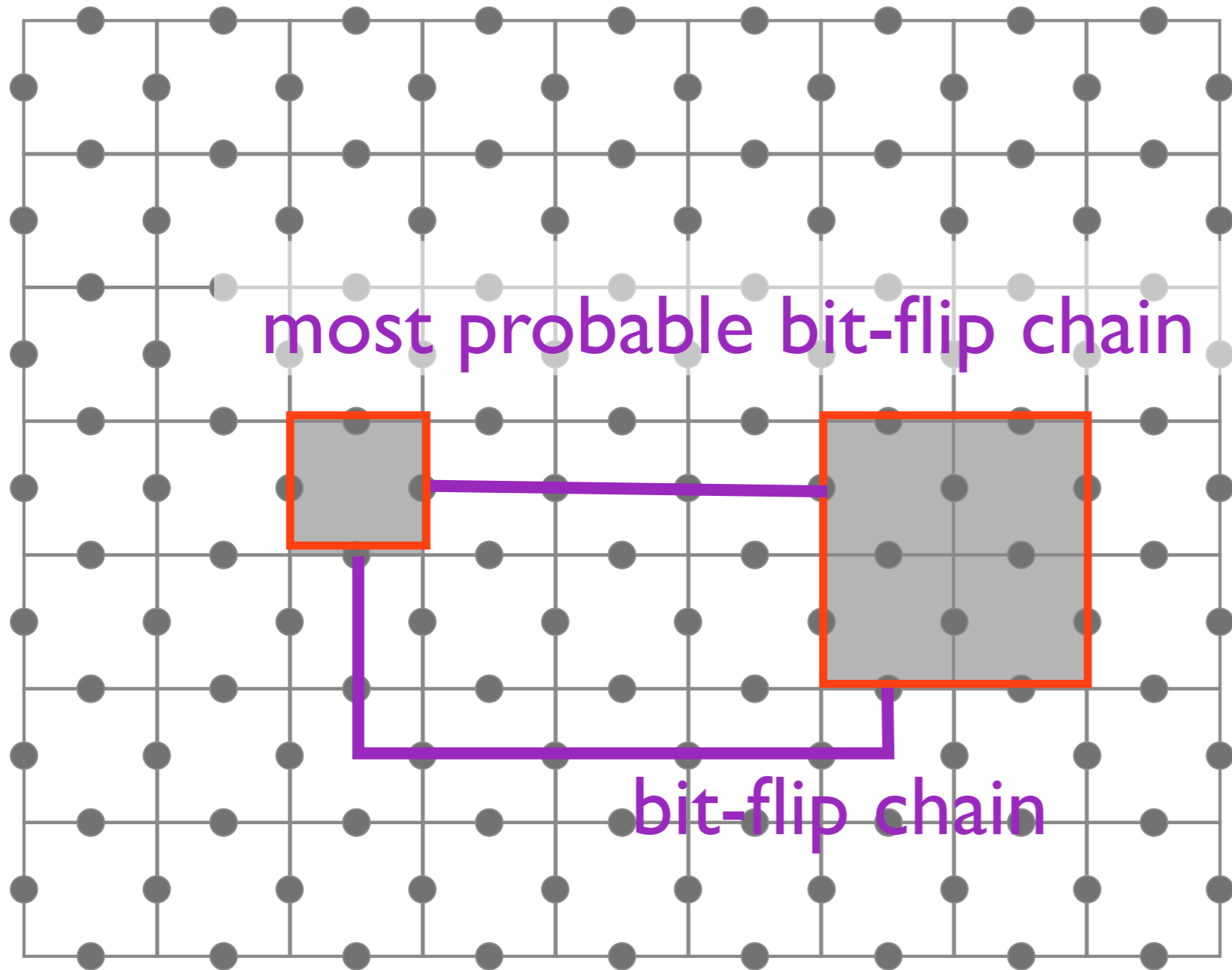
Surface Code

~Code distance~



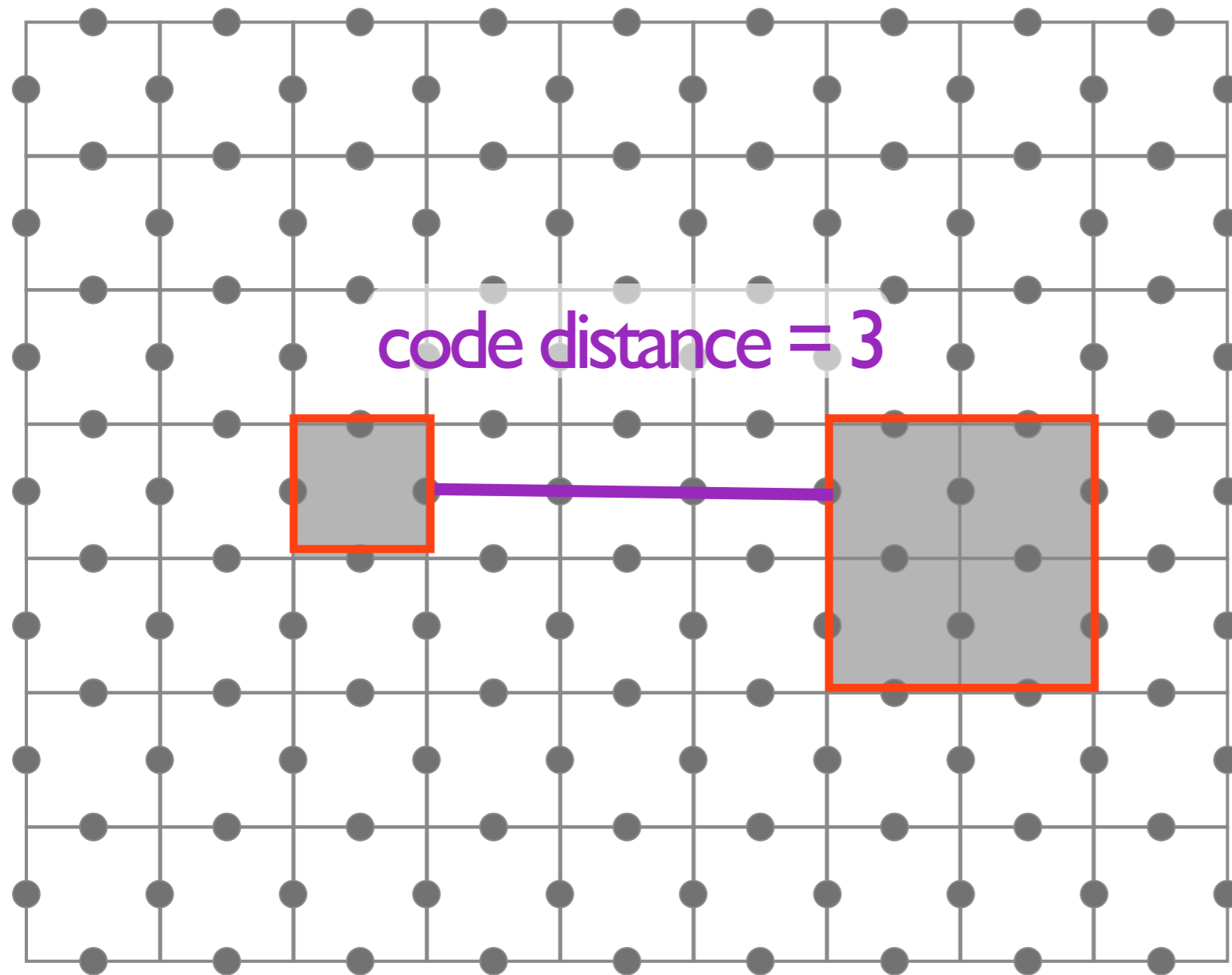
Surface Code

~Code distance~



Surface Code

