# Computational approaches to analyze complex dynamic systems: model-checking and its applications.

Part 4: Models and algorithms to analyze large-scale concurrent systems: approaches inspired by pi-calculus and static analysis

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## Introduction

### Modeling biological regulatory networks

- Thomas' framework
- From Thomas' framework to discrete-event systems
- From Thomas' framework to timed systems
- Common limits of current models for biological analyses

3 The Process Hitting: a framework well suited to concurrent systems

- Definition
- From biological models to Process Hitting and refining
- Tool for analyzing Process Hitting: pint
- Inferring information on the biological model thanks to the Process Hitting
  - Interaction Graph Inference
  - Parametrization Inference
- 5 Summary & Conclusion

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# Motivations

## Objective: formal verification of properties

- Model the system S :
  - Discrete models: finite state automata, Petri nets,  $\ldots \Rightarrow$  Lecture 1
  - Timed models:
    - Timed extensions of finite state automata: timed/hybrid automata  $\Rightarrow$  Lecture 2
    - Timed extensions of Petri nets: time/stopwatch Petri nets  $\Rightarrow$  Lecture 3
- Formalize the specification  $\varphi$  :
  - Observers
  - Temporal logics: LTL, CTL,  $\ldots \Rightarrow$  Lecture 1
  - Timed extensions of temporals logics  $\Rightarrow$  Lectures 2 & 3

• Does  $S \models \varphi$  ?

#### Model-checking algorithms

 $\Rightarrow$  State space exploration

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## Some major issues

## Need for modeling tasks with suspending/resuming features

**Expressivity**/**Decidability** compromise to discuss  $\Rightarrow$  Lectures 2 & 3

#### State space combinatorial explosion

- Need for symbolic approaches  $\Rightarrow$  Lectures 2 & 3
- $\bullet\,$  Need for new models and abstracted algorithms  $\Rightarrow\,$  Lecture 4

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## Context and Aims

**MeForBio** team: Algebraic modeling to study complex dynamical biological systems



1) Two main models

- Historical model: Biological Regulatory Network (René Thomas)
- Recently designed model: Process Hitting

2) Allow efficient translation from one model to the other

## Context and Aims

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#### 1) Two main models

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## Today's issue

### Tricky question

How can we study complex dynamical biological systems, **involving up to 1.000 interacting components**?

### Observation

- Classical model-checking approaches suffer from state space explosion
- Leads:
  - Taking profit for Process Algebra structure, based on a **compact** representation of the interactions
  - Develop **static analysis approaches** to verify some crucial properties, e.g. stable states, reachability, key processes, ...

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# Contribution

### Scientific challenge

How can we cope with the analysis of **large-scale systems**, involving up to 1.000 interacting components?

### Objectives of this talk

- Introduce a Process Algebra inspired framework based on a compact representation of the interactions
- Develop efficient static analysis approaches to answer most common problems
- Apply the methodology to large-scale biological regulatory networks

#### Joint work with

- L. Paulevé (ETH Zurich), M. Folschette, O. Roux (IRCCyN)
- K. Inoue (NII)

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# Short introduction to Biological Regulatory Networks

## Principle of R. Thomas' discrete modeling [TGL76]

- Activations and inhibitions between genes
- Gene/protein couples
- Genes expression is associated to a set of discrete logic levels
- Effective control beyond a given threshold; opposite effect below.

## Interaction graph

- Nodes = Genes
- Directed edges = Interactions
- But what is the evolutionary tendency of a when a is at level 1 and b at level 1? ⇒ Need for parametrization



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Proposed by René Thomas in 1973, several extensions since then

**Historical bio-informatics model** for studying genes interactions Widely used and well-adapted to represent dynamic gene systems



Interaction Graph

### Interaction Graph: structure of the system (genes & interactions)

Nodes: genes
Name a, b, z
Possible values (levels of expression) 0..1, 0..2
Edges: interactions
Threshold 1
Type (activation or inhibition) + / -



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Interaction Graph: structure of the system (genes & interactions)

#### Nodes: genes

•Name *a*, *b*, *z* 

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**Parametrization**: strength of the influences (cooperations)