~Code distance~

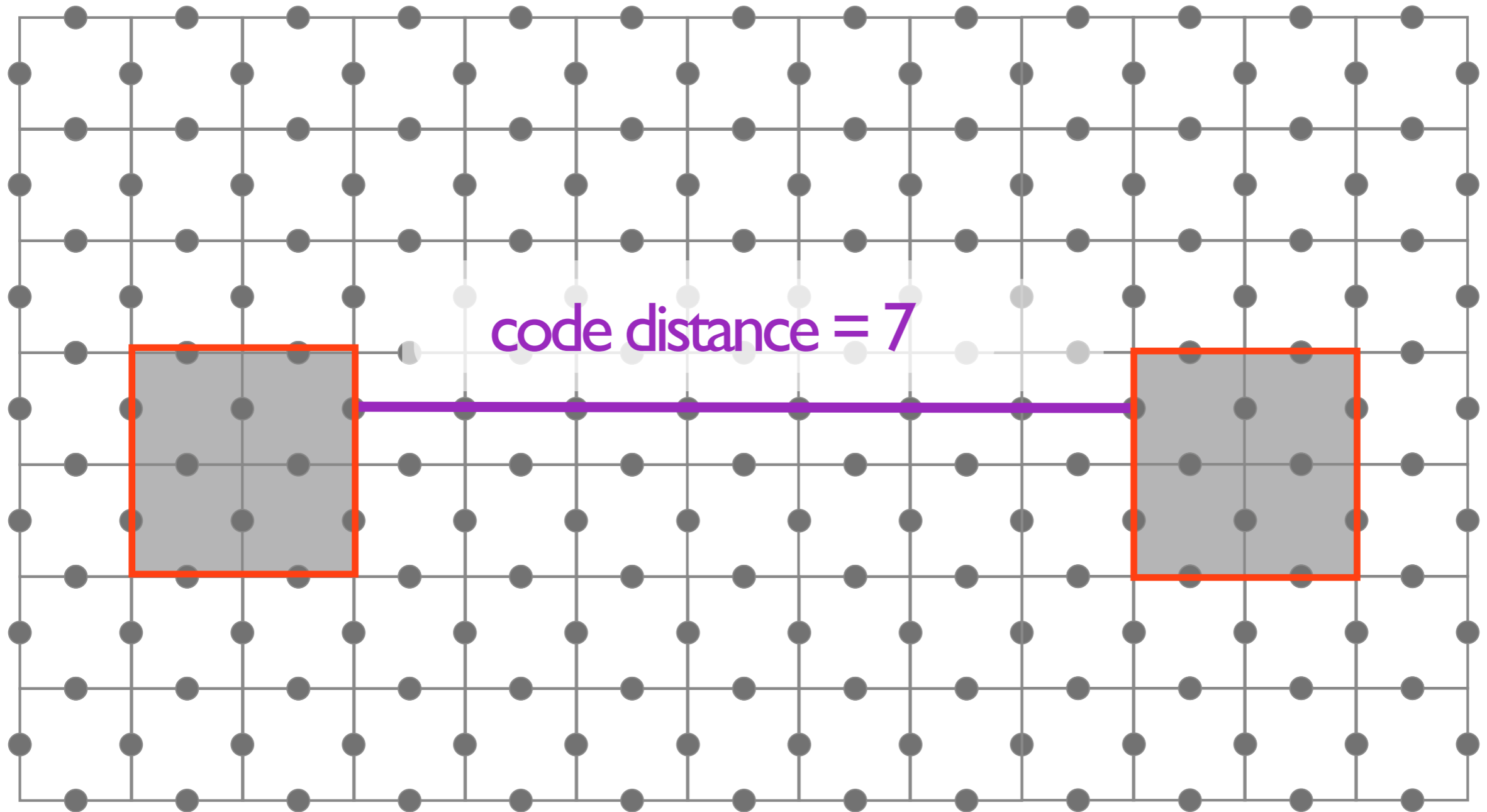


code distance = 3

$$P_{\text{logical memory error rate}} = P_{\text{physical memory error rate}}^4$$