Maps of tendencies for each gene

- $\rightarrow$  To any influences of predecessors  $\omega$
- $\rightarrow$  Corresponds a **parameter**  $k_{x,\omega}$

 $k_{z,\{a^+,b^+\}} = 2$  means: z tends to 2 when  $a \ge 1$  and b < 1



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Biological Regulatory Network

 $\rightarrow$  All needed information to run the model or study its dynamics:

- Build the State Graph
- Find reachability properties, fixed points, attractors
- Other properties...
- $\rightarrow$  Strengths: well adapted for the study of biological systems
- → **Drawbacks**: inherent complexity; needs the full specification of cooperations

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Modeling biological regulatory networks Thom

Thomas' framework

## Limits of discrete modeling



Figure: Motif inside biological segmentation networks (e.g. drosophila)

(Rmq.: boolean network)

## Limits of discrete modeling



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## Issues related to the synthesis of timed parameters

#### Problems

- Infer the production and degradation rates
- Consider accumulation mechanisms (on/off oscillations)

#### Adaptation of the R. Thomas model

The logic interaction graph is enriched with two functions. These functions give, for any discrete state of the network and for any gene:

- The **production delay** of the gene, depending on its set of resources;
- The degradation delay of the gene, depending on its set of resources.

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## Increasing the models expressivity

### Adapt the model to biological issues

- Introduction of delays  $\Rightarrow$  timed transition systems
- Need for modeling tasks with suspend / resume  $\Rightarrow$  introduce the notion of stopwatches
- Balance between expressivity and decidability

## Problem

### Choosing an appropriate time model for S

- Dense time?
- Discrete time?

### Inference and verification of quantitative temporal properties

 $\Rightarrow$  Efficient state-space exploration algorithms  $\Rightarrow$  Compact data structures for **storage** and **computation** of the state space

## Problem

### Choosing an appropriate time model for S

- Dense time?
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## Inference and verification of quantitative temporal properties

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# Analysis of Biological Regulatory Networks

### Motivation: apply modal logics

- Compute state graph and check properties
- Translate the model into a **discrete-events** (resp. **timed**) model, e.g. (time) Petri nets, and check properties

### Motivation: why Petri nets?

- Mathematical and graphic formalism
- Easy representation concurrence and parallelism
- Structural properties (P-invariants, T-invariants, ...)
- Dynamical properties (liveness, boundedness, reachability, ...)
- Mature tools : Snoopy, ginSIM, ROMÉO, etc.

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 $\{P_1, P_2, P_4\}$ 

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$$\{P_1, P_2, P_4\} \xrightarrow{t_2} \{P_1, P_3, P_4\} \xrightarrow{t_1} \dots$$

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#### Figure: An other Petri net

$$\{P_1, P_2, P_4\}$$

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Figure: An other Petri net

$$\{P_1, P_2, P_4\} \xrightarrow{t_2} \{P_3, P_4\} \dots$$

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## Petri net with reset arcs - Reminder



Figure: A Petri net with reset arcs

$$\{P_1, P_2, 5 \times P_4\} \stackrel{t_2}{\rightarrow} \dots$$

## Petri net with reset arcs - Reminder



Figure: A Petri net with reset arcs

$$\{P_1, P_2, 5 \times P_4\} \xrightarrow{t_2} \{P_1, P_3\} \xrightarrow{t_1} \dots$$

## Petri nets with read arcs



#### Figure: A Petri net with read arcs

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# (Logic) Inhibitor Hyperarc Petri nets : AND inhibition



#### Figure: An inhibitor hyperarc Petri net

$$\{P_1, P_2, P_4\}$$

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## (Logic) Inhibitor Hyperarc Petri nets : AND inhibition



Figure: An inhibitor hyperarc Petri net :  $t_1$  inhibited iff  $(M(P_3) \ge 1$  and  $M(P_4) \ge 1)$ 

$$\{P_1, P_2, P_4\} \xrightarrow{t_2} \{P_1, P_3, P_4\}$$

## (Logic) Inhibitor Hyperarc Petri nets : AND inhibition



Figure: An inhibitor hyperarc Petri net

$$\{P_1, P_2, P_4\} \xrightarrow{t_2} \{P_1, P_3, P_4\} \xrightarrow{t_3} \{P_1, P_4\} \xrightarrow{t_1} \dots$$

## From regulatory networks to Petri nets

#### Principle

- One place per gene
- Marking: discrete level of concentration



#### Critical issues

- How to test the concentration level without decrementing it?
- How to model an action that takes place only **below** a given concentration?
- $\rightarrow$  Use read and inhibitor (hyper)arcs

#### From regulatory networks to Petri nets



#### Figure: Translation towards Petri nets

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## From regulatory networks to Petri nets

#### Analysis

- Automated translation
- **Bounded** networks  $\rightarrow$  reduced cost of read and logic inhibitor hyperarcs

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## From regulatory networks to time extensions of Petri nets

#### Principle

- Thinly **discretize** the concentration levels (thus thresholds) of each gene
- Associate the **production and degradation delays** to transitions resulting from the discrete translation

#### From regulatory networks to time extensions of Petri nets



#### Figure: Translation towards time Petri nets

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## From regulatory networks to time extensions of Petri nets

#### Analysis

- Allows the model-checking of (parametric) TCTL formulae
- Possibility to infer the time parameters associated to a transition
- $\bullet$  Automate the translation and export to  $\operatorname{ROM\acute{E}O}$  software

## Model validation

#### Objective : formal verification of a model properties

- Model the *S* system:
  - $\rightarrow$  Petri nets, Time Petri nets, Stopwatch Petri nets,  $\ldots$
- Formalize the specification  $\varphi$  :
  - $\rightarrow$  observers, timed logics (LTL, CTL, **TCTL**),...
- Does  $S \models \varphi$  ?

Algorithms implemented in  ${
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## Application : p53-MdM2 network



Figure: Network from [WAjK09]

## Application : p53-MdM2 network



#### Figure: Translation into time Petri net

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## **Biological analysis**

#### Model validation

- Verify properties (sustained oscillations, damped oscillations, ...)
- Model-checking of TCTL formulae

#### Delays inference

Model-checking of parametric TCTL formulae

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## **Biological analysis**

#### Limits

- Undecidability of TCTL model-checking, even for bounded parametric TPN [TLR09]
- State space combinatorial explosion
- Limitation in the size of the nets and number of parameters

#### Methodology

- Identification of relevant sub-problems
- Progressive inference of time delays

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## Tools for Interaction Graphs Study

#### Interaction Graphs [RCB05]

- No positive circuit  $\Rightarrow$  only 1 attractor
- No negative circuit  $\Rightarrow$  no cyclic attractor
- Positive circuits  $\Rightarrow$  criterion for max. number of attractors
- Temporal logics  $\Rightarrow$  check properties (needs State Graph)
  - $\rightarrow$  SM-BIONET [KCRB09], ginSIM [CNT12], Biocham [CFS06]
  - $\rightarrow$  Translate models into discrete-event systems and run model-checkers
- Some recent works focus on boolean networks topological fixed points: [PR10]

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#### Problem

Combinatorial explosion when computing the State Graph

 $\rightarrow$  Need for static analysis  $\rightarrow$  introduction of the Process Hitting

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## Intuitive principle of the Process Hitting framework

Process	=	component a at level i
Interaction	=	
		a at level i makes b at level j increase or
		decrease to level k
denoted		$a \rightarrow b \neq b$ (bit and bounce)
		$a_i \rightarrow b_j + b_k$ (int and bounce)

#### Definition (Interaction and Retroaction)

Interaction  $(a_i \rightarrow b_j \stackrel{r}{\vdash} b_k)$ , where  $a_i$  is the level of a process a and  $b_j \neq b_k$ , Retroaction  $(a_i \rightarrow a_i \stackrel{r}{\vdash} a_k)$ : when  $a_i = b_j$ . Process Hitting

Definition

## The Process Hitting modeling



#### **Sorts**: components a, b, z

**Actions**: dynamics  $b_1 \rightarrow z_0 \lor z_1$ ,  $a_0 \rightarrow a_0 \lor a_1$ ,  $a_1 \rightarrow z_1 \lor z_2$ 