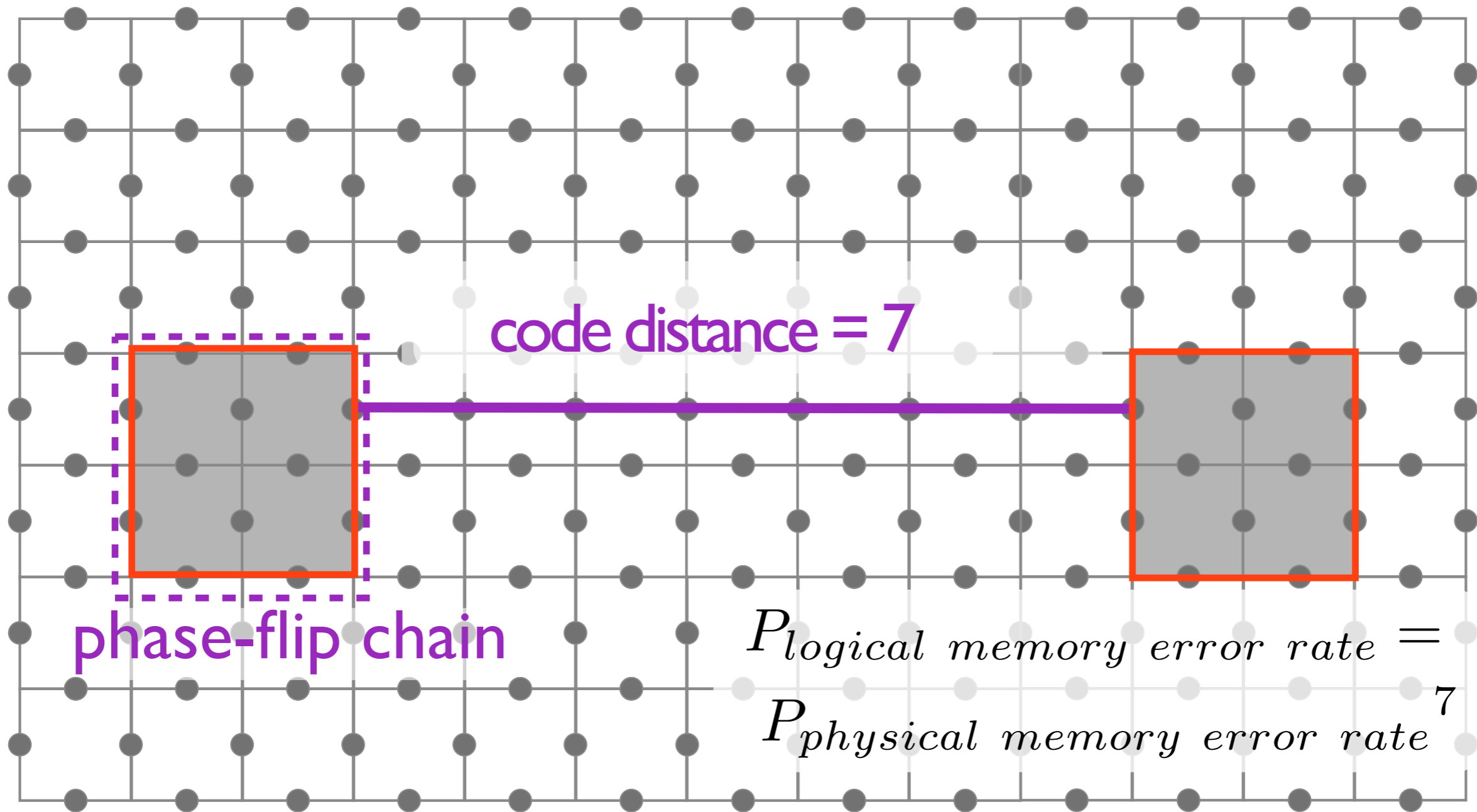
Surface Code

~Code distance~



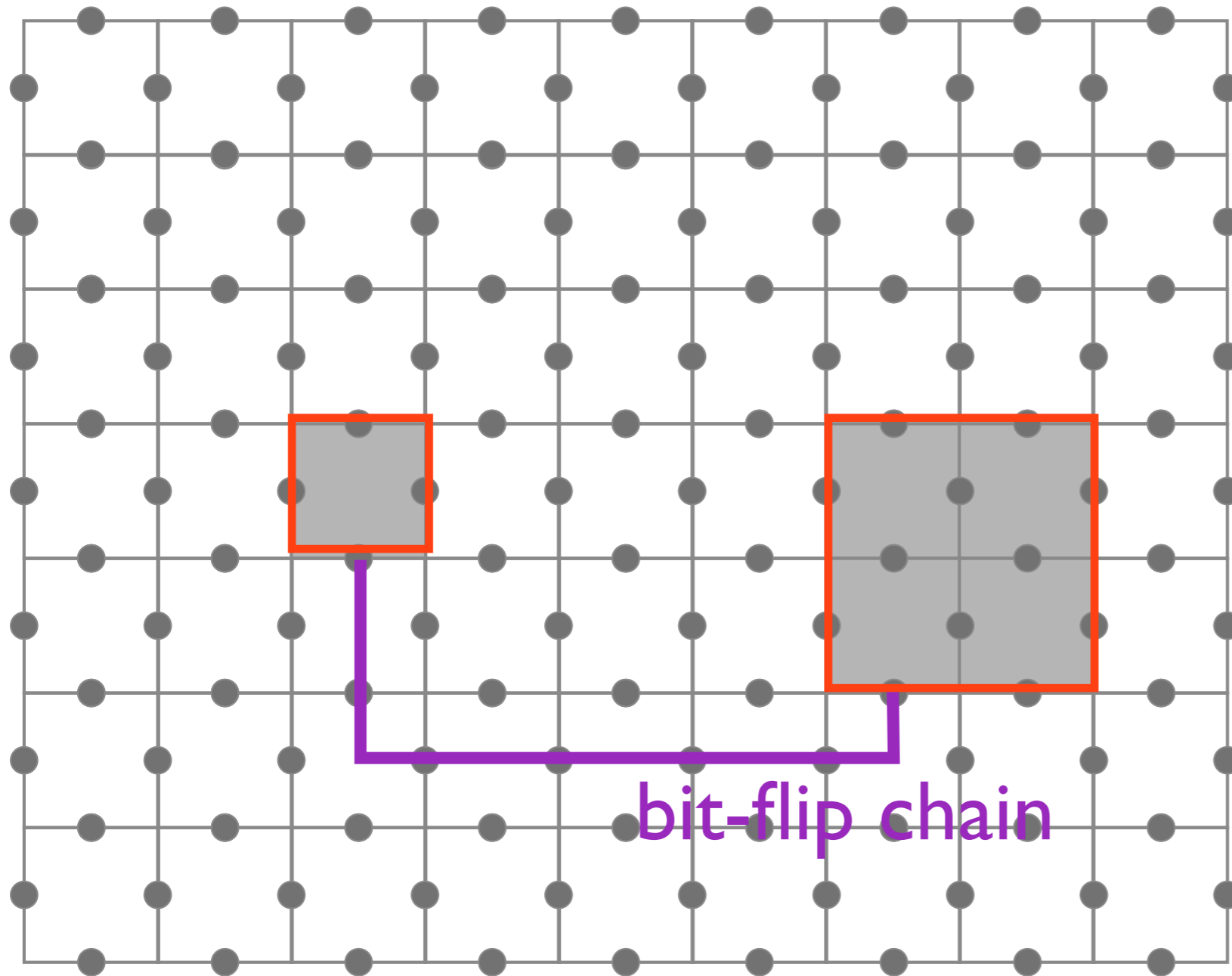
Surface Code

~Code distance~



Surface Code

~Code distance~