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2013/04/17 38 / 62 Process Hitting Definition

## The Process Hitting modeling



**Sorts:** components a, b, z **Processes:** local states / levels of expression  $z_0, z_1, z_2$  **States:** sets of active processes **Actions:** dynamics  $b_1 \rightarrow z_0 \upharpoonright z_1, a_0 \rightarrow a_0 \upharpoonright a_1, a_1 \rightarrow z_1 \upharpoonright z_2$  $a_1 \rightarrow a_1 \rightarrow a_1 \rightarrow a_1 \rightarrow a_2 \rightarrow a_$ 

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How to introduce some **cooperation** between sorts?  $a_1 \wedge b_0 \rightarrow z_1 \uparrow z_2$ Solution: a **cooperative sort** abConstraint: each configuration is represented by one process  $\langle a_1, b_0 \rangle \Rightarrow ab_{10}$ Advantage: regular sort; drawbacks: complexity, temp@al shift.  $z \in z = 200$ M. MAGNIN (IRCCyN-NII) Lecture Series - Lecture 4 / NII 2013/04/17 39 / 62



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## Adding cooperations [PMR12]



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### Static Analysis: Fixed Points [PMR11a]

**Fixed point** = state where no action can be fired

 $\rightarrow$  avoid couples of processes bounded by an action



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### Static Analysis: Fixed Points [PMR11a]

**Fixed point** = state where no action can be fired

- $\rightarrow$  avoid couples of processes bounded by an action
- $\rightarrow$  Hitless Graph





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Exponential complexity w.r.t. the number of sorts

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#### Static analysis: successive reachability [PMR12]

#### Problem

Given an initial state of a Process Hitting, is it possible to reach successively  $a_i$ , then  $b_i$ , then  $a_k$ , then  $c_1, \ldots$ ?  $\Rightarrow$  Combinatorial explosion of the dynamics to explore

#### Key idea

Instead of checking the successive reachability  $\mathcal{R}$ , which is complex, we will check:

- an under-approximation  $\mathcal{P}$ : if  $\mathcal{P}$  is not satisfied, then  $\mathcal{R}$  neither
- an over-approximation Q: if Q is satisfied, then  $\mathcal{R}$  too.

EL OQO

# Static analysis: successive reachability [PMR12]



 $\begin{array}{c} \rightarrow \text{ Concretization of the objective} = \text{scenario} \\ a_0 \rightarrow c_0 \mathrel{\upharpoonright} c_1 :: \quad b_0 \rightarrow d_0 \mathrel{\upharpoonright} d_1 :: \quad c_1 \rightarrow b_0 \mathrel{\upharpoonright} b_1 :: \quad b_1 \rightarrow d_1 \mathrel{\upharpoonright} d_2 \end{array}$ 

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= 200

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 $\begin{array}{c} \rightarrow \text{ Concretization of the objective} = \text{scenario} \\ a_0 \rightarrow c_0 \mathrel{\sc l} c_1 :: \quad b_0 \rightarrow d_0 \mathrel{\sc l} d_1 :: \quad c_1 \rightarrow b_0 \mathrel{\sc l} b_1 :: \quad b_1 \rightarrow d_1 \mathrel{\sc l} d_2 \end{array}$ 

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= 200

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 $\begin{array}{c} \rightarrow \text{ Concretization of the objective} = \text{scenario} \\ a_0 \rightarrow c_0 \mathrel{\mbox{$\stackrel{f}{$}$}} c_1 :: \ b_0 \rightarrow d_0 \mathrel{\mbox{$\stackrel{f}{$}$}} d_1 :: \ c_1 \rightarrow b_0 \mathrel{\mbox{$\stackrel{f}{$}$}} b_1 :: \ b_1 \rightarrow d_1 \mathrel{\mbox{$\stackrel{f}{$}$}} d_2 \end{array}$ 

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 $\begin{array}{c} \rightarrow \text{ Concretization of the objective} = \text{scenario} \\ \underline{a_0 \rightarrow c_0 \restriction' c_1} :: \quad b_0 \rightarrow d_0 \restriction' d_1 :: \quad c_1 \rightarrow b_0 \restriction' b_1 :: \quad b_1 \rightarrow d_1 \restriction' d_2 \end{array}$ 

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 $\begin{array}{rl} \rightarrow \text{ Concretization of the objective} = \text{scenario} \\ a_0 \rightarrow c_0 \mathrel{\baseline rel} c_1 :: & b_0 \rightarrow d_0 \mathrel{\baseline rel} d_1 :: & \underline{c_1 \rightarrow b_0} \mathrel{\baseline rel} b_1 \mathrel{\baseline rel} :: & b_1 \rightarrow d_1 \mathrel{\baseline rel} d_2 \end{array}$ 

# Static analysis: successive reachability [PMR12]



 $\begin{array}{c} \rightarrow \text{ Concretization of the objective} = \text{scenario} \\ a_0 \rightarrow c_0 \mathrel{\upharpoonright} c_1 \mathrel{::} \hspace{0.2cm} b_0 \rightarrow d_0 \mathrel{\upharpoonright} d_1 \mathrel{::} \hspace{0.2cm} c_1 \rightarrow b_0 \mathrel{\upharpoonright} b_1 \mathrel{::} \hspace{0.2cm} \underline{b_1 \rightarrow d_1 \mathrel{\upharpoonright} d_2} \end{array}$ 

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- ightarrow Directly checking an objective sequence R is hard
- $\rightarrow$  Rather check the approximations *P* and *Q*, where *P*  $\Rightarrow$  *R*  $\Rightarrow$  *Q*:



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Static analysis by abstractions:

- ightarrow Directly checking an objective sequence R is hard
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Linear w.r.t. the number of sorts and exponential w.r.t. the number of processes in each sort

 $\rightarrow$  Efficient for big models with few levels of expression



 $\rightarrow$  New abstract structure **Sufficient condition**:





 $\rightarrow$  New abstract structure Sufficient condition:

- no cycle
- each objective has a solution





### $\rightarrow$ New abstract structure **Sufficient condition**:

- no cycle
- each objective has a solution

### R is true





 $\rightarrow$  New abstract structure **Sufficient condition**:

- no cycle
- each objective has a solution



= 900



 $\rightarrow$  New abstract structure **Sufficient condition**:

- no cycle
- each objective has a solution

### Inconclusive





**Necessary condition**:



1= 990



#### Necessary condition:

There exists a traversal with no cycle

- $\bullet~$  objective  $\rightarrow~$  follow one solution
- $\bullet$  solution  $\rightarrow$  follow all processes
- process  $\rightarrow$  follow all objectives





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#### **Over-approximation**



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- $\bullet$  solution  $\rightarrow$  follow all processes
- $\bullet~\mbox{process} \to \mbox{follow}$  all objectives



#### **Over-approximation**



Necessary condition:

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#### **Over-approximation**



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There exists a traversal with no cycle

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- Introduces time features
- Parameters: either (r, sa), or the firing interval [d; D].

ightarrow Tests by simulation and model-checking

- Introduces time features
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#### Refining with Actions Removal

#### Prevent behaviors by deleting unrealistic actions



#### Refining with Actions Removal

#### Prevent behaviors by deleting unrealistic actions



#### Refining with Actions Removal

#### Prevent behaviors by deleting unrealistic actions



- How to express  $(a_1 \wedge b_1) \rightarrow z_0 \stackrel{\scriptstyle{\uparrow}}{\rightarrow} z_1$ ?
- $\rightarrow$  Add a **cooperative sort** reflecting the state of *a* and *b*



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- How to express  $(a_1 \wedge b_1) \rightarrow z_0 \lor z_1$ ?
- $\rightarrow$  Add a **cooperative sort** reflecting the state of *a* and *b*



Allow cooperation between two genes

- How to express  $(a_1 \wedge b_1) \rightarrow z_0 \stackrel{\scriptscriptstyle ?}{\vdash} z_1$ ?
- $\rightarrow$  Add a **cooperative sort** reflecting the state of *a* and *b*



 $\rightarrow$  Introduces a temporal shift (over-approximation)

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# Using Process Hitting for Interaction Graphs Study