$$P_{\text{logical gate error rate}} = P_{\text{physical gate error rate}}^7$$

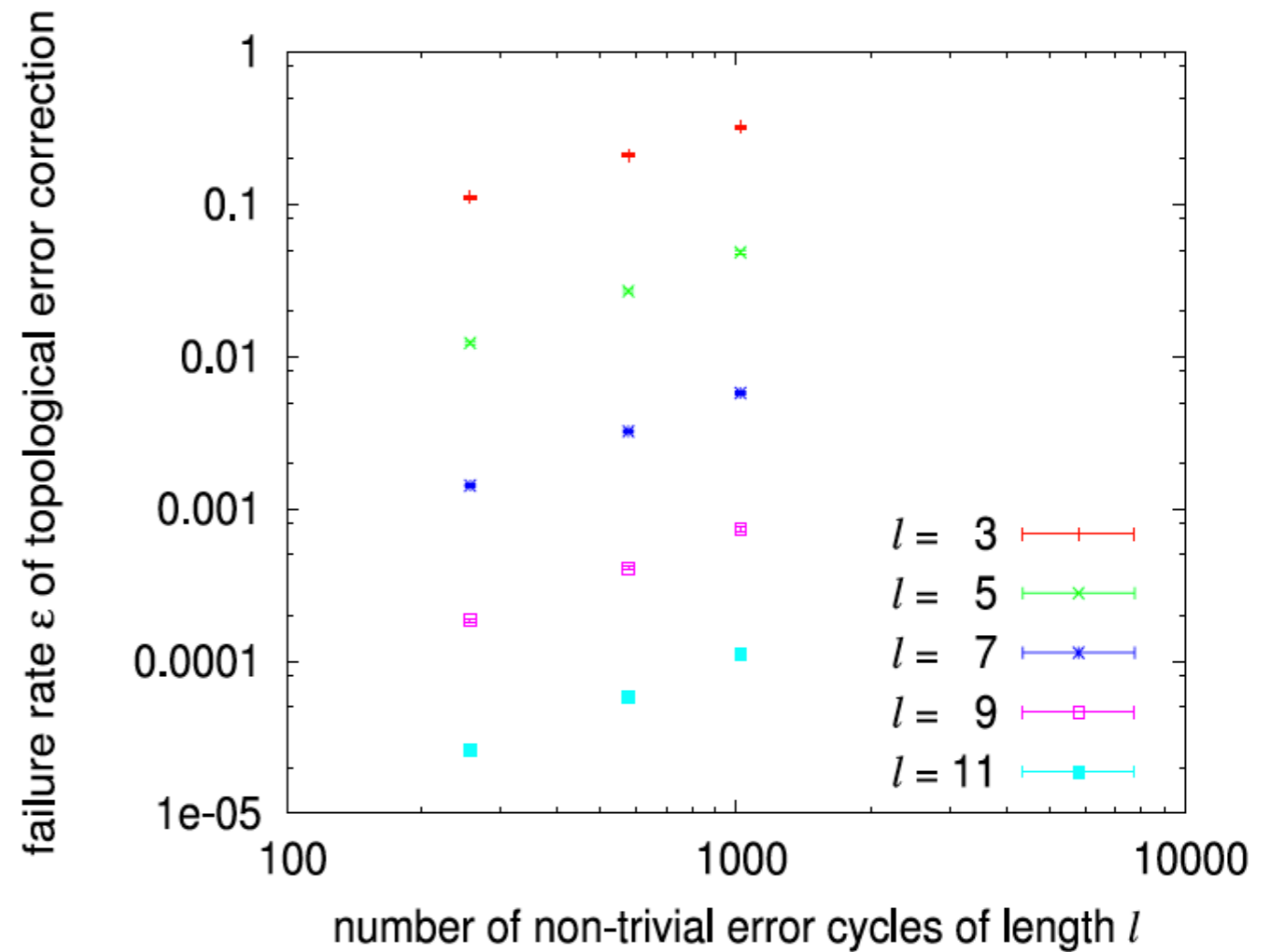


Exponential Suppression

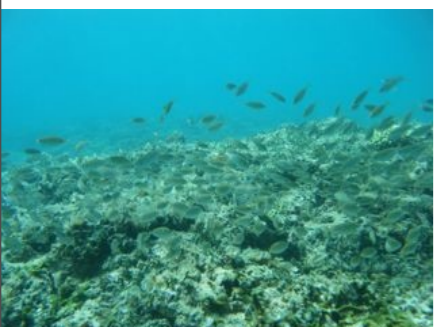
- Classical simulation of small lattices
- Error rate

$$\epsilon_{\text{top}} \sim \exp(-\kappa(p)l)$$

- p = phy err rate
- $K \sim 0.9$
- l = min err chain length

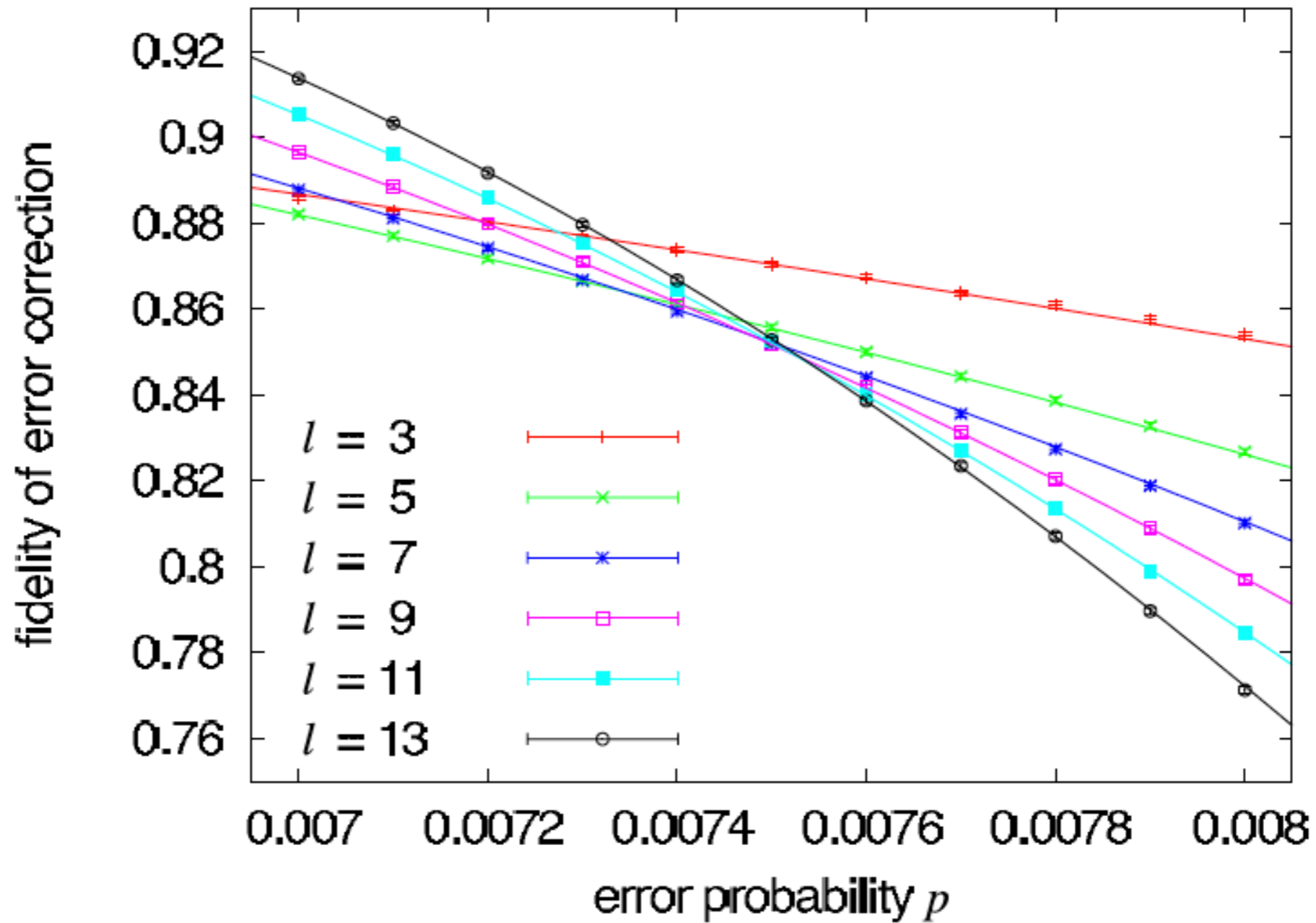


- Still a polynomial correction needed
- Raussendorf et al., NJP 9, 2007

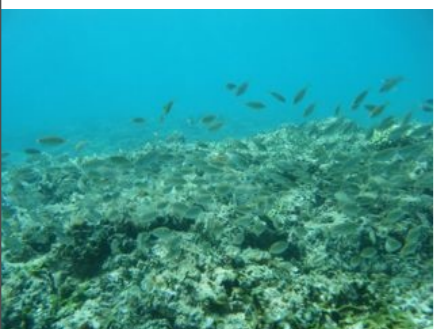




Threshold



- Gate = Memory = Meas errors = 0.75%