#### Motivation

- Interaction Graph is the **historical discrete model** (suitable and widespread in biological research)
- Several tools exist of the analysis of interaction graphs, but the state graph is needed for some results ⇒ combinatorial explosion

# Contribution: Process Hitting to study large Biological Regulatory Networks

- Translation from Interaction Graphs + Refining
- Efficient static analysis

# The Process Hitting modeling

#### Key features

- Dynamic modeling with an atomistic point of view
  - $\rightarrow$  Independent actions
  - $\rightarrow$  Cooperation modeled with cooperative sorts
- Efficient static analysis
  - $\rightarrow$  Reachability of a process can be computed in **linear time** in the number of sorts
- Useful for the study of large biological models
  - $\rightarrow$  Up to hundreds of sorts

#### (Future) extensions

- Actions with stochasticity
- Actions with priorities
- Continuous time with clocks?

#### The Pint Tool

[http://processhitting.wordpress.com]

#### Features

- Free software (API available for future developments)
- Textual language to describe a Process Hitting (GUI currently under development)

#### • Implemented tools:

- Translations from and to various other models
- Fixed points research
- Stochastic simulation
- Reachability checker

#### The Pint Tool

[http://processhitting.wordpress.com]

#### **Results and performance** (reachability analysis):

Model	sorts	procs	actions	states	Biocham <sup>1</sup>	libddd <sup>2</sup>	PINT
egfr20	35	196	670	2 <sup>64</sup>	[3s-KO]	[1s-150s]	0.007s
tcrsig40	54	156	301	2 <sup>73</sup>	[1s-KO]	[0.6s-KO]	0.004s
tcrsig94	133	448	1124	2 <sup>194</sup>	KO	KO	0.030s
egfr104	193	748	2356	2 <sup>320</sup>	KO	KO	0.050s

<sup>1</sup> [Inria Paris-Rocquencourt/Contraintes] <sup>2</sup> [LIP6/Move]

# The Mobyle portal

[http://cardioserve.nantes.inserm.fr/cgi-bin/mobyle/portal.py]

#### Presentation

- Web application unifying tools for systems biology analysis
- Powered by the Mobyle framework
- Project led in the context of the French ANR "BIOTempo" project



Figure: General architecture of the BIOtempo Mobyle server

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#### Presentation

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Figure: Screenshot from the BIOtempo Mobyle server: cardioserve.nantes.inserm.fr/cgi-bin/mobyle/portal.py

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## Overview

#### Introduction

#### 2 Modeling biological regulatory networks

- Thomas' framework
- From Thomas' framework to discrete-event systems
- From Thomas' framework to timed systems
- Common limits of current models for biological analyses

3 The Process Hitting: a framework well suited to concurrent systems

- Definition
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Inferring information on the biological model thanks to the Process Hitting

- Interaction Graph Inference
- Parametrization Inference
- Summary & Conclusion

Information inference

#### Inferring a BRN with Thomas' parameters







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Information inference

#### Inferring a BRN with Thomas' parameters



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Information inference

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Information inference Interaction Graph Inference

# Inferring the Interaction Graph [FPI<sup>+</sup>12]





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Information inference Interaction Graph Inference

# Inferring the Interaction Graph [FPI<sup>+</sup>12]





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# Inferring the Interaction Graph [FPI<sup>+</sup>12]





 $\rightarrow$  Exhaustive search in all possible configurations

- Pick one regulator [a], and choose an active process for all the others [b<sub>0</sub>].
- 2. Change the active process of this regulator [*a*<sub>0</sub>, *a*<sub>1</sub>] and watch the **focal processes**.
- 3. Conclude locally:  $(a_0 \lor a_1 \Rightarrow z_0 \lor z_2)$  $\Rightarrow$  activation (+) & threshold = 1.
- 4. Iterate

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- 3. Conclude locally:  $(a_0 \upharpoonright a_1 \Rightarrow z_0 \vDash z_2)$  $\Rightarrow$  activation (+) & threshold = 1.
- 4. Iterate and conclude globally.

### Inferring the Interaction Graph [FPI<sup>+</sup>12]





#### Problematic cases:

- → No focal processes (cycle) → Opposite influences (+ & -)  $\}$  ⇒ Unsigned edge



- 1. For each configuration of resources  $[\omega = \{a^+, b^-\}]$ find the **focal processes**. If possible, conclude.  $[k_{z,\{a^+,b^-\}} = 1]$ Inconclusive cases:
  - Behavior cannot be represented as a BRN
  - Lack of cooperation (no focal processes)



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### Inferring Parameters



- 1. For each configuration of resources  $[\omega = \{a^+, b^-\}]$ find the **focal processes**. If possible, conclude.  $[k_{z,\{a^+,b^-\}} = 1]$ Inconclusive cases:
  - Behavior cannot be represented as a BRN
  - Lack of cooperation (no focal processes)
- 2. If some parameters could not be inferred, enumerate all admissible parametrizations, regarding **biological constraints** and **the dynamics** of the Process Hitting  $\Rightarrow k_{z,\{a^+,b^-\}} \in \{0;1;2\}; k_{z,\{a^-,b^+\}} \in \{0;1;2\}$

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### Implementation

Workflow:

- Read and translate the models with **OCaml** 
  - $\rightarrow$  Uses the existing free library Pint
  - $\rightarrow$  Documentation + examples:

http://processhitting.wordpress.com/

- Express the problem in ASP (logic programming)
  - $\rightarrow$  Solve with Clingo (Gringo + Clasp)

Model specifications				IG inference		Parameters inference	
Name	S+CS	Р	А	$\Delta t$	Edges	$\Delta t$	Paramet
[EGFR20]	<b>20</b> +22	152	399	1s	50	1s	191
[TCRSIG40]	<b>40</b> +14	156	301	1s	54	1s	143
[TCRSIG94]	<b>94</b> +39	448	1124	13s	169	$\infty$	2.10 <sup>9</sup>
[EGFR104]	<b>104</b> +89	748	2356	4min	241	1min 30s	1.10 <sup>6</sup> /2.
S = Sorts CS	= Cooperat	ive sor	ts P =	= Process	ses A =	= Actions	

 EGFR20]: Epidermal Growth Factor Receptor, by Özgür Sahin et al.

 EGFR104]: Epidermal Growth Factor Receptor, by Regies Samagazet al = 아직 ~

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[EGFR20]	<b>20</b> +22	152	399	1s	50	1s	191		
[TCRSIG40]	<b>40</b> +14	156	301	1s	54	1s	143		
[TCRSIG94]	<b>94</b> +39	448	1124	13s	169	$\infty$	2.10 <sup>9</sup>		
[EGFR104]	<b>104</b> +89	748	2356	4min	241	1min 30s	$1.10^{6}/2.$		
S = Sorts $CS = Cooperative sorts$ $P = Processes$ $A = Actions$									

 [EGFR20]: Epidermal Growth Factor Receptor, by Özgür Sahin et al.

 [EGFR104]: Epidermal Growth Factor Receptor, by Regina Samaga et al.

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### Summary

#### Process Hitting and ASP

- Inference of the complete Interaction Graph
- Inference of the possibly partial Parametrization
- Enumerate all full & admissible Parametrizations
  - $\rightarrow$  Exhaustive approaches

**Complexity**: linear in the number of genes, exponential in the number of regulators of one gene

### Summary

#### Contribution: new translation Process Hitting ~>> René Thomas

- $\rightarrow$  New formal link between the two models
- $\rightarrow$  More visibility to the Process Hitting
- $\rightarrow$  Inference approach that takes benefit from both the Process Hitting compact structure and the power of ASP

### Further work

#### Models and algorithms

- Add priorities in the Process Hitting framework and adapt the static analyses approaches for this enriched model (⇒ paper currently submitted at CS2Bio'13)
- From priorities to quantitative timing information
- Connect Process Hitting compact structure with decomposition techniques in continuous approaches [ACC12] (⇒ paper currently submitted at CMSB'13)

#### Application

- Use the approach for the analysis of larger biological networks
- Contribute to the **discovery** of biological regulatory networks based on biological data
- Study key properties (e.g. concept of resilience)

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