- Raussendorf et al., NJP 9, 2007





Advanced Topics

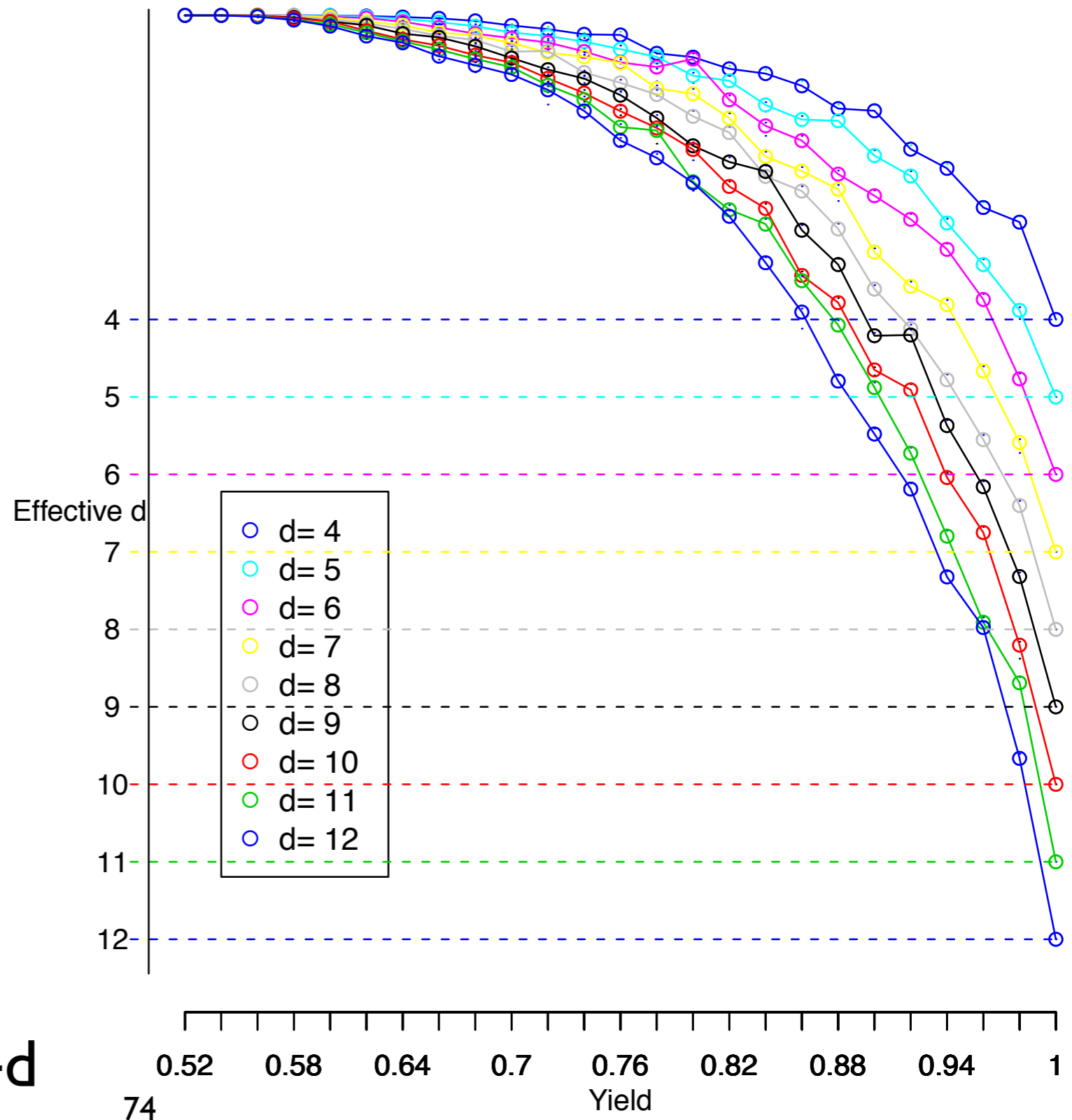
- 3-D version, talk to Simon
- Defective lattice
- Planar code
- Surface code communications
- Uses in distributed quantum computation



Defective lattice

Current estimate is that
yield of 90% *halves* the
effective code distance.

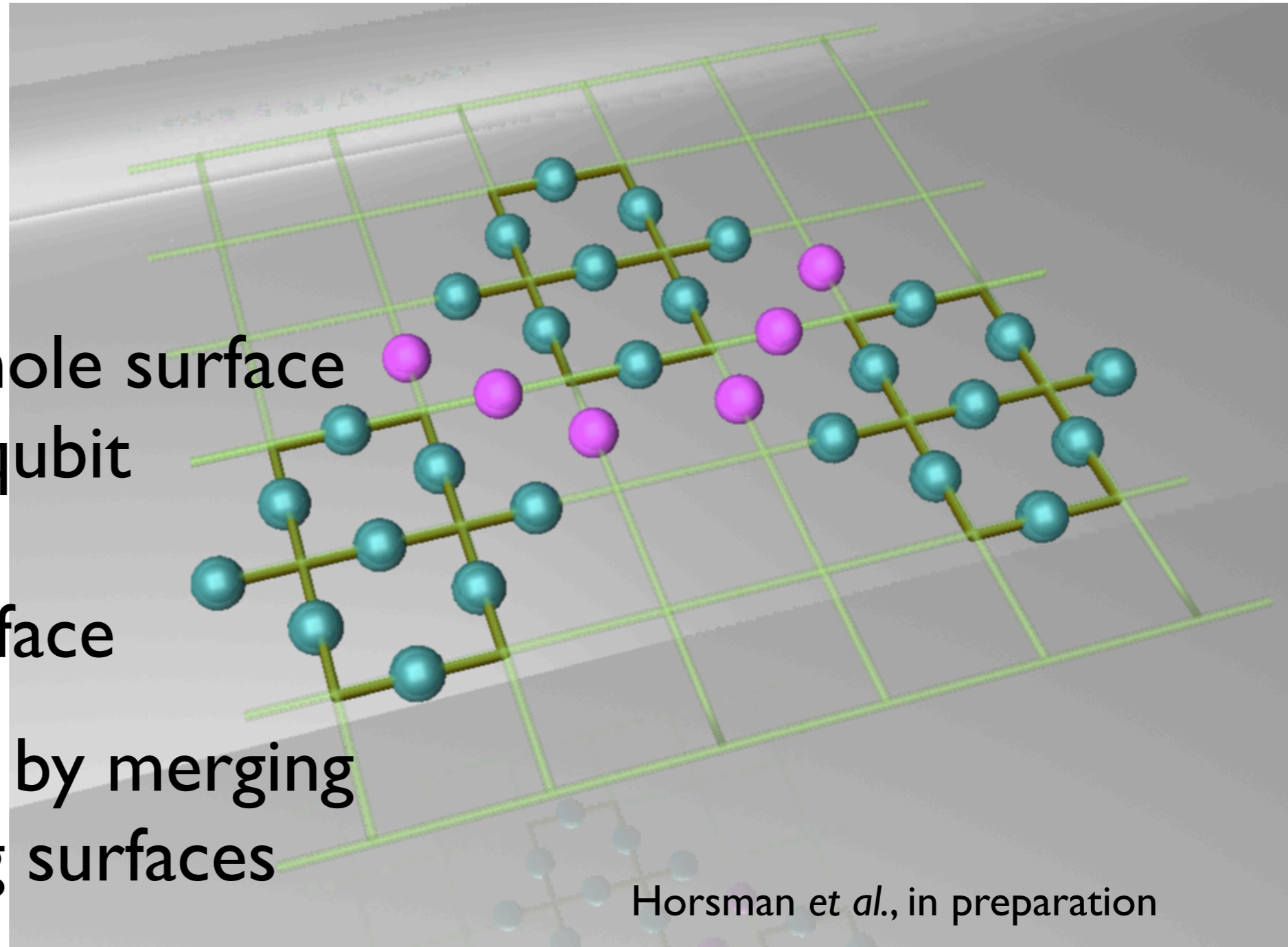
Nagayama *et al.*, in preparation



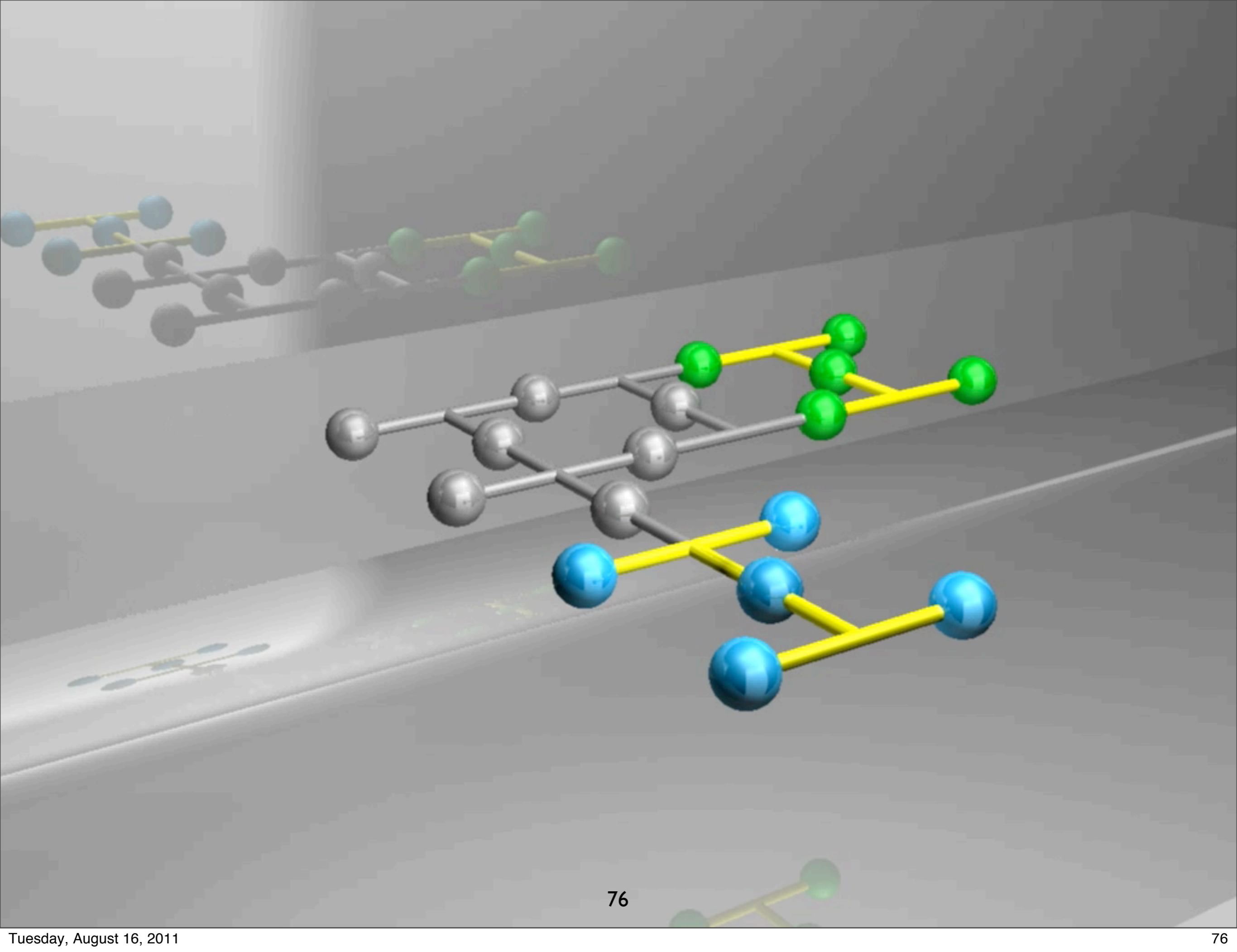
code distance = $4d$

Planar Code

- Use one whole surface per logical qubit instead of holes in surface
- Gates done by merging and splitting surfaces
- Useful for small-scale experiments



Horsman *et al.*, in preparation





Surface Code Strengths





Surface Code Strengths

- Simple, 2-D or 3-D nearest-neighbor-only operation (physical feasibility high!)





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 - strength of EC grows incrementally (compare to concatenated CSS codes)
 - software-assigned resources
 - easy movement of logical qubits





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- High threshold: 1.4% for gate, memory, measurement errors
- Flexible:
 - strength of EC grows incrementally (compare to concatenated CSS codes)
 - software-assigned resources
 - easy movement of logical qubits
- Supporting classical processing achievable



Key References

- clearest explanation:
Fowler *et al.*, PRA 80, 052312 (2009)
(source of many of the figures)
- detailed paper:
Raussendorf *et al.*, NJP 9, 199 (2007)
- cryptic seminal paper:
Raussendorf and Harrington, PRL 98, 190504 (2007)
- Surface code communication:
Fowler *et al.*, PRL 104, 180503 (2010)
- Defects in the surface code:
Stace *et al.*, PRL 102, 200501 (2009)



8 Hours of Lecture on Surface Code



Surface code 量子誤り訂正に関するチュートリアル・ワークショップ



お知らせ

Surface code 量子誤り訂正に関するチュートリアル・ワークショップ
 FIRST/Quantum Cybernetics/CREST Joint 1.5-day Surface Code Quantum Error Correction Tutorial/Workshop

日時 2011年2月23日 (水) 10:00-17:00, 24日 (木) 10:00-12:00
 場所 大阪大学豊中キャンパス・基礎工学研究科
 担当者 Rodney Van Meter (慶應義塾大学)、北川勝浩 (大阪大学)
 講師 永山翔太、Rodney Van Meter、Clare Horsman (慶應義塾大学)
[FIRST 最先端研究開発支援プログラム量子情報処理プロジェクト](#)

授業マテリアル

受講したい回をクリックしてください。

第 01 回 2011/02/23 Day One (1)

- 📄 [Lecture Material 1 \(pdf\)](#) (3096445バイト, 5/4/2011)
- 📄 [Lecture Material 2 \(pdf\)](#) (1670839バイト, 5/4/2011)
- Introductions/Plan for the two days [Rodney Van Meter]
- Basic surface code concepts [Shota Nagayama]
- The lattice and cluster state, stabilizers, qubit state

📺 [Start Video](#)

第 02 回 2011/02/23 Day One (2)

- 📄 [Lecture Material \(pdf\)](#) (2416589バイト, 5/4/2011)
- Intermediate topics [Rodney Van Meter]
- Lattice operation



• <http://aqua.sfc.wide.ad.jp/Publications.html>

Preview: What kind of system can run surface code effectively?

- Billions of qubits
- GHz physical gates
- Millisecond memory lifetimes
- Error rate $\sim 0.1\%$
- ~ 1 year to factor 2048-bit number
- Stay tuned...